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# Cisco Software-Defined Access Technology Deep Dive

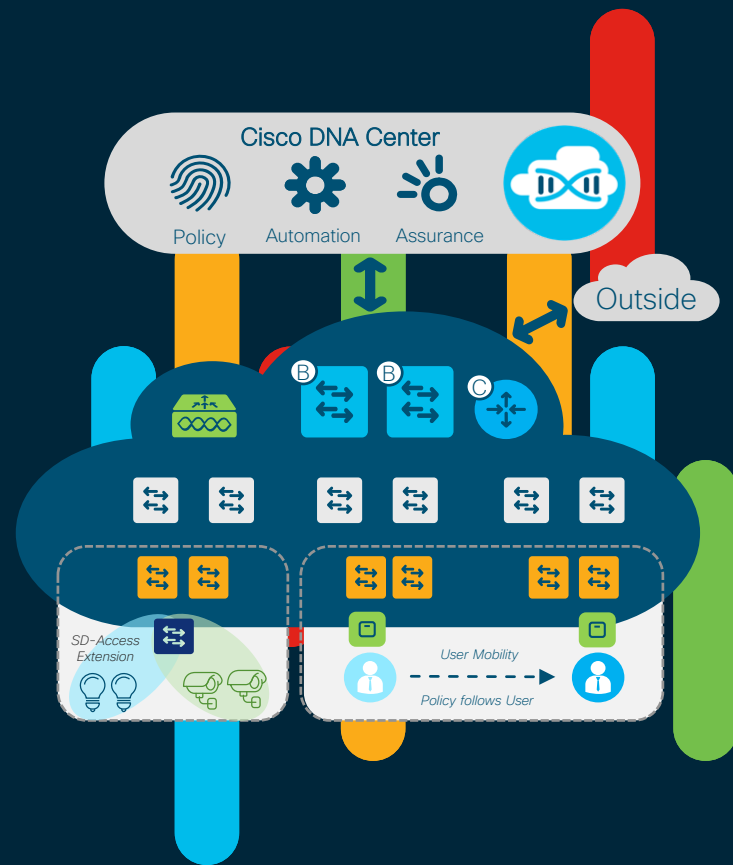
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Distinguished Engineer - Technical Marketing

BRKCRS-3810

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Barcelona | January 27-31, 2020



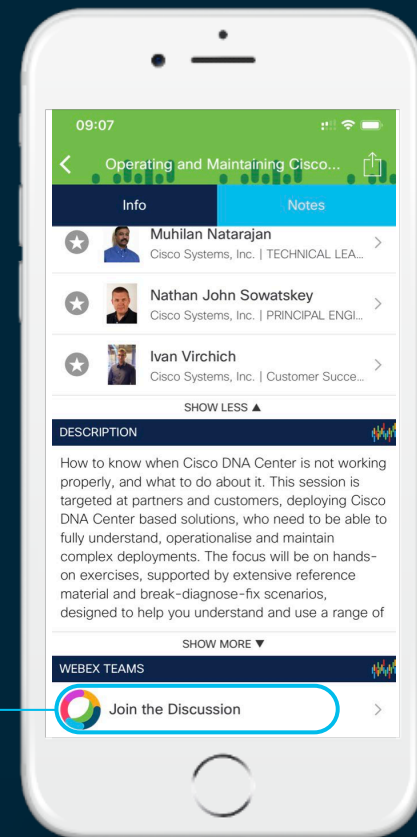
# Cisco Webex Teams

## Questions?

Use Cisco Webex Teams to chat with the speaker after the session

## How

- 1 Find this session in the Cisco Events Mobile App
- 2 Click “Join the Discussion”
- 3 Install Webex Teams or go directly to the team space
- 4 Enter messages/questions in the team space



# Assumptions



This session assumes you have received Cisco DNA, SD-Access & ISE Training

If not... please complete one or all of the following training materials:

- [CiscoLive](#)
- [dCloud Lab](#)
- [Learning@Cisco](#)
- [SDA Design CVD](#)
- [SDA Deploy CVD](#)
- [DNAC Guides](#)

This session is based on Cisco DNAC / SDA 1.3.1, ISE 2.6 and IOS-XE 16.11

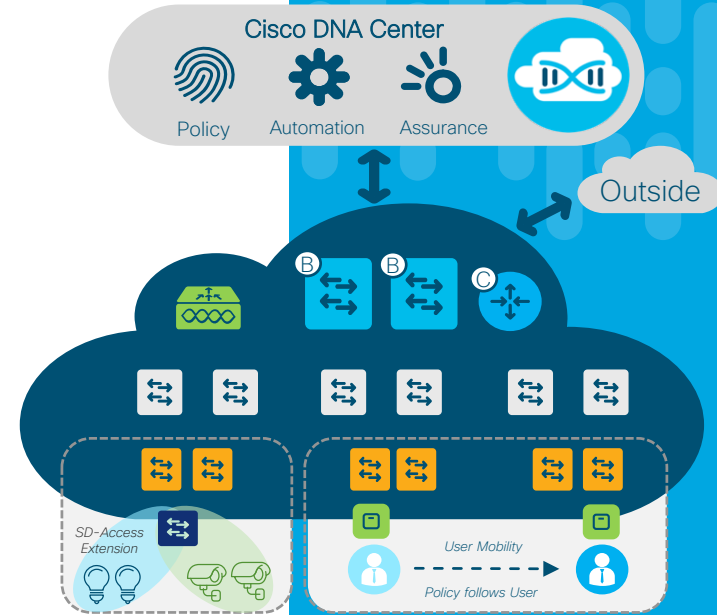
- [Product Compatibility Matrix](#)

For a full list of current capabilities, restrictions, limitations & caveats refer to:

- [DNAC Release Notes](#)

# Agenda

- Cisco SD-Access Recap
- Host Onboarding
  - Endpoint Classification
  - DHCP in SD-Access Fabric
- Connectivity & Access Control
  - Unicast Forwarding
  - Access Control Policy
- Advanced Topics
  - Multicast Forwarding
  - Broadcast Forwarding



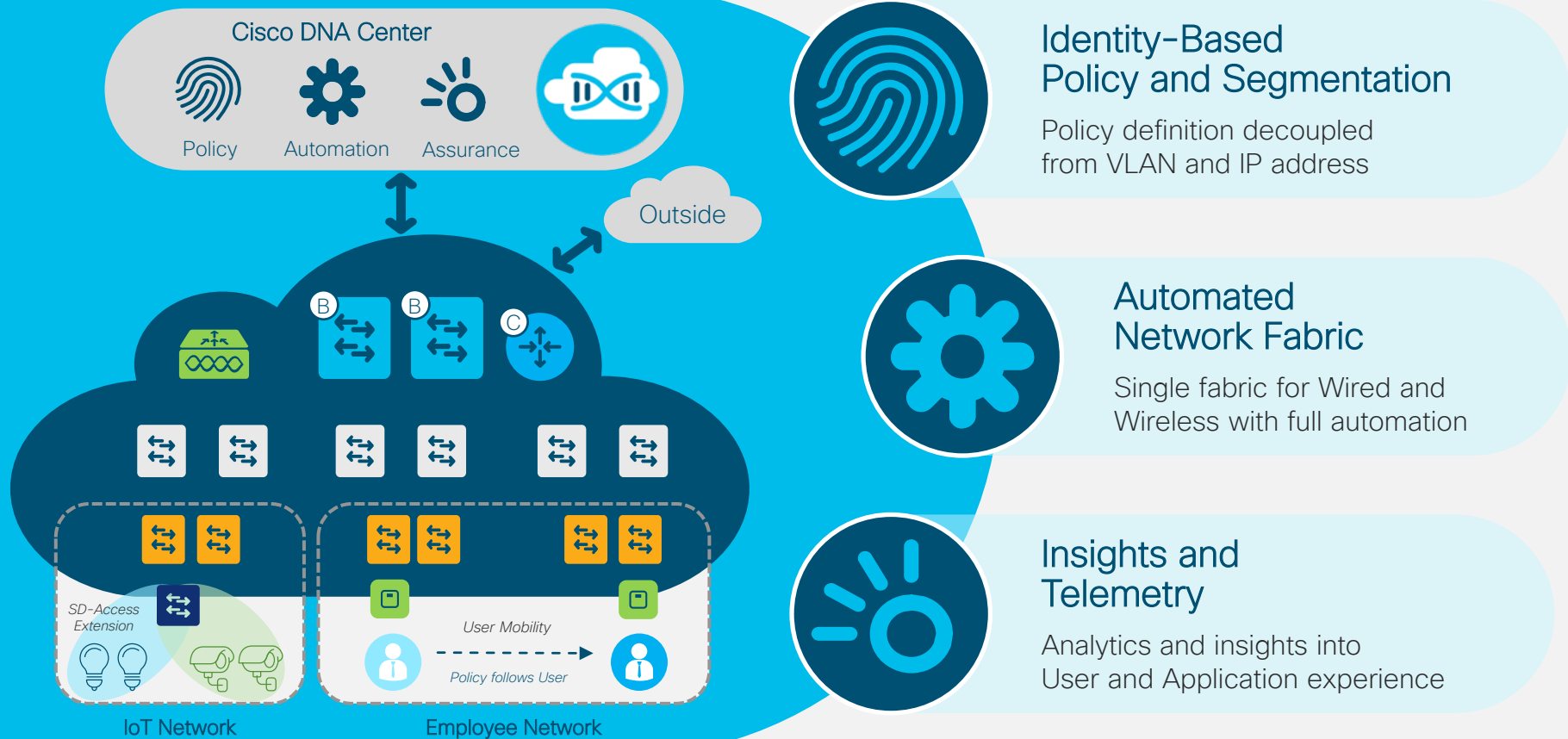
What is Campus Fabric?

# Fabric Fundamentals

1. Architecture
2. Key Components
3. Fabric Constructs

# Cisco Software Defined Access

## The Foundation for Cisco's Intent-Based Network



### Identity-Based Policy and Segmentation

Policy definition decoupled from VLAN and IP address

### Automated Network Fabric

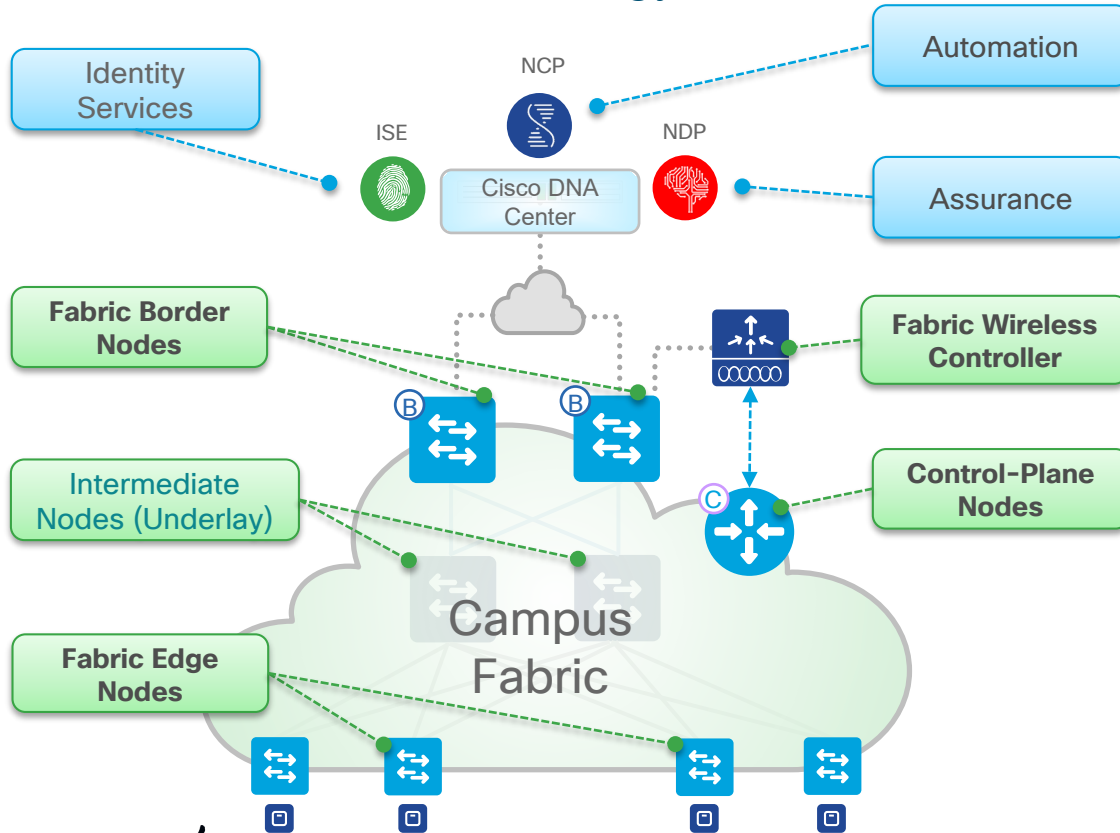
Single fabric for Wired and Wireless with full automation

### Insights and Telemetry

Analytics and insights into User and Application experience

# SD-Access Architecture

## Fabric Roles & Terminology



- **Network Automation** – Simple graphical user interface and intent based automation (e.g. NCP) of fabric devices
- **Network Assurance** – Data Collectors (e.g. NDP) analyze Endpoint to App flows and monitor fabric status
- **Identity Services** – NAC & ID Systems (e.g. ISE) for dynamic Endpoint to Group mapping and Policy definition
- **Control-Plane Nodes** – Map System that manages Endpoint to Device relationships
- **Fabric Border Nodes** – A Fabric device (e.g. Core) that connects External L3 network(s) to the SD-Access Fabric
- **Fabric Edge Nodes** – A Fabric device (e.g. Access or Distribution) that connects Wired Endpoints to the SD-Access Fabric
- **Fabric Wireless Controller** – A Fabric device (WLC) that connects APs and Wireless Endpoints to the SD-Access Fabric

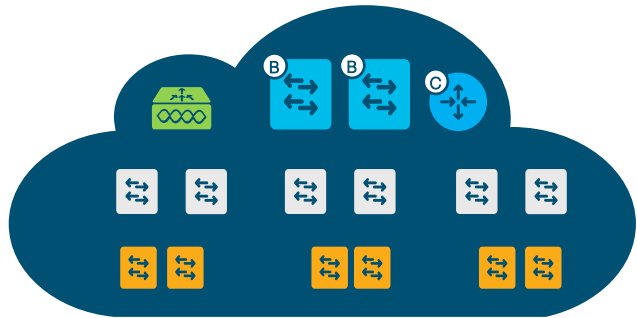


# SD-Access Fabric

## Campus Fabric - Key Components



1. **Control-Plane** based on LISP
2. **Data-Plane** based on VXLAN
3. **Policy-Plane** based on CTS

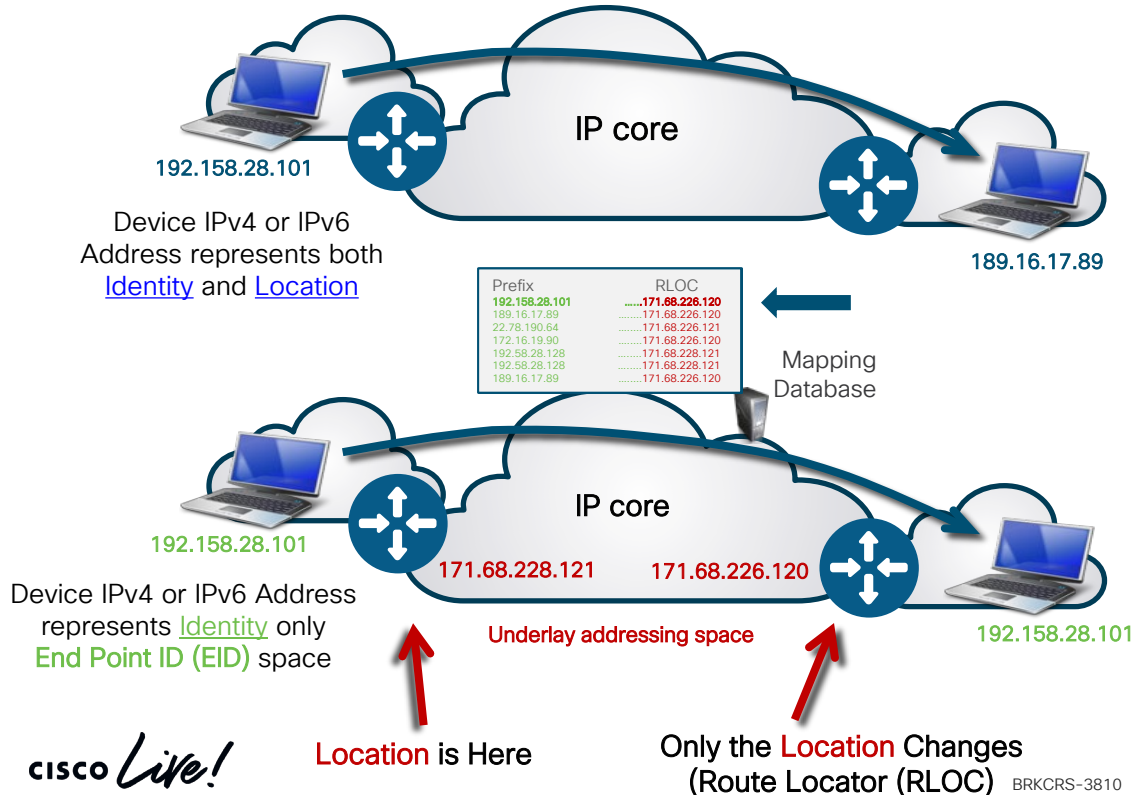


### Key Differences

- L2 + L3 Overlay -vs- L2 or L3 Only
- Host Mobility with Anycast Gateway
- Adds VRF + SGT into Data-Plane
- Virtual Tunnel Endpoints (Automatic)
- NO Topology Limitations (Basic IP)

# Locator / ID Separation Protocol

## Location and Identity separation



### Traditional Behavior - Location + ID are "Combined"

When the Device moves, it gets a new IPv4 or IPv6 Address for its new Identity and Location

### Overlay Behavior - Location & ID are "Separated"

When the Device moves, it keeps the same IPv4 or IPv6 Address. It has the Same Identity

# SD-Access Fabric

## LISP Control Plane

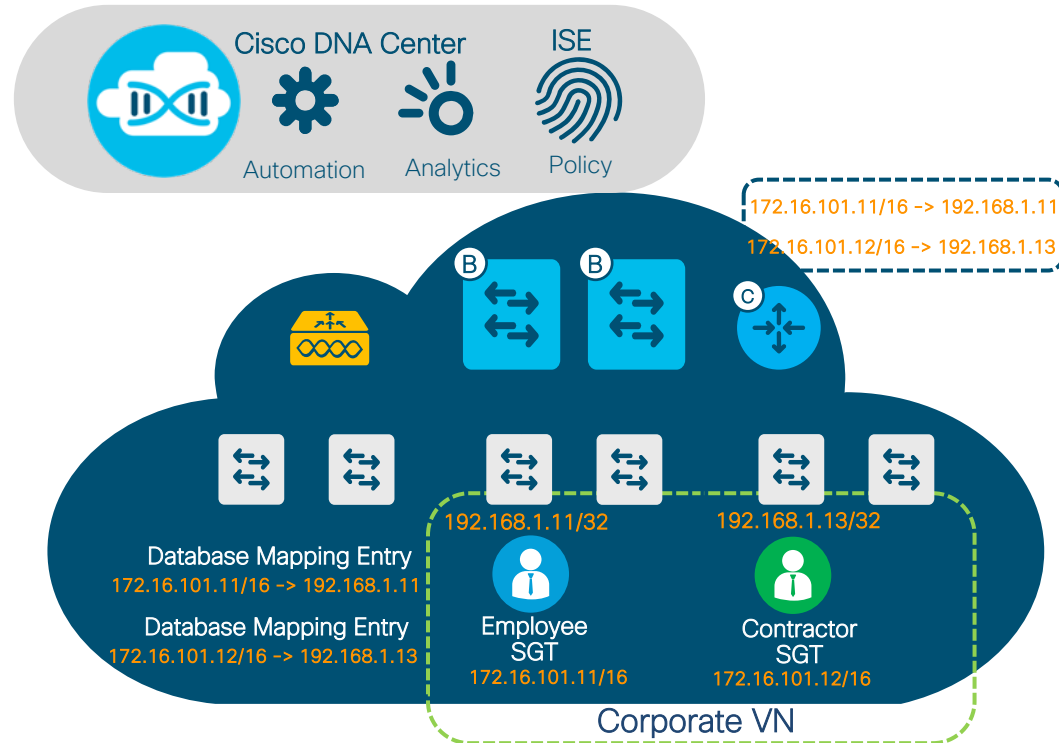


Fabric nodes use LISP as a control plane for Endpoint Identifier (EID) and Routing Locator (RLOC) info

Fabric Control Plane node acts as a Map Server / Resolver for EID to RLOC mappings

Fabric Edge and Internal Border devices registers EIDs to the Map Server.

External Border node acts as PXTR (LISP Proxy Tunnel Router) and provides default gateway when no mapping exists.



# SD-Access Fabric

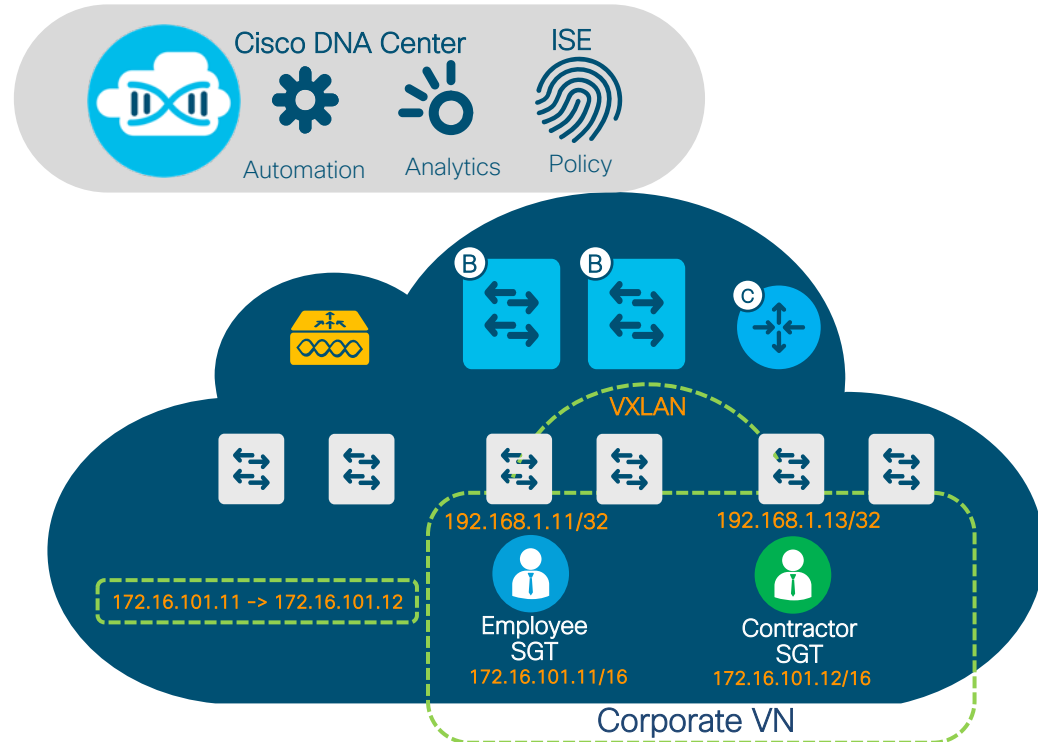
## VXLAN Data Plane



Fabric nodes use VXLAN (Ethernet Based) as the data plane which supports both L2 and L3 overlay.

VXLAN header contains VNID (VXLAN Network Identifier) field which allows up to 16 million Virtual Networks (Layer 3 (VRFs) or Layer 2).

VXLAN header also has Group Policy ID, or Scalable Group Tags (SGTs), allowing 64,000 SGTs.



# SD-Access Fabric

## Cisco TrustSec Policy Plane

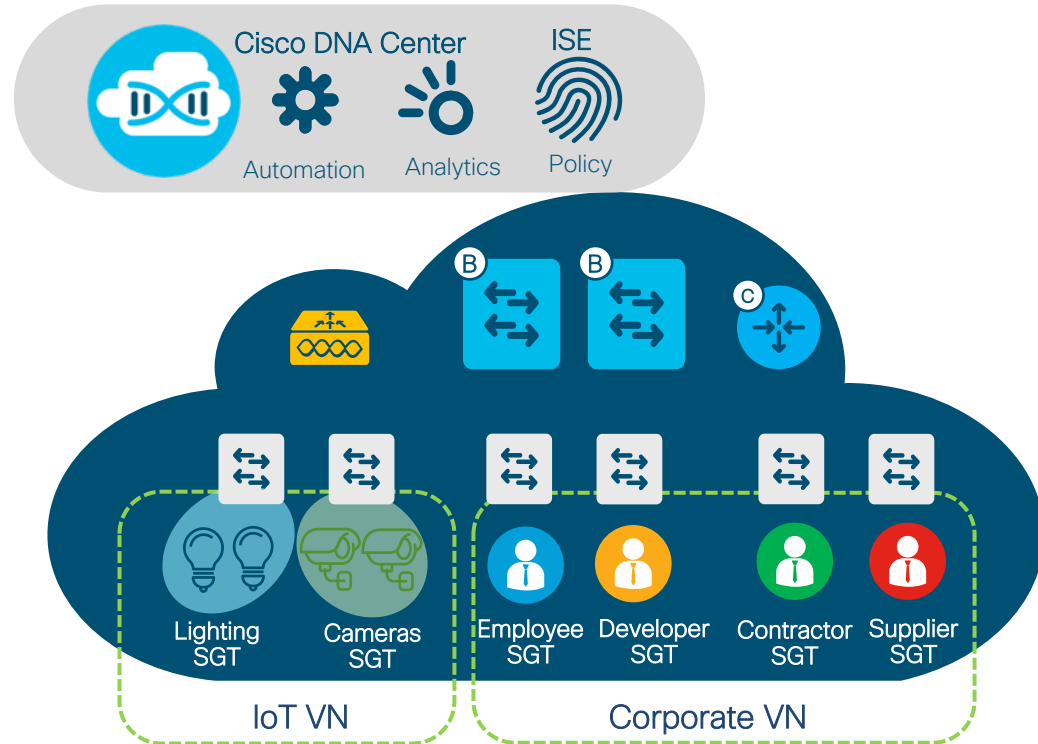


Scalable Groups (SG) are a logical construct defined/identified based on the user and/or device context.

ISE dynamically assign SGs to the users and devices coming to the network fabric.

Nodes add Scalable Group Tags (SGTs) to the fabric encapsulation when communicating between the users and devices.

Edge and border nodes enforce the SGACL policies for the SGTs they protect locally.

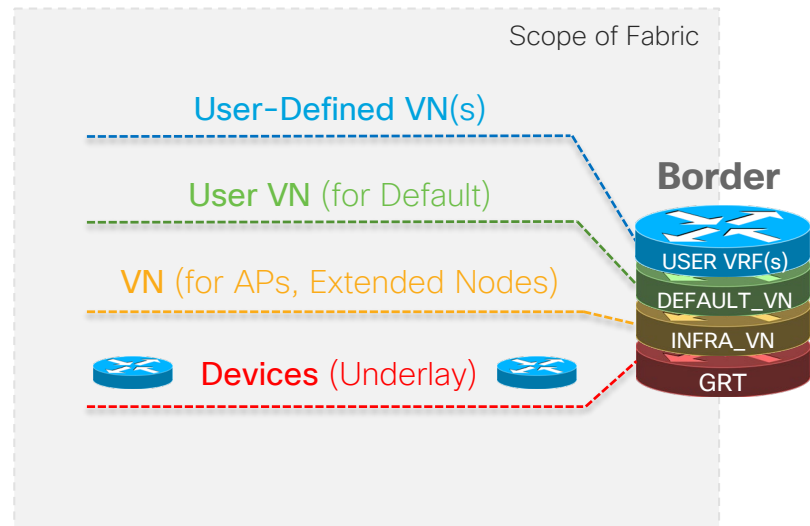


# SD-Access Fabric

## How VNs work in SD-Access



- **Fabric Devices (Underlay)** connectivity is in the **Global Routing Table**
- **INFRA\_VN** is only for **Access Points** and **Extended Nodes** in GRT
- **DEFAULT\_VN** is an actual “**User VN**” provided by default
- **User-Defined VNs** can be added or removed on-demand

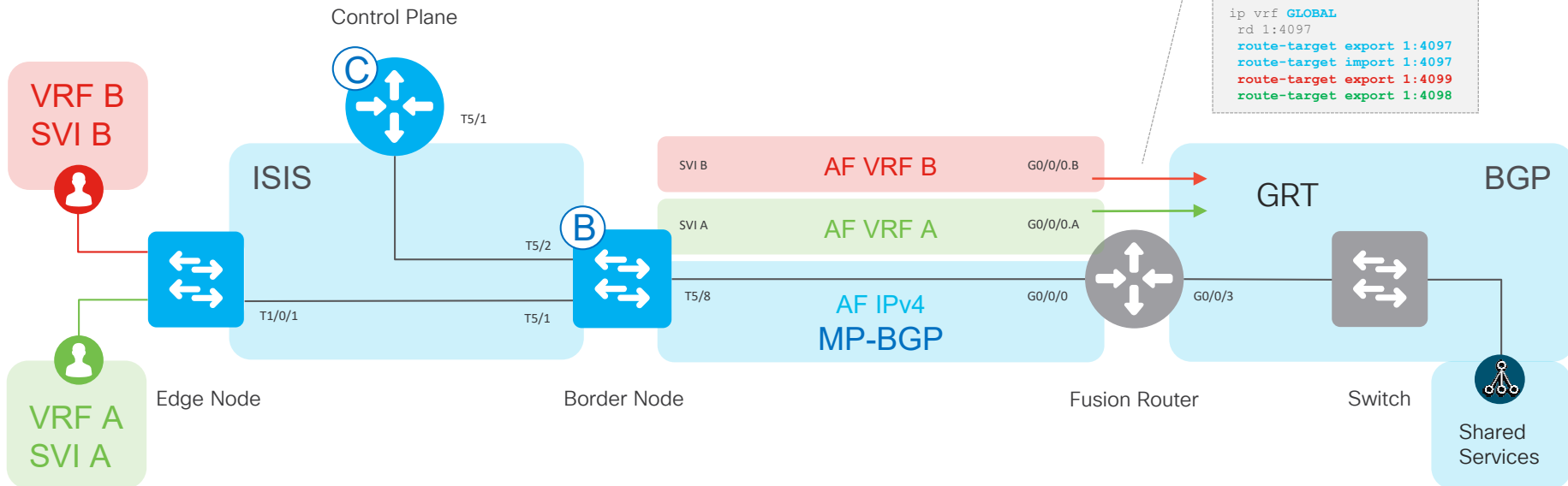


# SD-Access Fabric

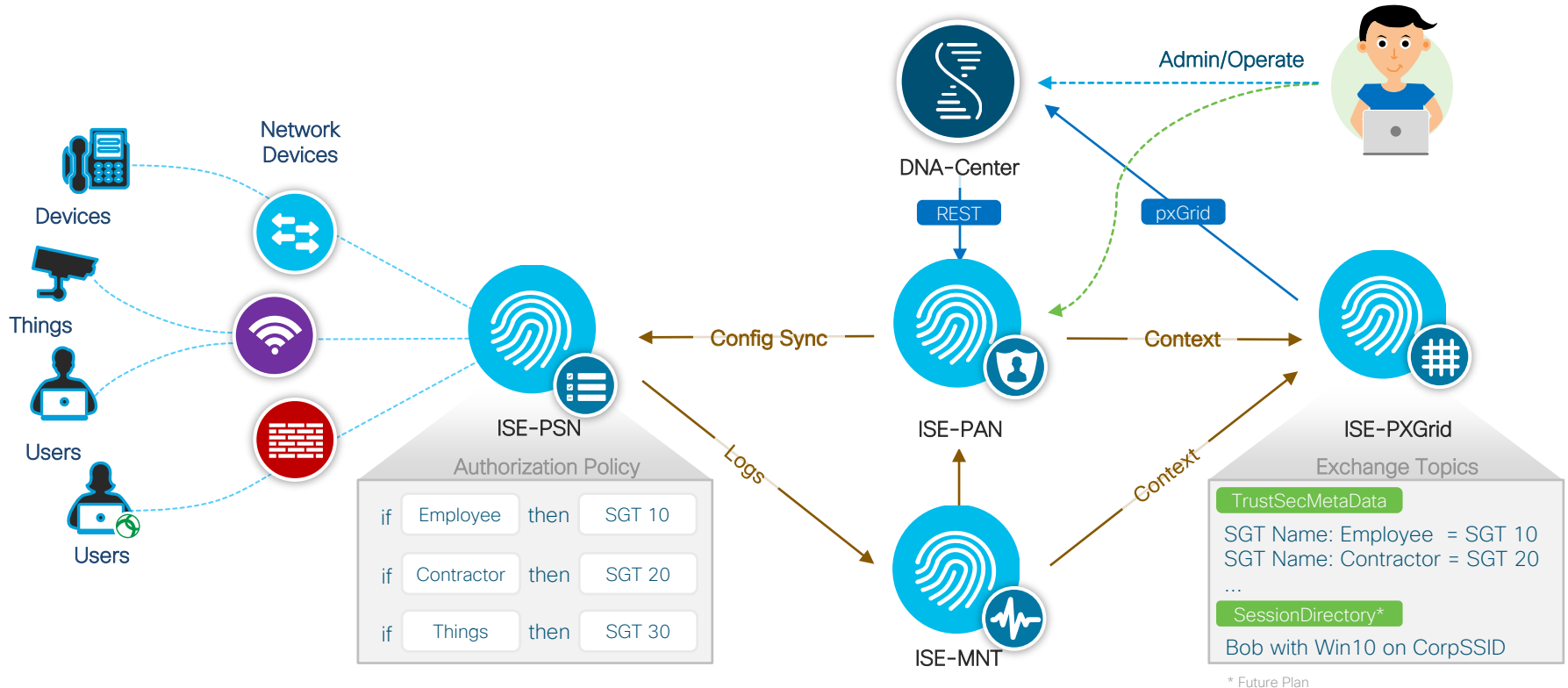
## How VNs work in SD-Access



SD-Access Designs connecting to existing Global Routing Table should use a “Fusion” router with MP-BGP & VRF import/export.



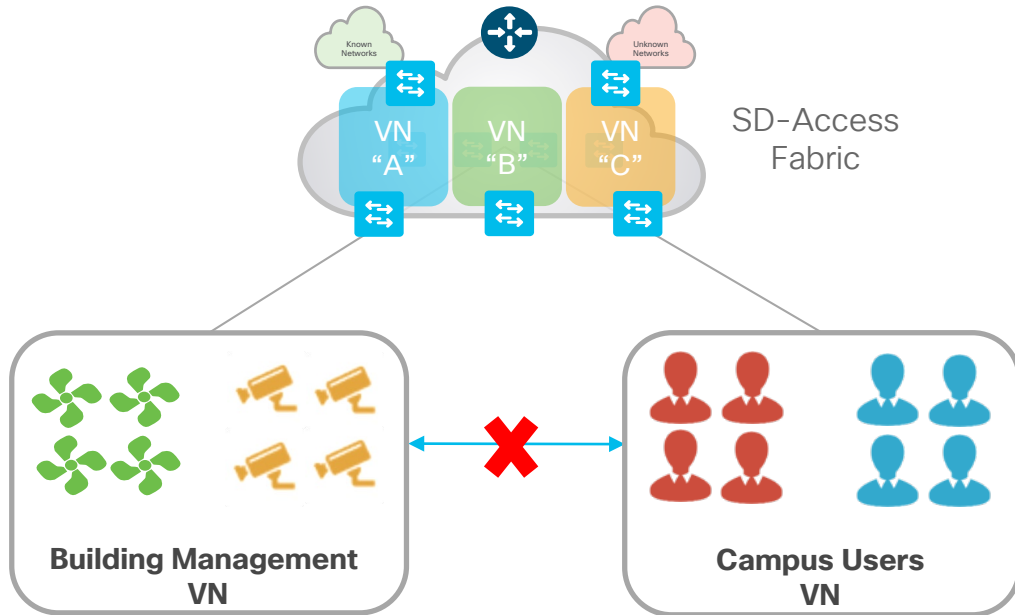
# ISE – Cisco DNA Center Operation





# SD-Access Policy

## Two Level Hierarchy - Macro Level

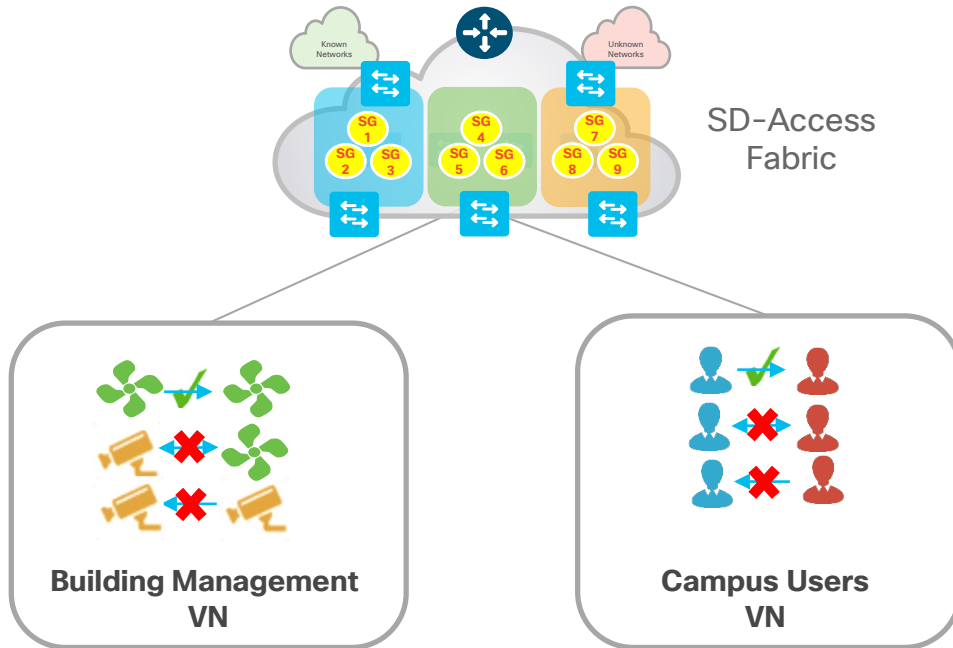


### Virtual Network (VN)

First level Segmentation ensures **zero communication** between forwarding domains. Ability to consolidate multiple networks into one management plane.

# SD-Access Policy

## Two Level Hierarchy - Micro Level



### Scalable Group (SG)

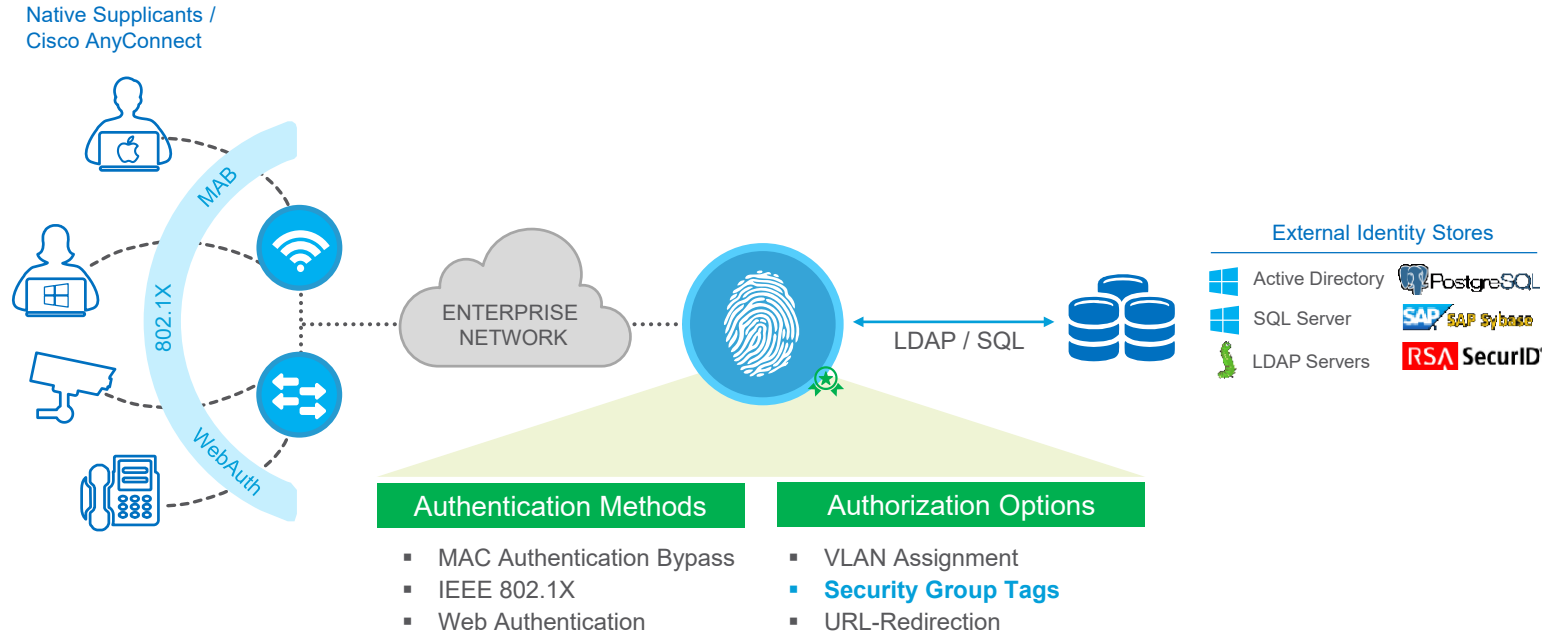
Second level Segmentation ensures **role based access control** between two groups within a Virtual Network. Provides the ability to segment the network into either line of businesses or functional blocks.

Classification and Group tagging

# Host Onboarding

1. Endpoint Classification
2. DHCP in Fabric

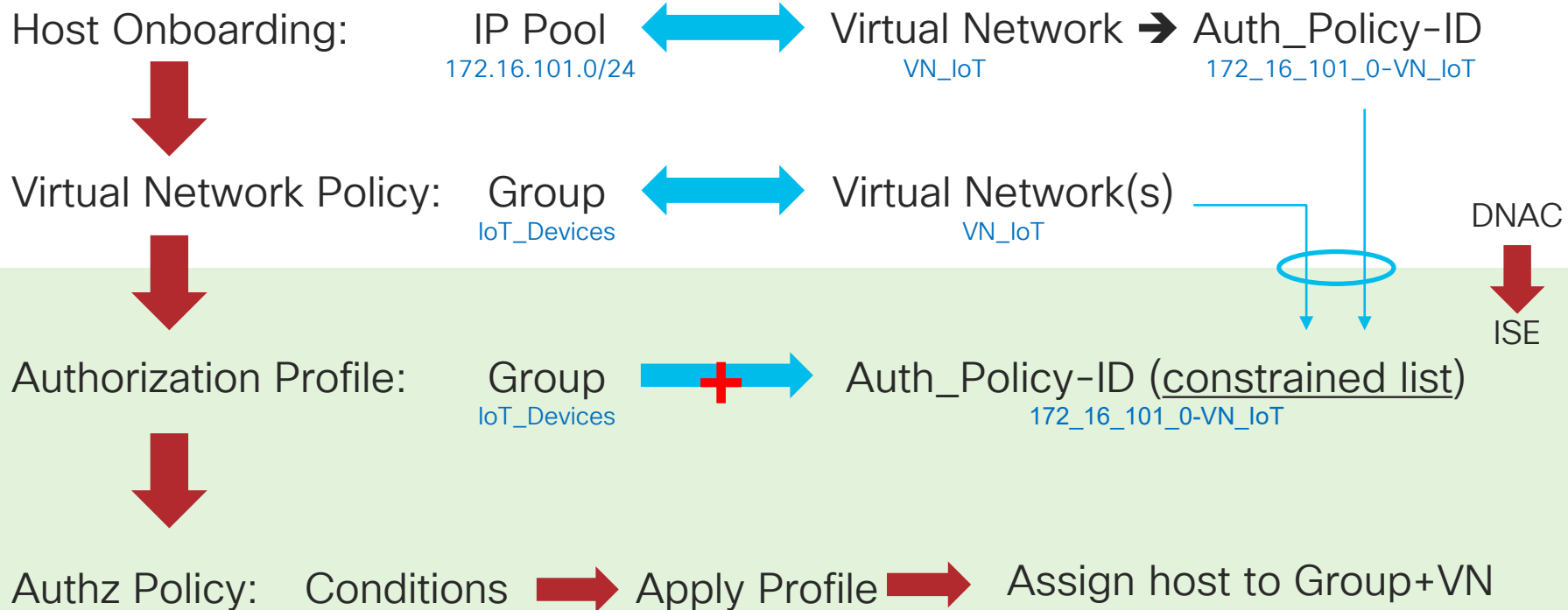
# Authentication and Authorization



Who are you?

What can you do?

# Authorization Workflow for Segmentation



# Authorization Policy – Authorization Profiles



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Policy Sets Profiling Posture Client Provisioning Policy Elements

Dictionarys Conditions Results

- Authentication
- Authorization
  - Authorization Profiles
  - Downloadable ACLs
- Profiling
- Posture
- Client Provisioning

Authorization Profiles > VN\_IoT\_w\_SGT

**Authorization Profile**

\* Name: VN\_IoT\_w\_SGT

Description: [Empty]

\* Access Type: ACCESS\_ACCEPT

Network Device Profile: Cisco

Service Template: [Unchecked]

Track Movement: [Unchecked]

Passive Identity Tracking: [Unchecked]

**Common Tasks**

- DAACL Name: [Unchecked]
- ACL (Filter-ID): [Unchecked]
- Security Group**: IOT\_Devices Virtual Network: VN\_IoT
- VLAN: [Unchecked]

**Advanced Attributes Settings**

Select an item = [Empty] - +

**Attributes Details**

```
Access Type = ACCESS_ACCEPT
cisco-av-pair = cts:security-group-tag=0010-0
cisco-av-pair = cts:sgt-name=IOT_Devices
cisco-av-pair = cts:vn= VN_IoT
Tunnel-Private-Group-ID = 1:
Tunnel-Type = 1:13
Tunnel-Medium-Type = 1:6
```

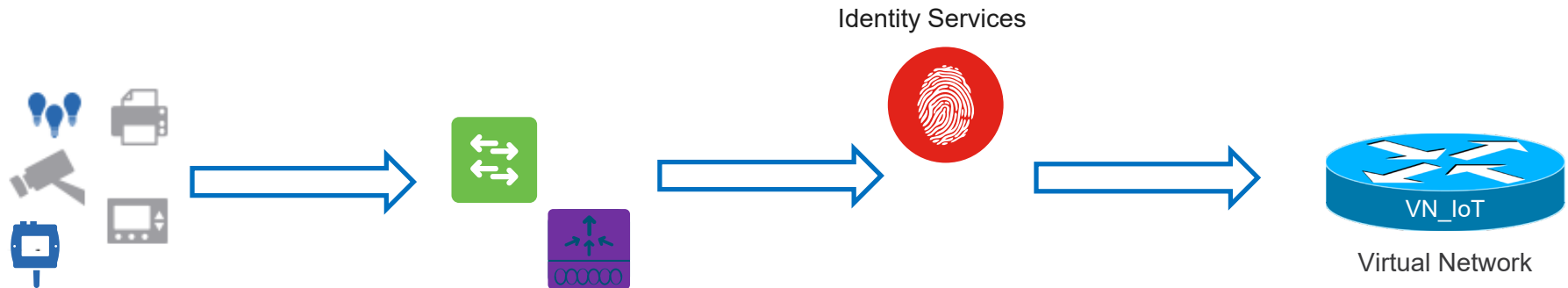
# ISE Authorization for VN assignment

Authorization Result = Virtual Network + Security Group

Work Centers > TrustSec > Authorization Policy

Status	Rule Name	Conditions	Results	Profiles	Security Groups	Hits	Actions
	IoT_Devices_VN_Only	IdentityGroup-Name EQUALS Endpoint Identity Groups:Profiled:IoT_DEVICES	<input type="text" value="*VN_IoT"/>		Select from list	0	

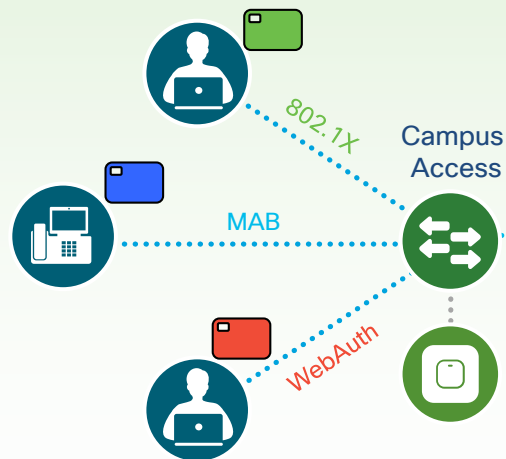
**VN+SG**



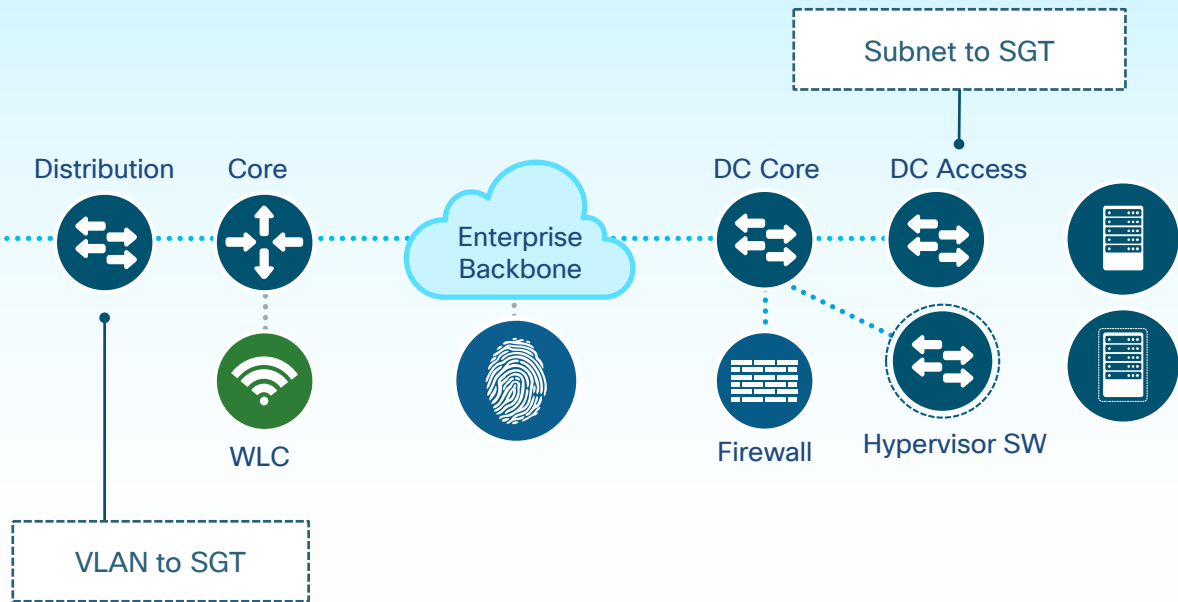
# Group Assignment

Two ways to assign SGT

## Dynamic Classification



## Static Classification





# Defining Security Groups

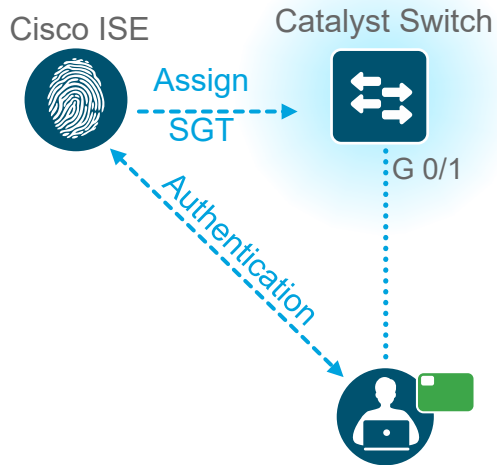
Define SGs under 'Components' section in **TrustSec Work Center** (from ISE 2.0)

The screenshot shows the Cisco Identity Services Engine (ISE) TrustSec Work Center interface. The navigation bar includes 'Home', 'Context Visibility', 'Operations', 'Policy', and 'Administration'. The 'Administration' menu is expanded to show 'Work Centers', 'Network Access', 'Guest Access', and 'TrustSec'. Under 'TrustSec', there are sub-menus for 'BYOD', 'Profiler', and 'Posture'. The 'Components' section is selected, showing 'TrustSec Policy', 'Authentication Policy', 'Authorization Policy', 'SXP', 'Troubleshoot', 'Reports', and 'Settings'. The left sidebar lists 'Security Groups', 'IP SGT Static Mapping', 'Security Group ACLs', 'Network Devices', and 'Trustsec AAA Servers'. The 'Security Groups' section is expanded, and 'Device SGT' is highlighted with a blue arrow. The main content area displays the 'Security Groups' table, which is outlined in red. The table has columns for 'Icon', 'Name', 'SGT (Dec / Hex)', and 'Description'. The 'TrustSec\_Devices' group is selected with a checkmark in the first column.

**Security Groups**  
For Policy Export go to [Administration > System > Backup & Restore > Policy Export Page](#)

	Icon	Name	SGT (Dec / Hex)	Description
<input type="checkbox"/>	?	Unknown	0/0000	Unknown Security Group
<input checked="" type="checkbox"/>	🌐	TrustSec_Devices	2/0002	TrustSec Devices Security Group
<input type="checkbox"/>	🌐	Test_Servers	13/000D	Test Servers Security Group
<input type="checkbox"/>	🌐	Quarantined_Systems	255/00FF	Quarantine Security Group
<input type="checkbox"/>	🌐	Production_Users	7/0007	Production User Security Group
<input type="checkbox"/>	🌐	Production_Servers	11/000B	Production Servers Security Group
<input type="checkbox"/>	🌐	Point of Sale Systems	10/000A	Point of Sale Security Group

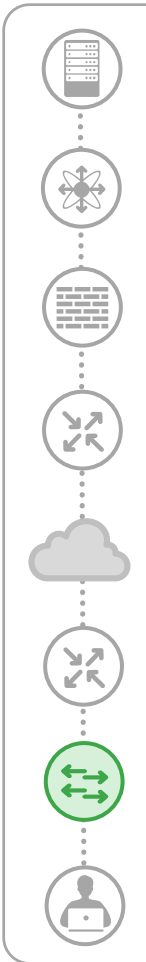
# SGT for wired 802.1X session



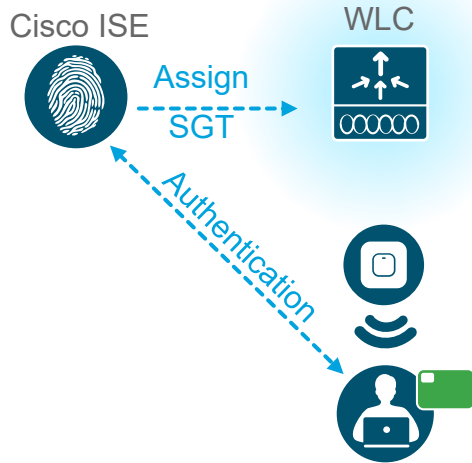
```
Switch# show authentication sessions int Gi 0/1 details
      Interface: GigabitEthernet1/0/23
          IIF-ID: 0x107AB4000000076
      MAC Address: 0005.0005.0005
      IPv6 Address: 2001:DB8:100:0:3809:A879:5197:16DB
      IPv4 Address: 172.20.100.2
      User-Name: bob@trustsec.lab
      Status: Authorized
      Domain: DATA
      Oper host mode: multi-auth
      Oper control dir: both
      Session timeout: N/A
      Common Session ID: 0A01010100000FC50BEC5800
      Acct Session ID: 0x00000FBE
      Handle: 0xD4000009
      Current Policy: POLICY_Gi1/0/23

Server Policies:
      SGT Value: 10

Method status list:
      Method          State
      mab             Authc Success
```



# SGTs for Wireless sessions



Works on AirOS and IOS-XE  
Wireless controllers.

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MONITOR WLANs CONTROLLER WIRELESS

Monitor

Summary

- Access Points
- Cisco CleanAir
- Statistics
- CDP
- Rogues
- Clients**
- Multicast

Client Properties

MAC Address	00:50:56:01:00:03
IPv4 Address	10.0.202.03
IPv6 Address	

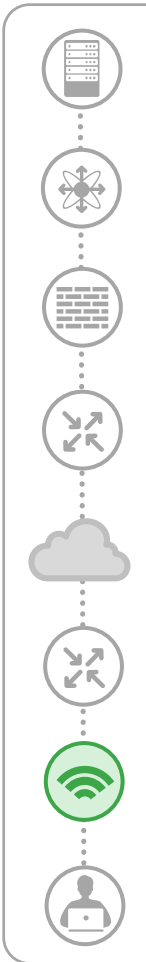
Security Information

Security Policy Completed	Yes
Policy Type	RSN (WPA2)
Encryption Cipher	CCMP (AES)
EAP Type	PEAP
SNMP NAC State	Access
Radius NAC State	RUN
<b>CTS Security Group Tag</b>	<b>5</b>

Client Type Regular

User Name **employee2**

Port Number 1



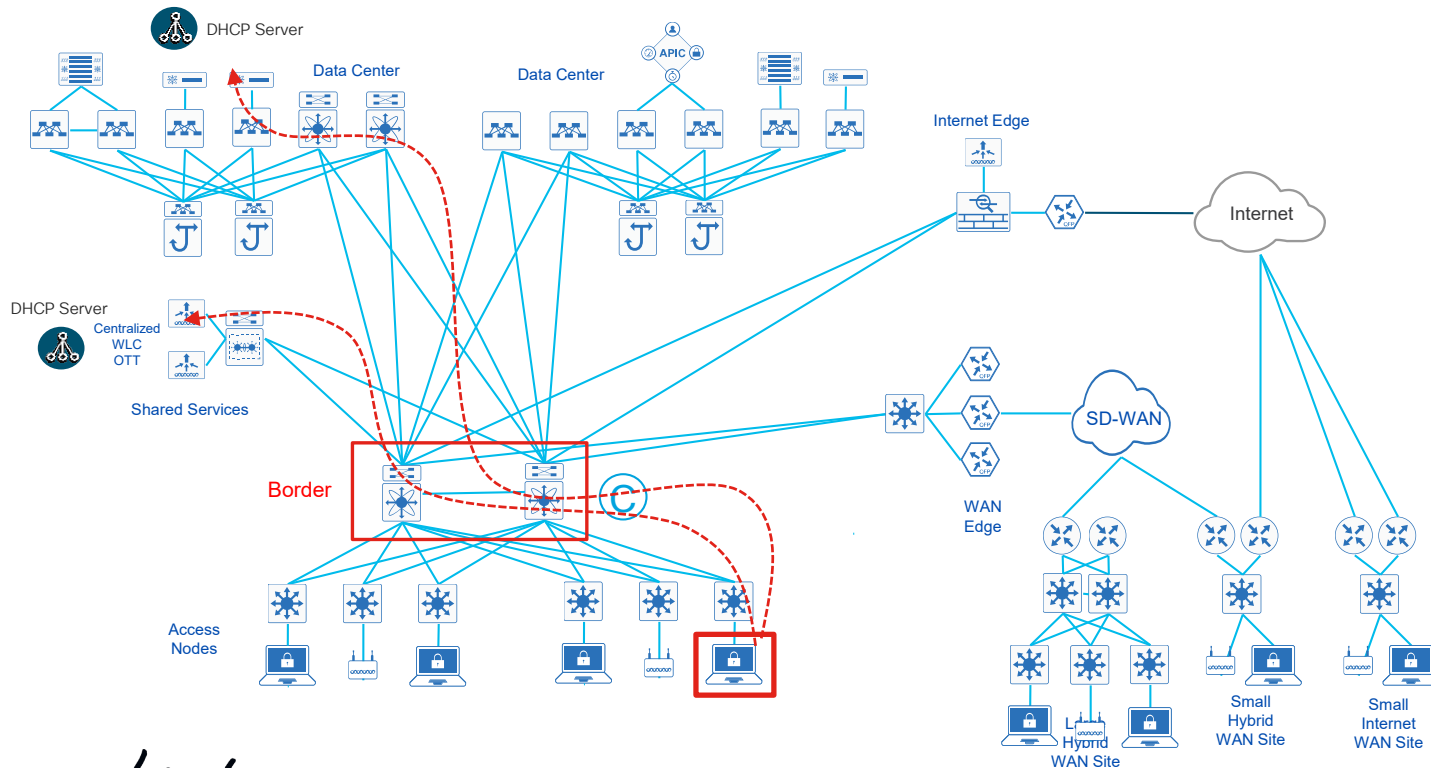
Gaining access to network fabric

# Host Onboarding

1. Endpoint Classification
2. DHCP in Fabric

# SD-Access Fabric Architecture

## DHCP reachability



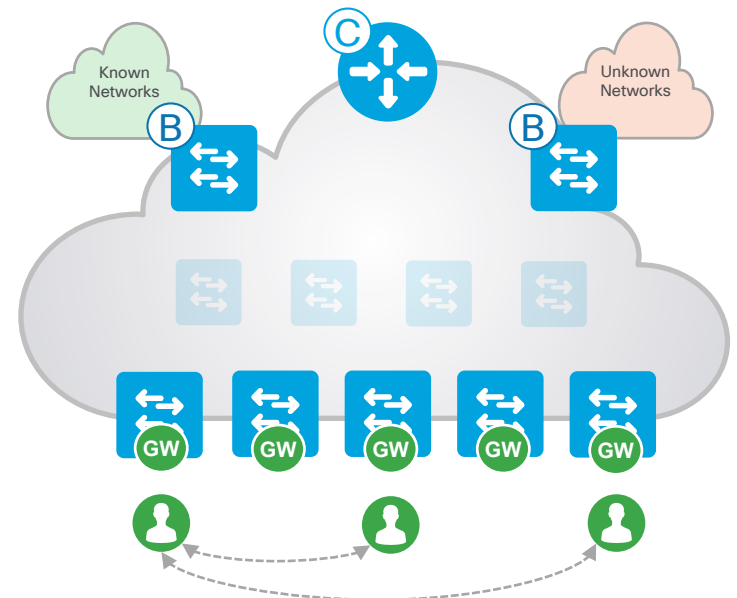
# SD-Access Fabric Architecture

## DHCP in an anycast Gateway environment



**Anycast GW** provides a single L3 Default Gateway for IP capable endpoints

- The same Switch Virtual Interface (SVI) is present on EVERY Edge with the SAME Virtual IP and MAC
- When a Host moves from Edge 1 to Edge 2, it does not need to change its Default Gateway 😊
- The SVI is also configured with an IP helper address for DHCP.

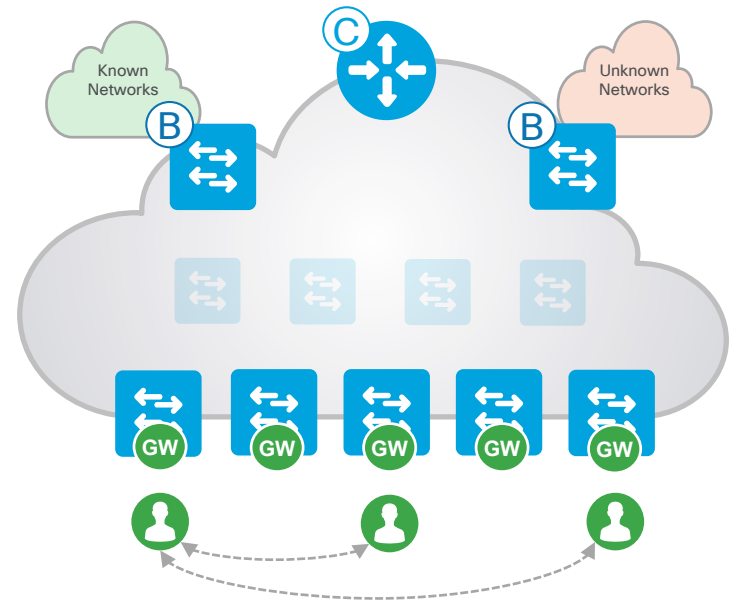


# SD-Access Fabric Architecture

## DHCP in an anycast Gateway environment

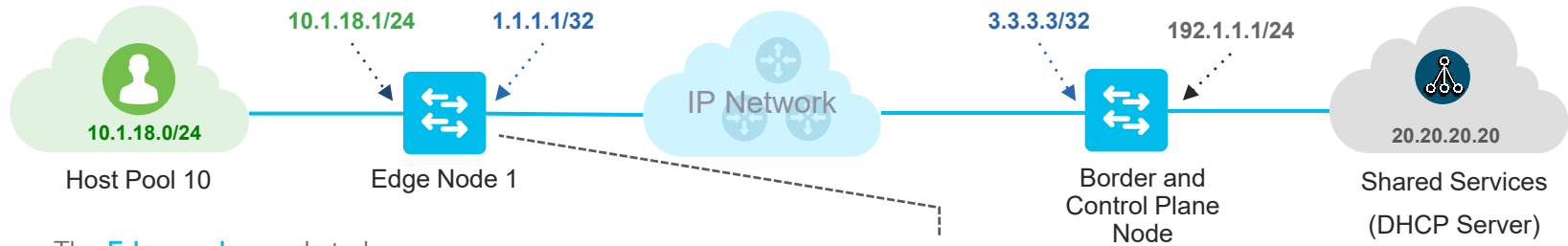
**DHCP reply** needs to come to the right edge node switch.

- But “we do not know on which edge node a host is located” as we don’t have an IP address for it yet.
- Once an IP address is assigned to the host, the control plane node learns where the host is located.
- The Control-Plane maintains the Host to Edge relationship (Fabric Dynamic EID mapping)



# SD-Access Fabric Architecture

## DHCP in Fabric – Enable LISP lookups for DHCP requests



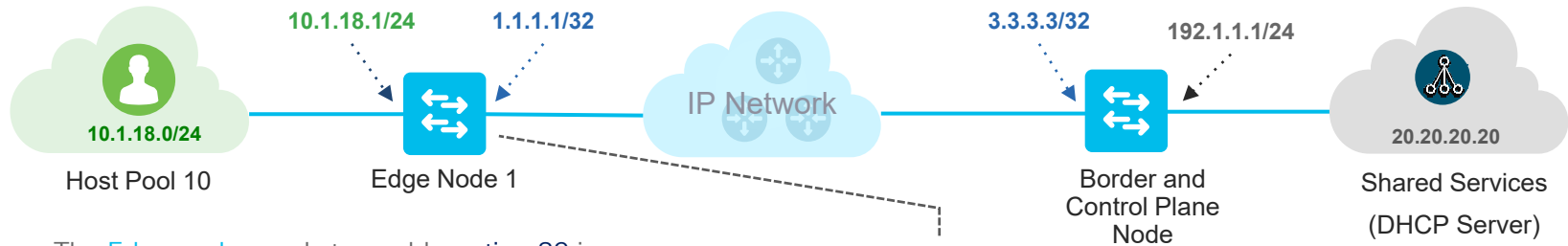
- The **Edge node** needs to be configured as a “**proxy-itr**” to avoid source EID validation. The source of the DHCP request doesn’t have an IP address yet.
- We also need a **0/0 map-cache** that triggers a LISP lookup for the DHCP helper address so that the DHCP request is sent in the overlay.

```
router lisp
instance-id 4098
dynamic-eid user
database-mapping 10.1.18.0/24 locator-set edge1
exit-dynamic-eid
!
service ipv4
eid-table vrf User
map-cache 0.0.0.0/0 map-request
itr map-resolver 3.3.3.3
proxy-itr 1.1.1.1
etr map-server 3.3.3.3 key uci
etr
use-petr 3.3.3.3
exit-service-ipv4
!
exit-instance-id
!
exit-router-lisp
```



# SD-Access Fabric Architecture

## DHCP in Fabric – Option 82 plus Snooping



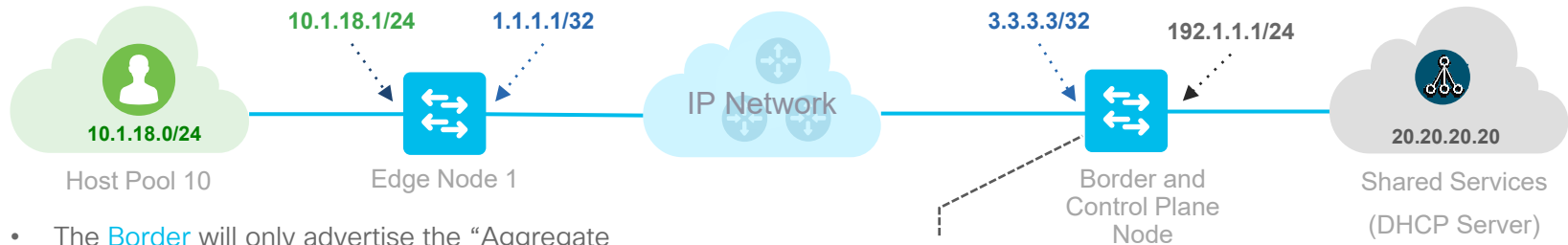
- The **Edge node** needs to enable option 82 in the DHCP request.
  - Option 82 will carry the VNID and RLOC.
- DHCP snooping needs to be enabled on all the VLANs in fabric.

```
ip dhcp relay information option
ip dhcp snooping vlan 1022
ip dhcp snooping

interface Vlan1022
ip vrf forwarding User
ip address 10.1.18.1 255.255.255.0
ip helper-address 20.20.20.20
lisp mobility user
```

# SD-Access Fabric Architecture

## DHCP in Fabric – “Pre-advertise” IP pools outside the fabric



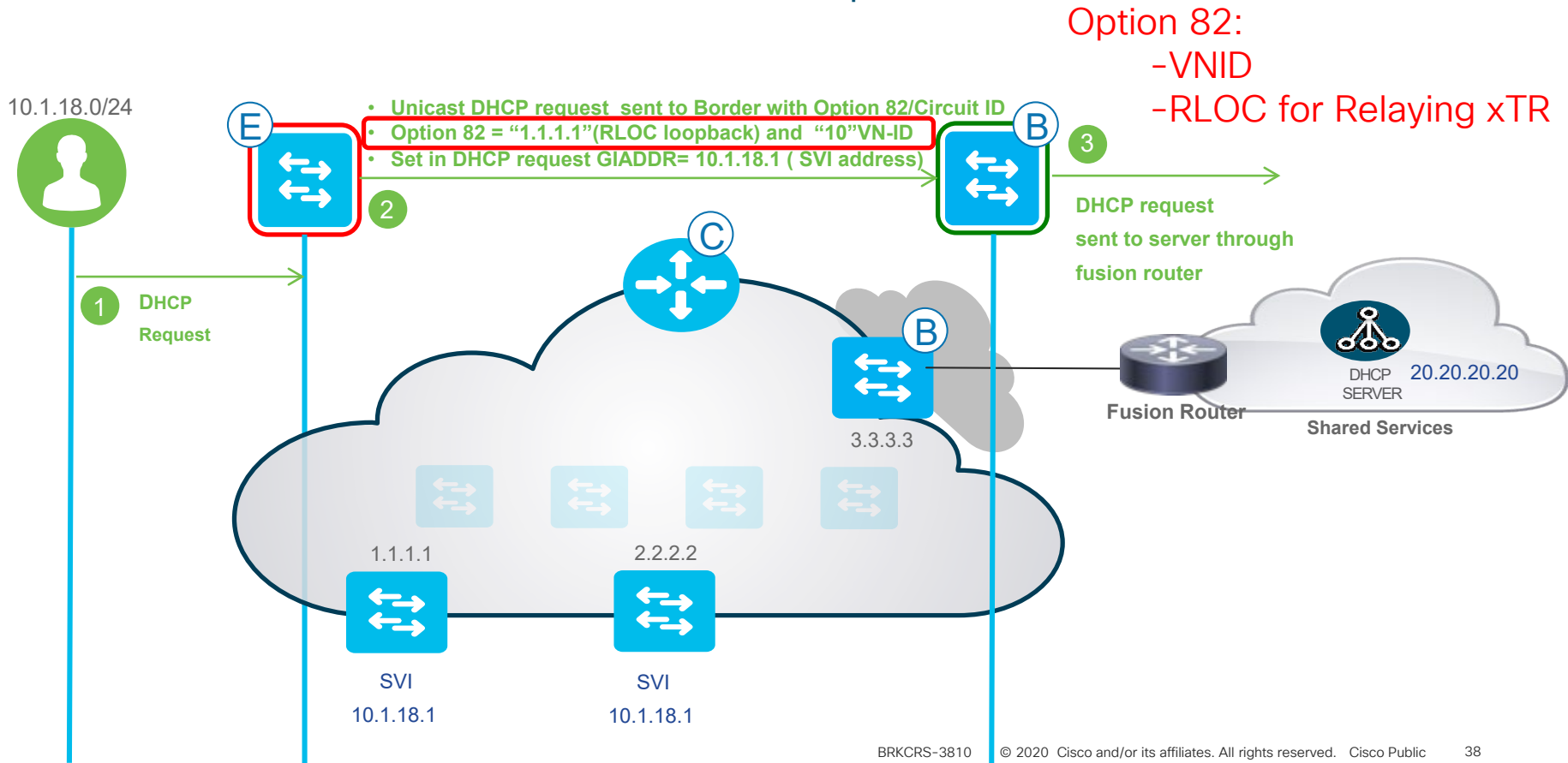
- The **Border** will only advertise the “Aggregate route(10.1.18.0/24) when there is a more specific prefix(host route)for that subnet in RIB”
- For the DHCP server to send the DHCP reply back we need a route to 10.1.18.0/24 at the DCHP server side.
- This will not happen until we have a HOST in the subnet 10.1.18.0/24 registered with LISP HTDB. This needs DHCP to happen first.

```
router bgp 65002
  bgp log-neighbor-changes
  !
  address-family ipv4 vrf USER
    aggregate-address 10.1.18.0 255.255.255.0 summary-only
    redistribute lisp metric 10
    neighbor x.x.x.x remote-as xxxxx
  exit-address-family

  interface Loopback1022
  vrf forwarding User
  ip address 10.1.18.1 255.255.255.255
```

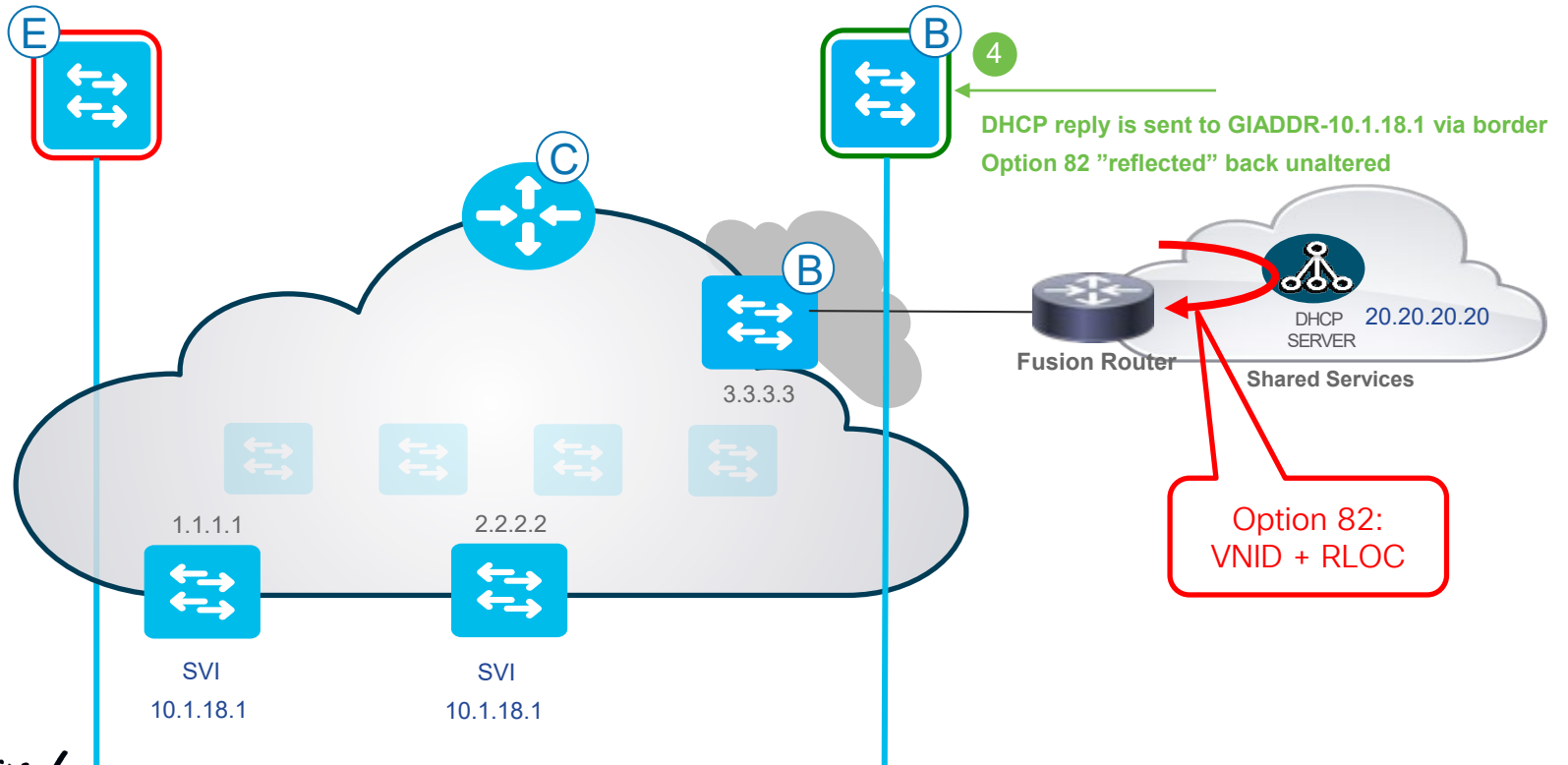
# SD-Access Fabric Architecture

## DHCP in Fabric – Include RLOC in DHCP request



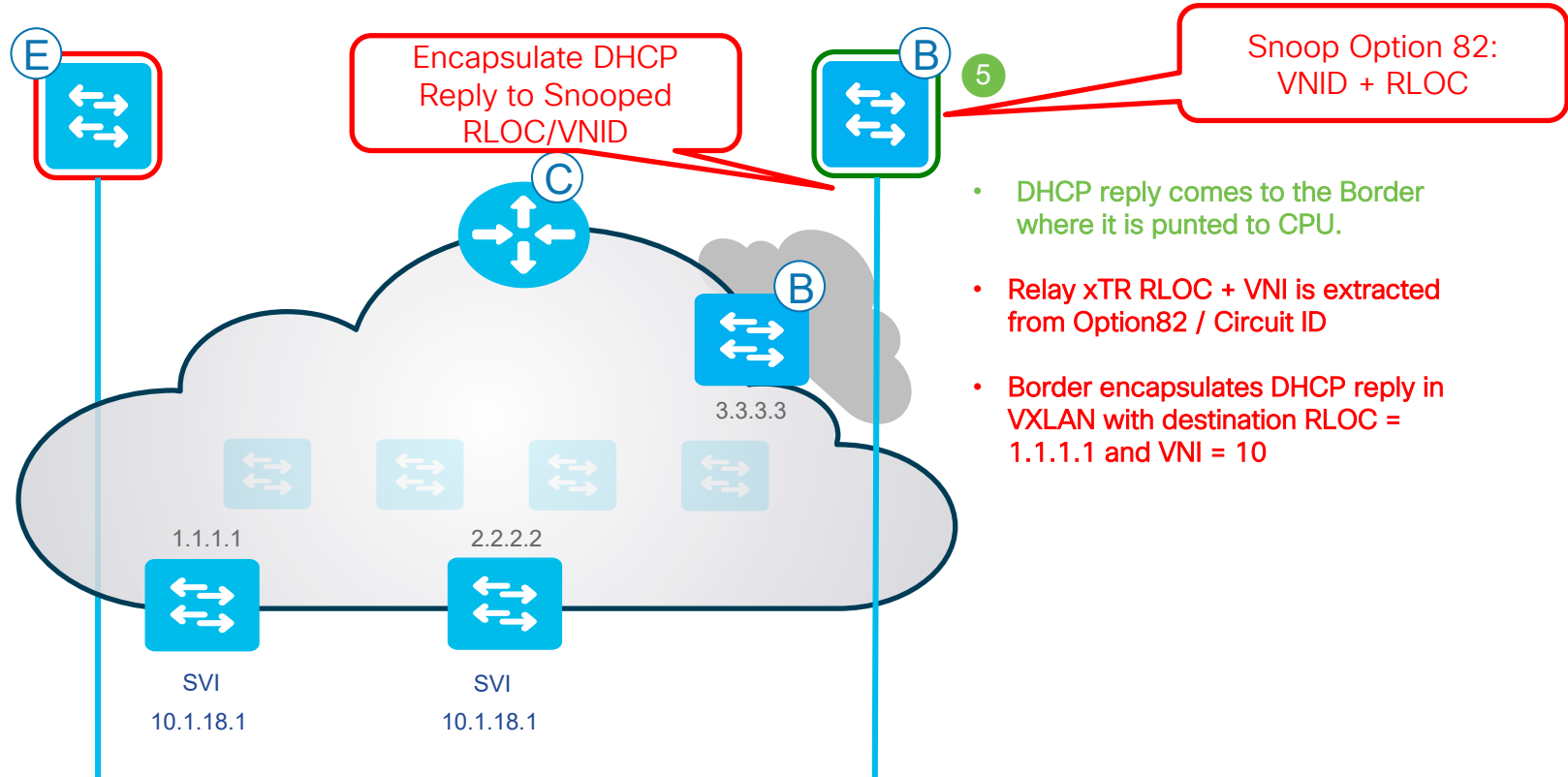
# SD-Access Fabric Architecture

## DHCP in Fabric – DHCP Reply



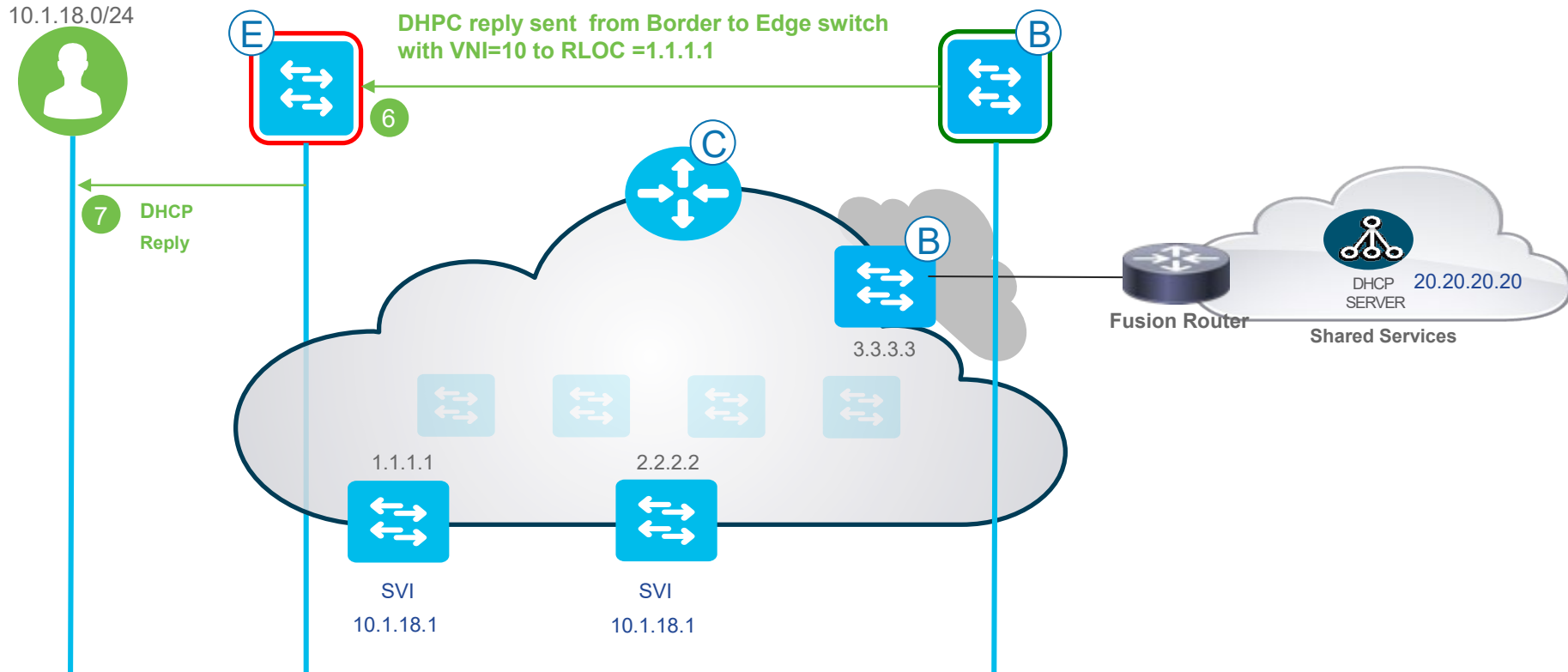
# SD-Access Fabric Architecture

## DHCP in Fabric – Snoop and encapsulate DHCP reply at the Border



# SD-Access Fabric Architecture

## DHCP in Fabric – Complete DHCP reply



# Unicast Packet Forwarding & Access Control

1. Control Plane Lookup
2. Access Control Policy

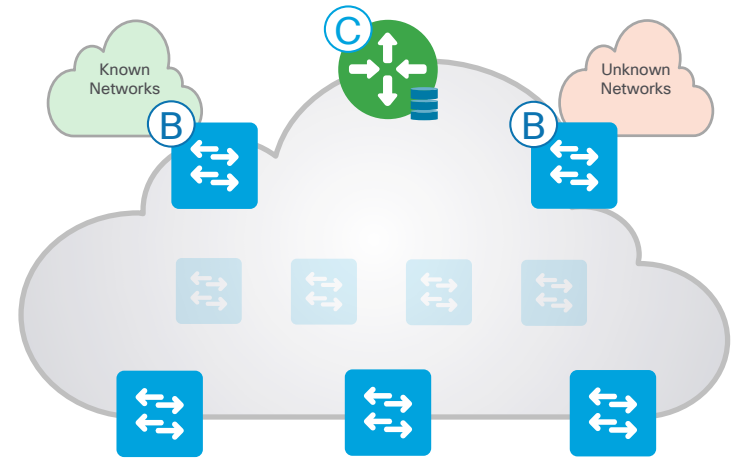
# Cisco SD-Access Architecture

## Control-Plane Nodes – A Closer Look



**Control-Plane Node** runs a Host Tracking Database to map location information

- A simple Host Database that maps Endpoint IDs to a current Location, along with other attributes
- Host Database supports multiple types of Endpoint ID lookup types (IPv4, IPv6 or MAC)
- Receives Endpoint ID map registrations from Edge and/or Border Nodes for “known” IP prefixes
- Resolves lookup requests from Edge and/or Border Nodes, to locate destination Endpoint IDs

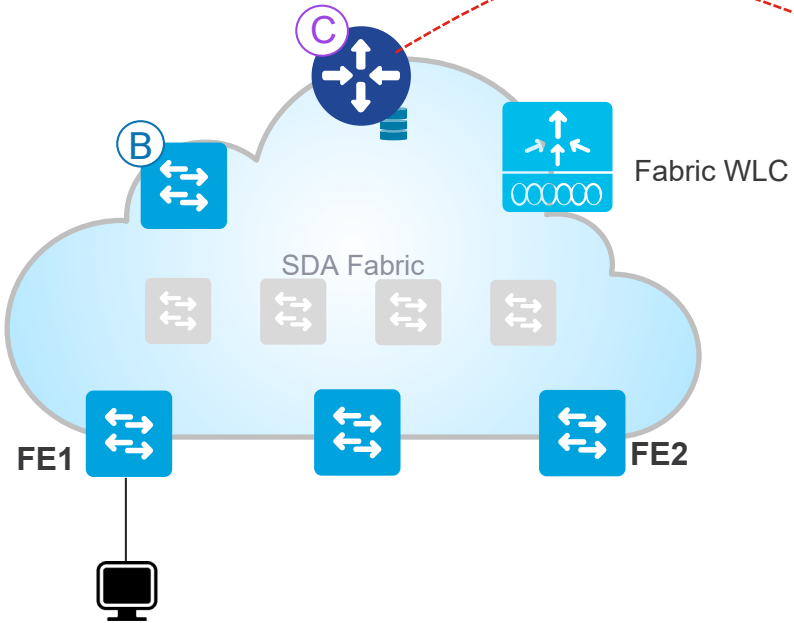




# Endpoint Registration in Fabric

# Cisco SD-Access Fabric Architecture

## Host Registration



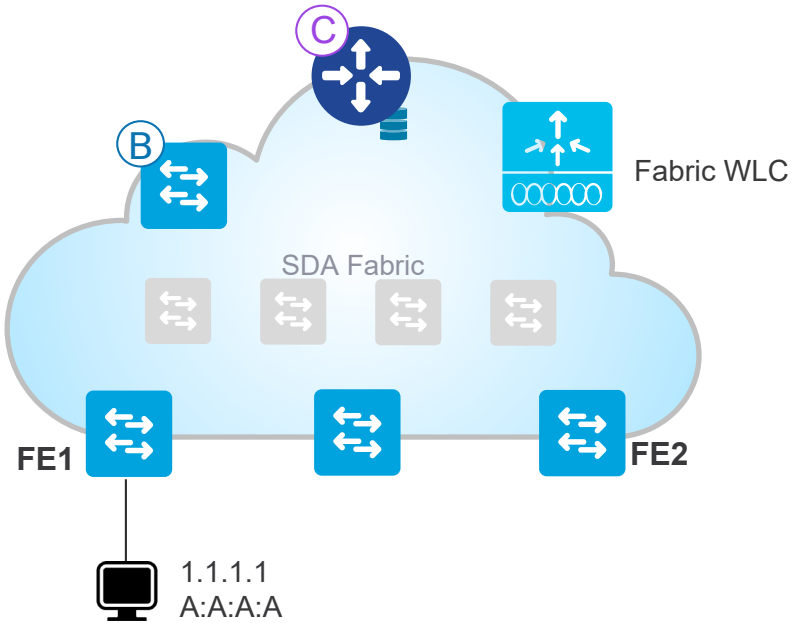
Control Plane state:

IP to RLOC Table	MAC to RLOC Table	Address Resolution Table

- The Control Plane node has three related tables:
  - **IP to RLOC** stores IP address of a Host and its corresponding location
  - **MAC to RLOC** stores MAC address of a Host and its corresponding location
  - **Address Resolution** Data from the above two tables are collated for IP to MAC bindings (ARP Table)

# Cisco SD-Access Fabric Architecture

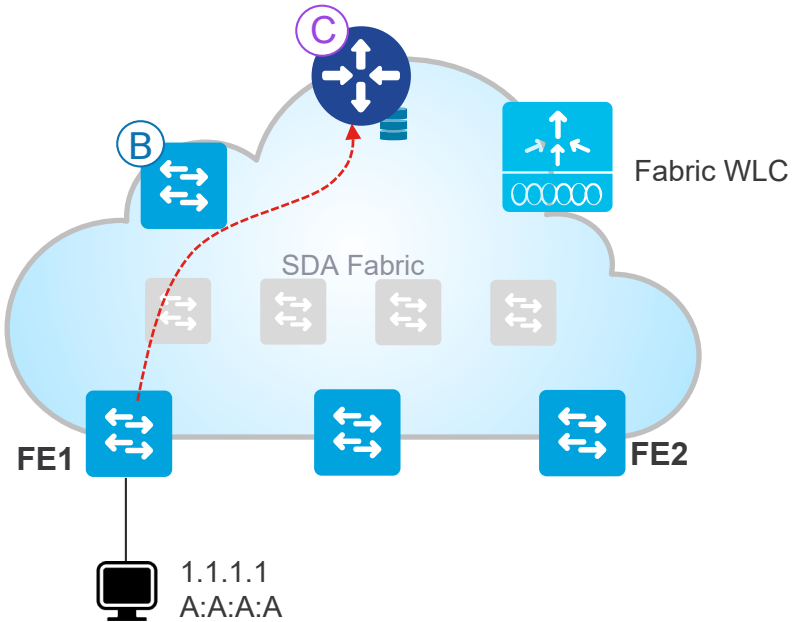
## Host Registration



- 1 Wired host attaches to the fabric network on an edge node

# Cisco SD-Access Fabric Architecture

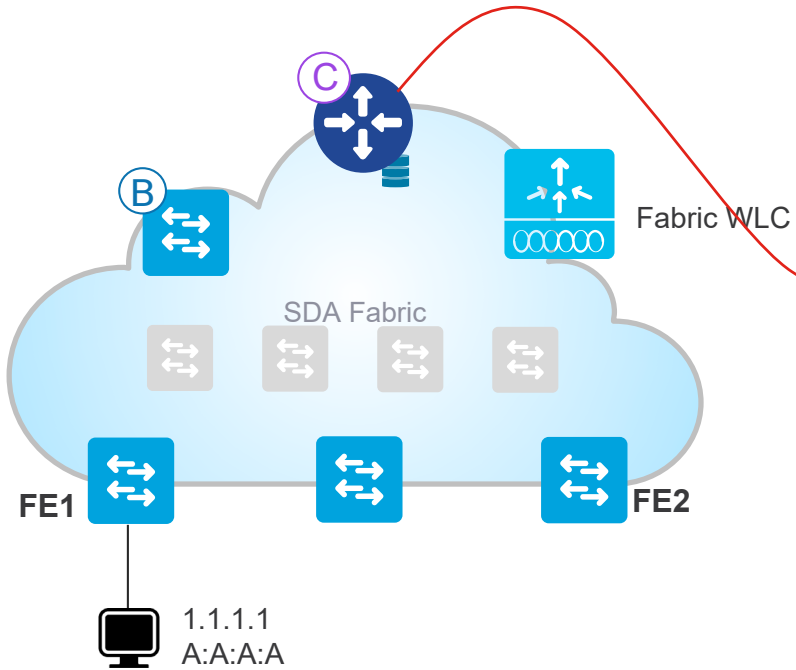
## Host Registration



- 2 After the host gets an IP address, the fabric edge sends a map-register to the control plane node with the IP and mac address of the host

# Cisco SD-Access Fabric Architecture

## Host Registration



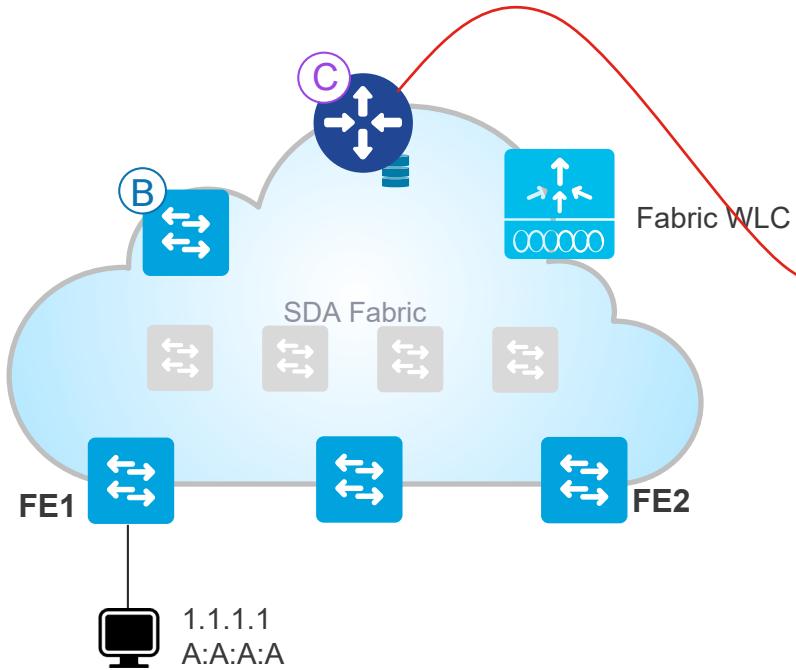
- 3 The control plane node upon receiving the map-register populates the database tables for the host

Control Plane state:

IP to RLOC Table	MAC to RLOC Table	Address Resolution Table
1.1.1.1 → FE1	AA:AA:AA:AA → FE1	

# Cisco SD-Access Fabric Architecture

## Host Registration



- 4 The Control plane then takes the information from the IP and MAC table and populates the ARP table

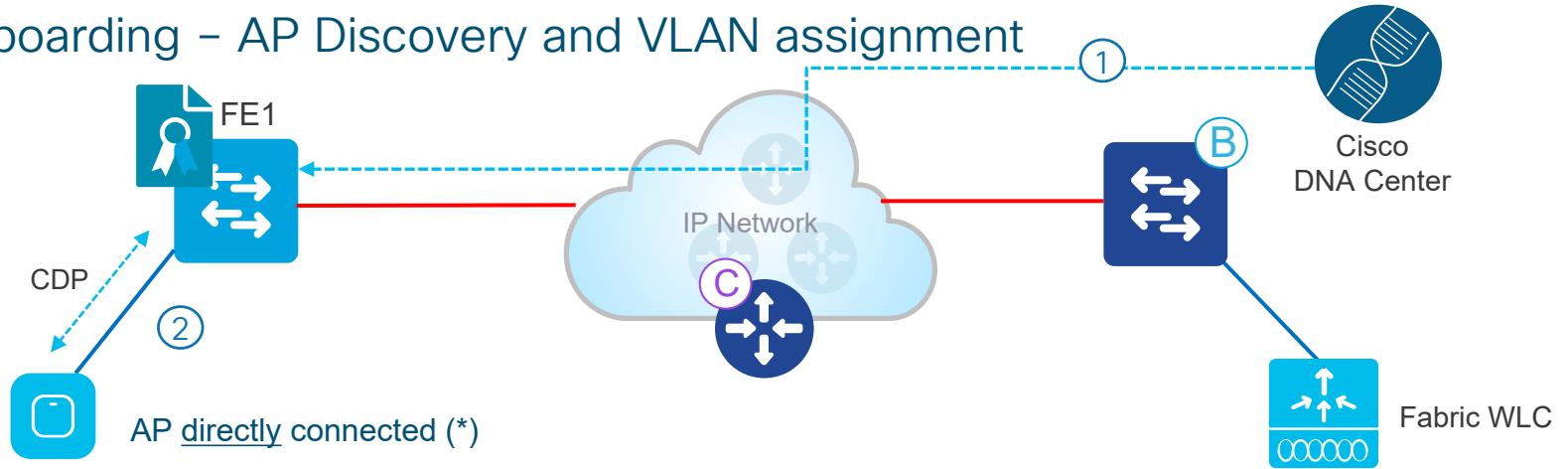
Control Plane state:

IP to RLOC Table	MAC to RLOC Table	Address Resolution Table
1.1.1.1 → FE1	AA:AA:AA:AA → FE1	1.1.1.1 = A:A:A:A

# Registering an Access Point in Fabric

# SD-Access Wireless Workflow

## AP onboarding – AP Discovery and VLAN assignment



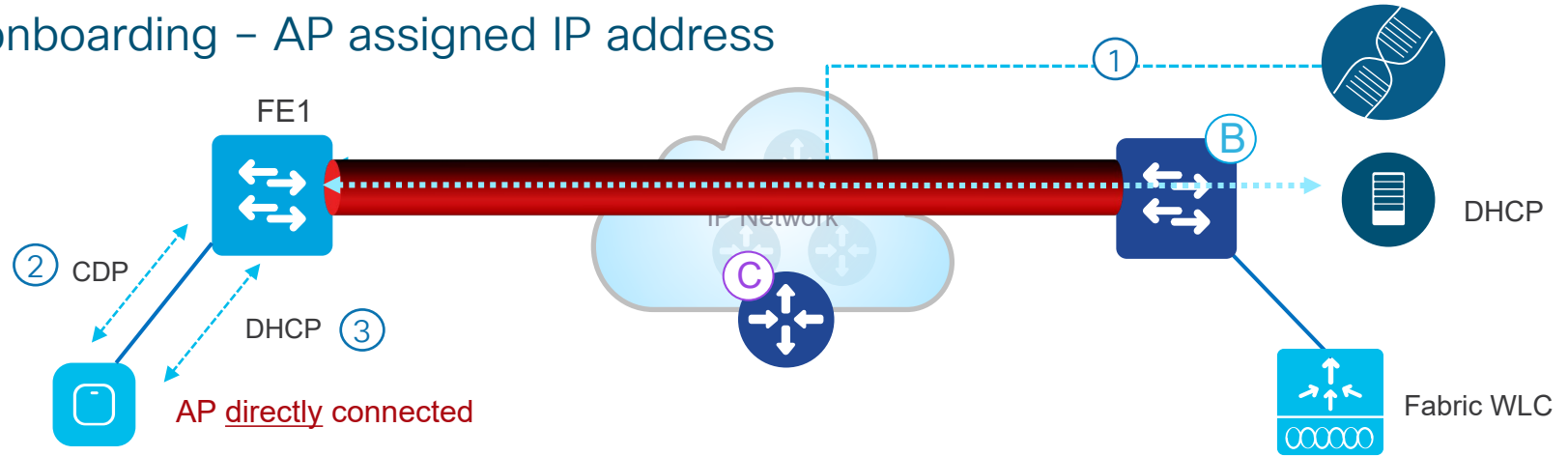
- ① Admin configures AP pool in Cisco DNA Center in INFRA\_VN. Cisco DNA Center pre-provision a configuration macro on all the FEs
- ② AP is plugged in and powers up. FE discovers it's an AP via CDP and applies the macro to assign the switch-port the right VLAN

(\*) AP can be connected also through an “Extended node” switch



# SD-Access Wireless Workflow

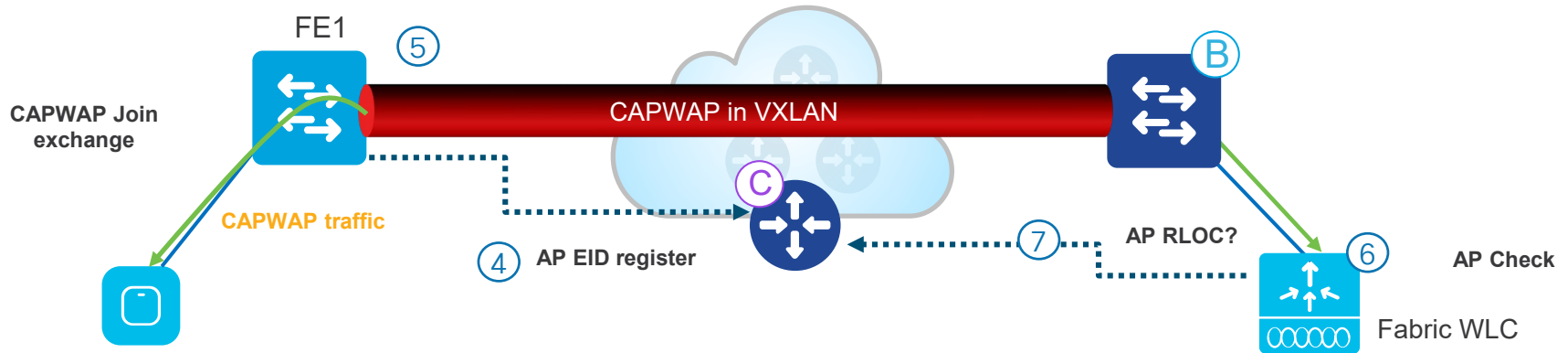
## AP onboarding – AP assigned IP address



- ① Admin configures AP pool in Cisco DNA Center in INFRA\_VN. Cisco DNA Center pre-provision a configuration macro on all the FEs
- ② AP is plugged in and powers up. FE discovers it's an AP via CDP and applies the macro to assign the switch port the the right VLAN
- ③ AP gets an IP address via DHCP in the overlay

# SD-Access Wireless Workflow

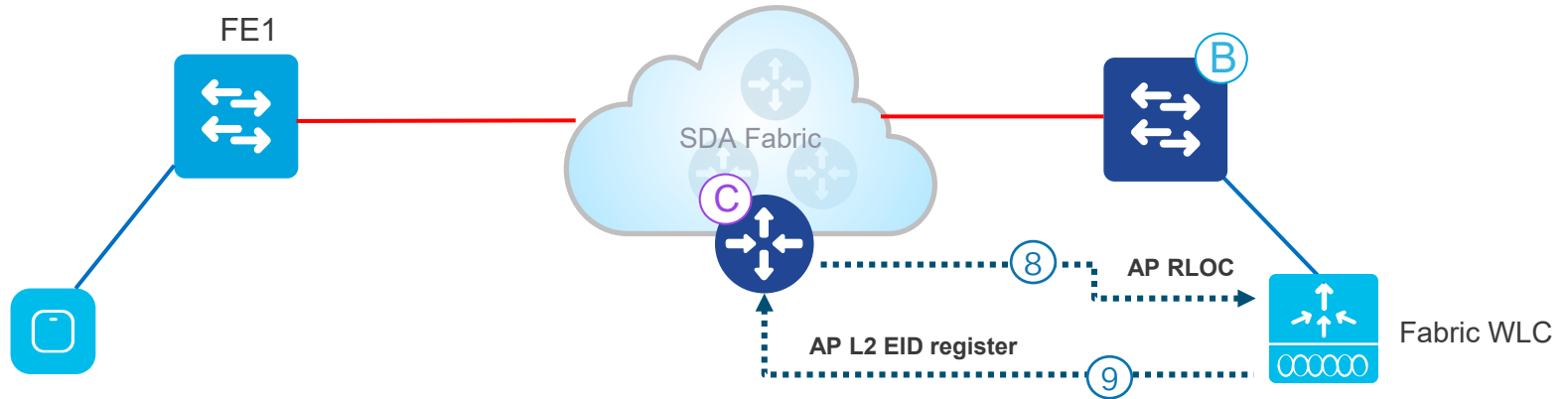
## AP onboarding – Register AP with Fabric and Connect to WLC



- ④ Fabric Edge registers AP's IP address and MAC (EID) and updates the Control Plane (CP)
- ⑤ AP learns WLC's IP and joins using traditional methods. Fabric AP joins in Local mode
- ⑥ WLC checks if AP is fabric-capable (11ax, Wave 2, Wave 1 APs)
- ⑦ If AP is supported, WLC queries the CP to know if AP is connected to Fabric

# SD-Access Wireless Workflow

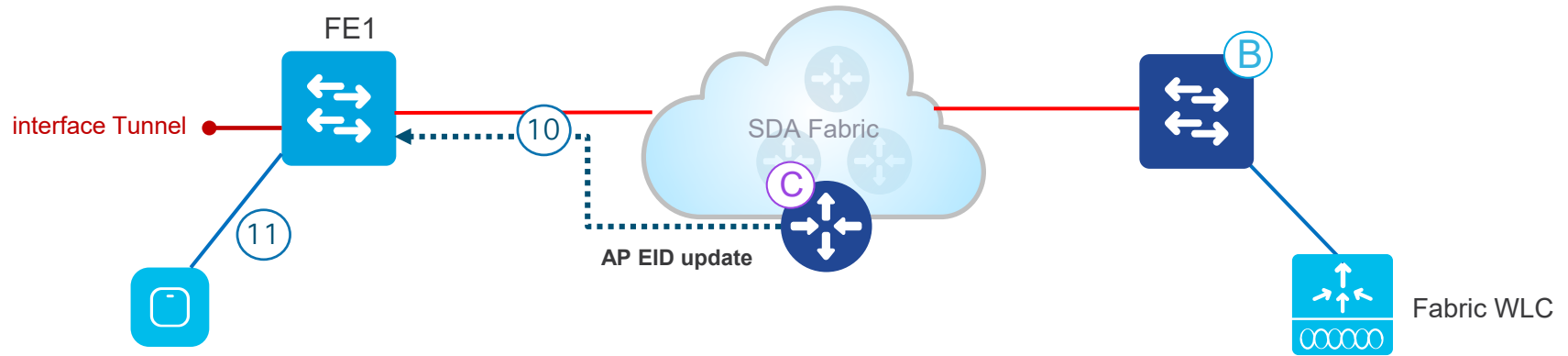
## AP onboarding – Switch AP to Fabric mode



- ⑧ Control Plane (CP) replies to WLC with RLOC. This means AP is attached to Fabric and will be shown as “Fabric enabled”
- ⑨ WLC does a L2 LISP registration for the AP in CP (a.k.a. AP “special” secure client registration). This is used to pass important metadata information from WLC to the FE

# SD-Access Wireless Workflow

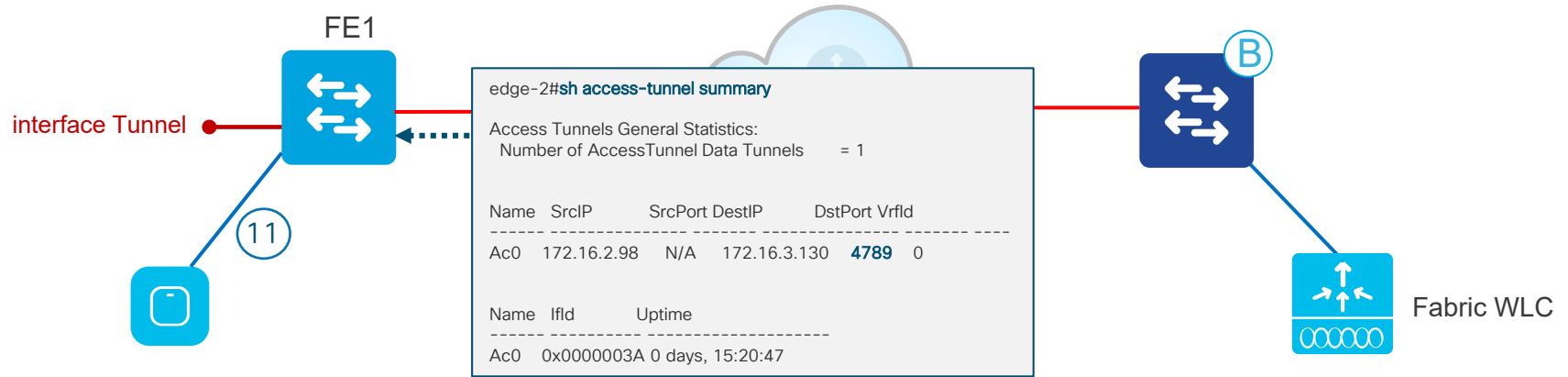
## AP onboarding – Program access ports for Fabric AP



- 10 In response to this proxy registration, Control Plane (CP) notifies Fabric Edge and pass the metadata received from WLC (flag that says it's an AP and the AP IP address)
- 11 Fabric Edge processes the information, it learns it's an AP and creates a VXLAN tunnel interface to the specified IP (optimization: switch side is ready for clients to join)

# SD-Access Wireless Workflow

## AP onboarding - Program access ports for Fabric AP

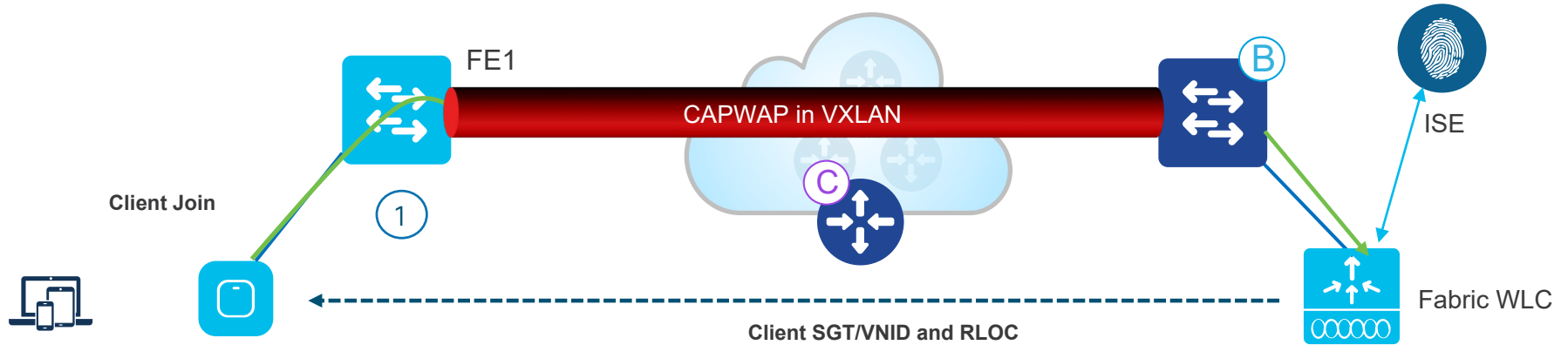


- 10 In response to this proxy registration, Control Plane (CP) notifies Fabric Edge and pass the metadata received from WLC (flag that says it's an AP and the AP IP address)
- 11 Fabric Edge processes the information, it learns it's an AP and creates a VXLAN tunnel interface to the specified IP (optimization: switch side is ready for clients to join)

# Registering a Wireless host in Fabric

# SD-Access Wireless Workflow

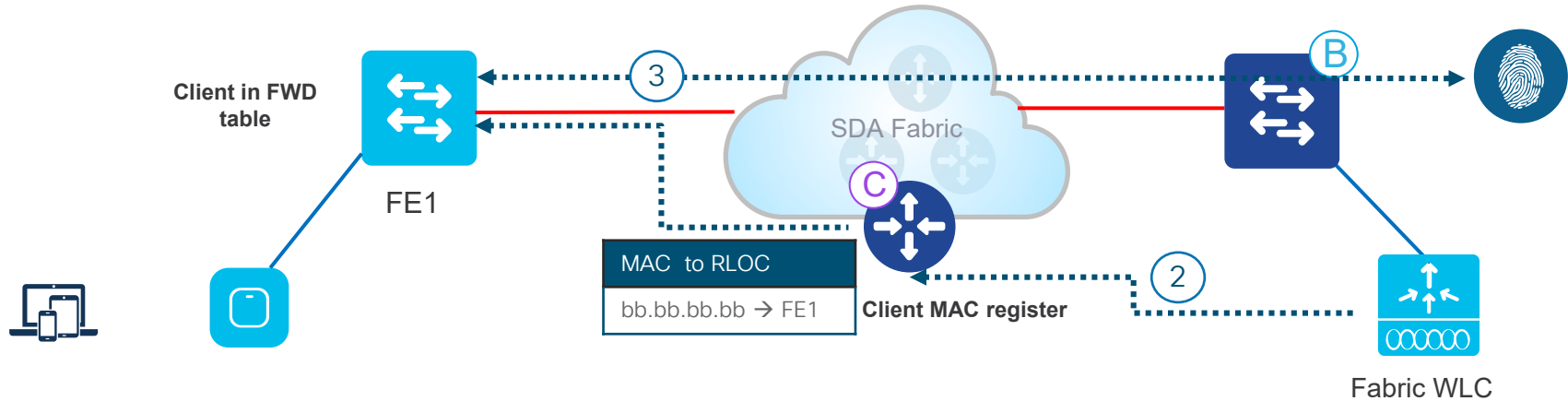
## Client Onboarding



- 1 Client authenticates to a Fabric enabled WLAN. WLC gets SGT from ISE, updates AP with client L2VNID and SGT

# SD-Access Wireless Workflow

## Client Onboarding

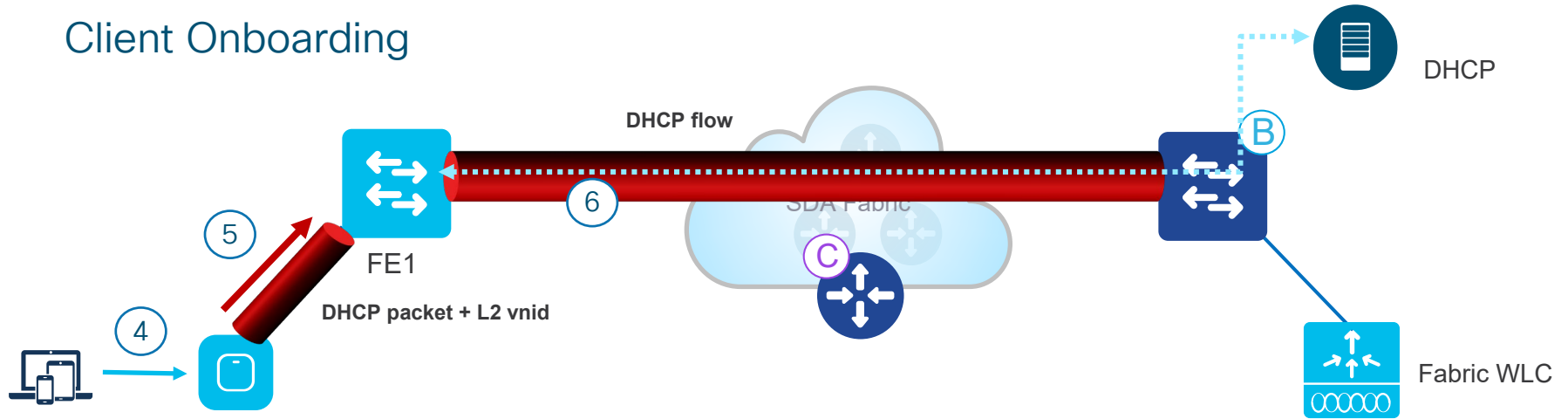


- 1 Client authenticates to a Fabric enabled WLAN. WLC gets SGT from ISE, updates AP with client L2VNID and SGT
- 2 WLC knows RLOC of AP from internal DB . WLC proxy registers Client L2 info in CP; this is LISP modified message to pass additional info, like the client SGT
- 3 FE gets notified by CP and knows it's a client; FE adds client MAC in L2 forwarding table and go and fetch the client policy from ISE based on the client SGT



# SD-Access Wireless Workflow

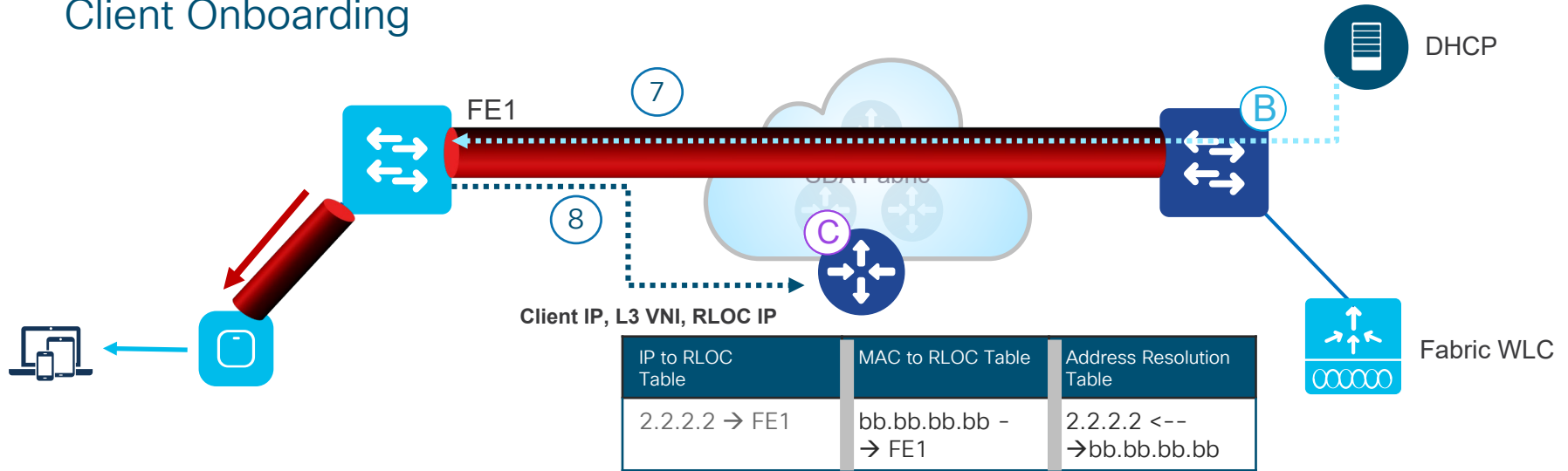
## Client Onboarding



- ④ Client initiates DHCP Request
- ⑤ AP encapsulates it in VXLAN with L2 VNI info (and SGT)
- ⑥ Fabric Edge maps L2 VNID to the VLAN interface and forwards the DHCP packet in the overlay (same as for a wired Fabric client)

# SD-Access Wireless Workflow

## Client Onboarding



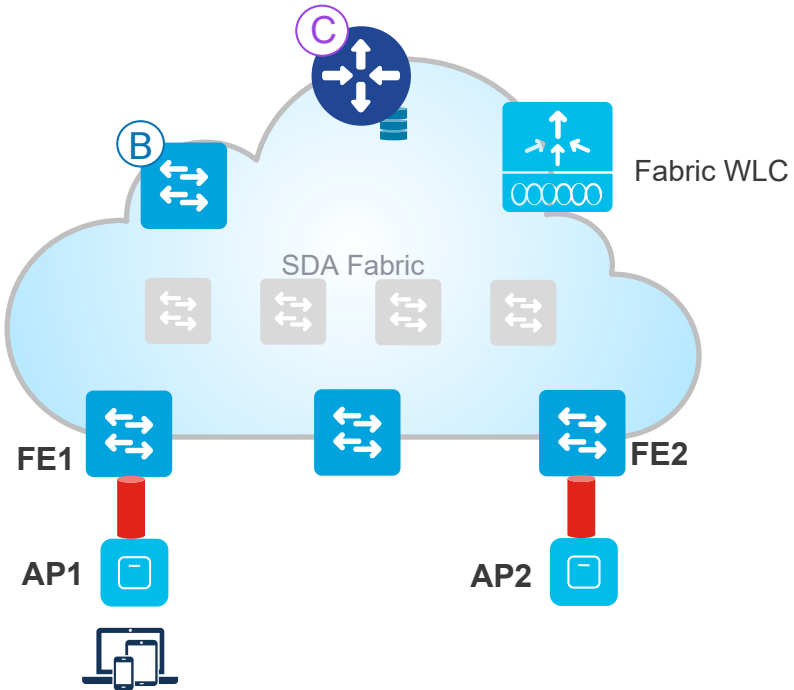
- 7 Client receives an IP address from DHCP
- 8 DHCP snooping triggers the client EID registration (MAC, IP address) by the Fabric Edge to the CP. *(If client has a static IP, then ARP or any other IP packet will trigger the registration)*

This completes client onboarding process

# Mobility Events in Fabric

# Cisco SD-Access Fabric Architecture

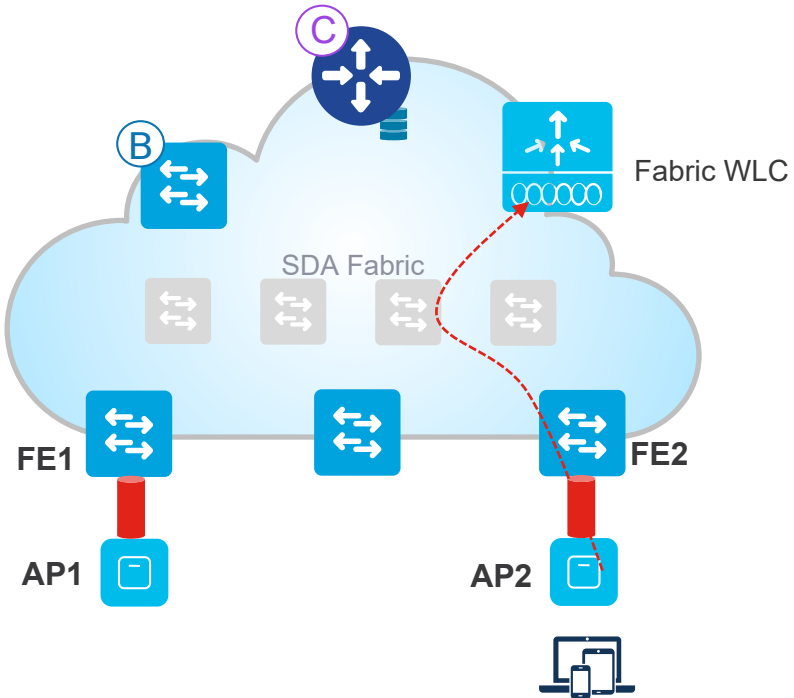
## Client Roam



- 1 Client roams from AP1 to AP2 (inter-switch roaming)

# Cisco SD-Access Fabric Architecture

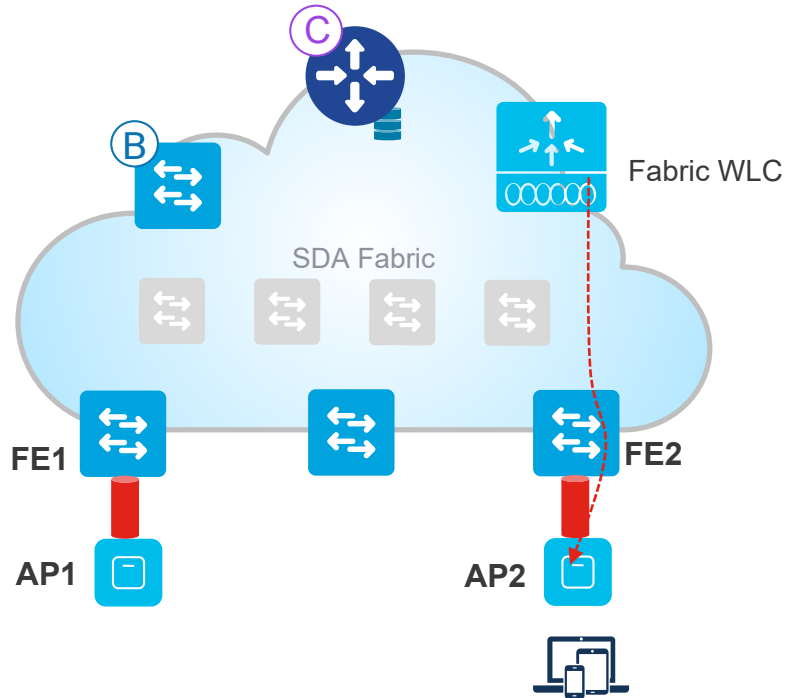
## Client Roam



- 2 AP2 detects the client move and registers its MAC address to the WLC as a mobility event

# Cisco SD-Access Fabric Architecture

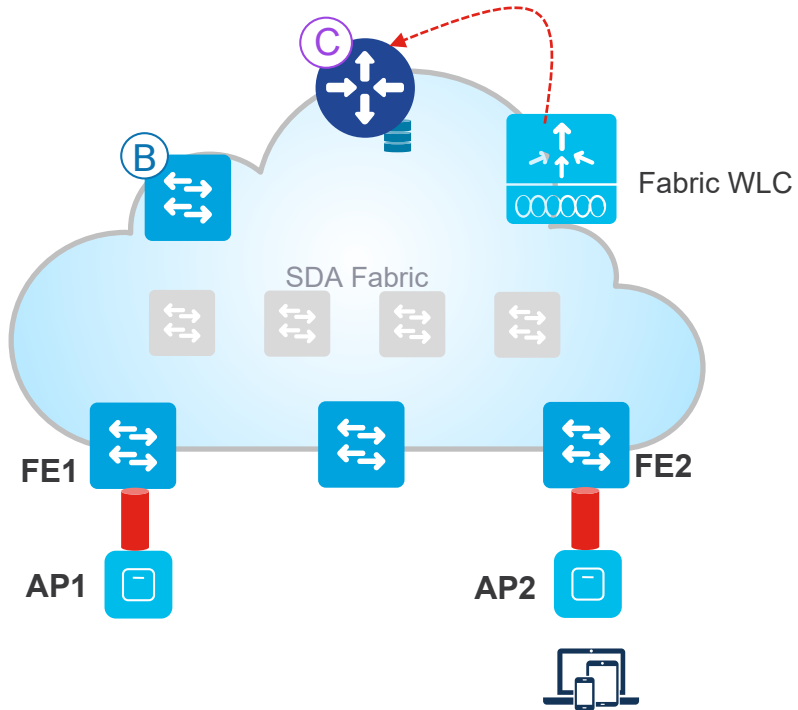
## Client Roam



- 3 WLC updates forwarding table on AP2 with client info (SGT, L2VNID, RLOC)

# Cisco SD-Access Fabric Architecture

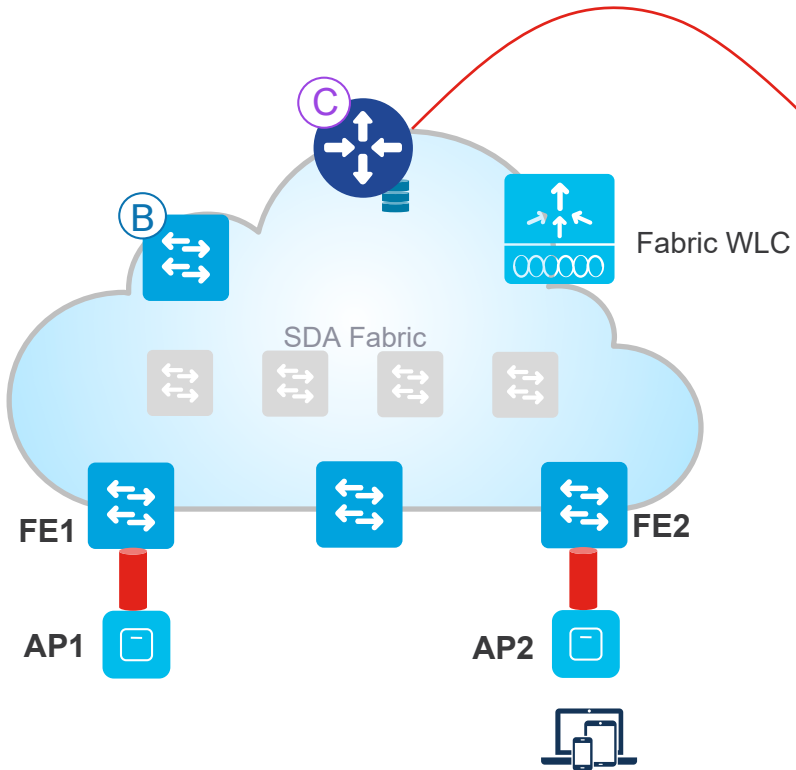
## Client Roam



- 4 WLC updates the L2 MAC entry in CP with new RLOC FE2.

# Cisco SD-Access Fabric Architecture

## Client Roam



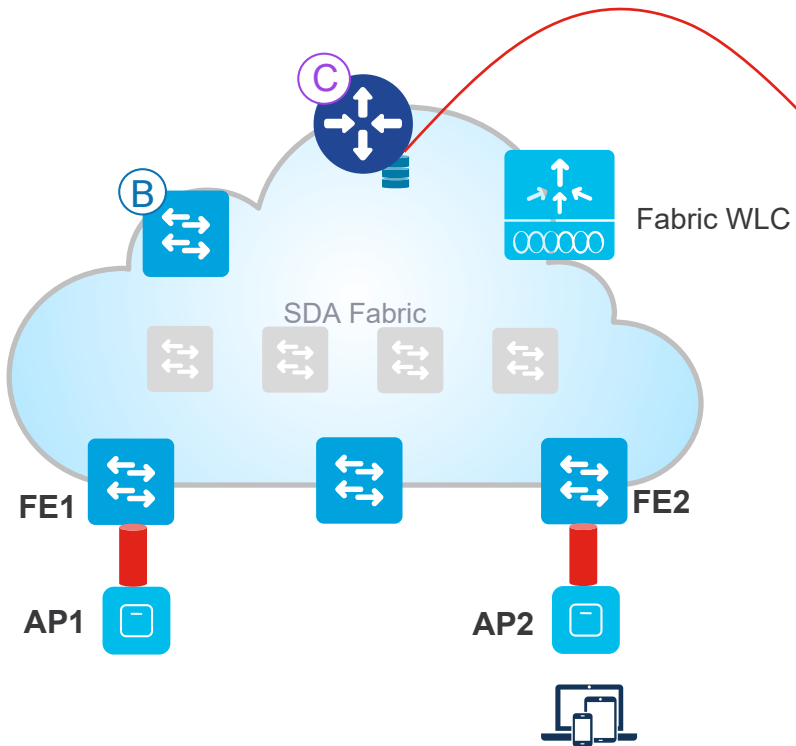
- 5 The control plane state will be mismatched as the IP address of host says it is located behind FE1 but MAC address says its behind FE2

IP to RLOC Table	MAC to RLOC Table	Address Resolution Table
2.2.2.2 → FE1	bb.bb.bb.bb - → FE2	2.2.2.2 <-- →bb.bb.bb.bb



# Cisco SD-Access Fabric Architecture

## Client Roam

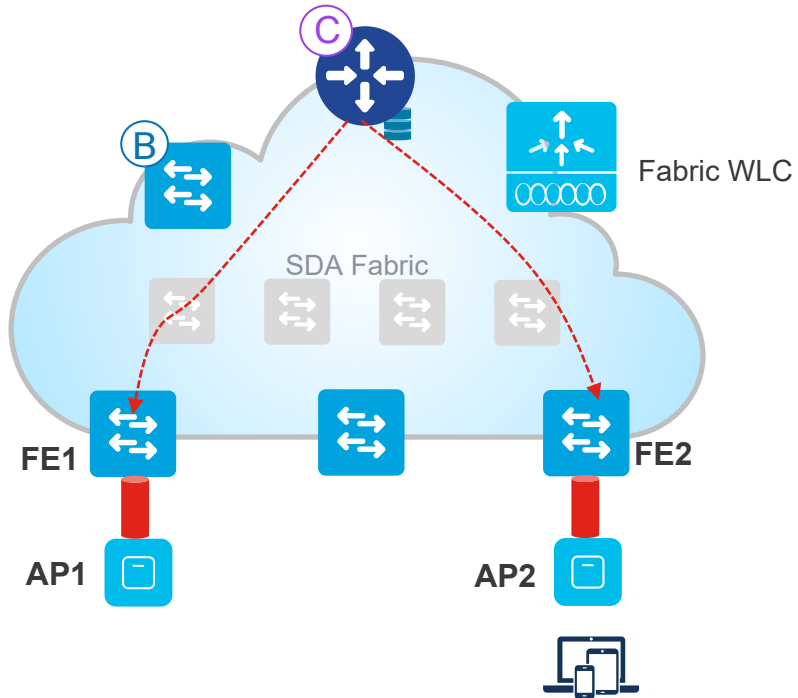


- 6 In this state the control plane node will trust the notification from the WLC and will correct its IP to RLOC table.

IP to RLOC Table	MAC to RLOC Table	Address Resolution Table
2.2.2.2 → FE2	bb.bb.bb.bb - → FE2	2.2.2.2 <-- → bb.bb.bb.bb

# Cisco SD-Access Fabric Architecture

## Client Roam



- 7 The control plane node after correcting its IP to RLOC table will send notifications to Edge 1 and Edge 2. CP then notifies
- Fabric Edge FE2 ("roam-to" switch) to add the client MAC to forwarding table pointing to VXLAN tunnel
  - Fabric Edge FE1 ("roam-from" switch) to do clean up for the wireless client

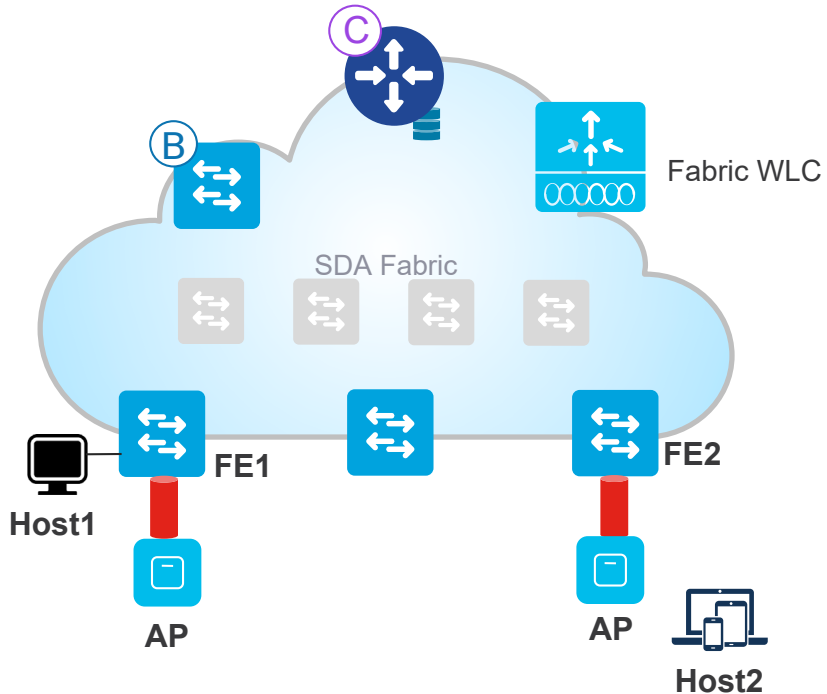
This ensures that when a host moves from one AP to other, the new edge node is waiting for it even before the move fully completes to ensure seamless roaming.

- 8 Roam is Layer 2 as FE2 has the same VLAN interface as FE1 (Anycast Gateway)

# Unicast Packet Forwarding in Fabric in the same subnet

# Cisco SD-Access Fabric Architecture

## Unicast Forwarding in the same subnet



### Assumptions:

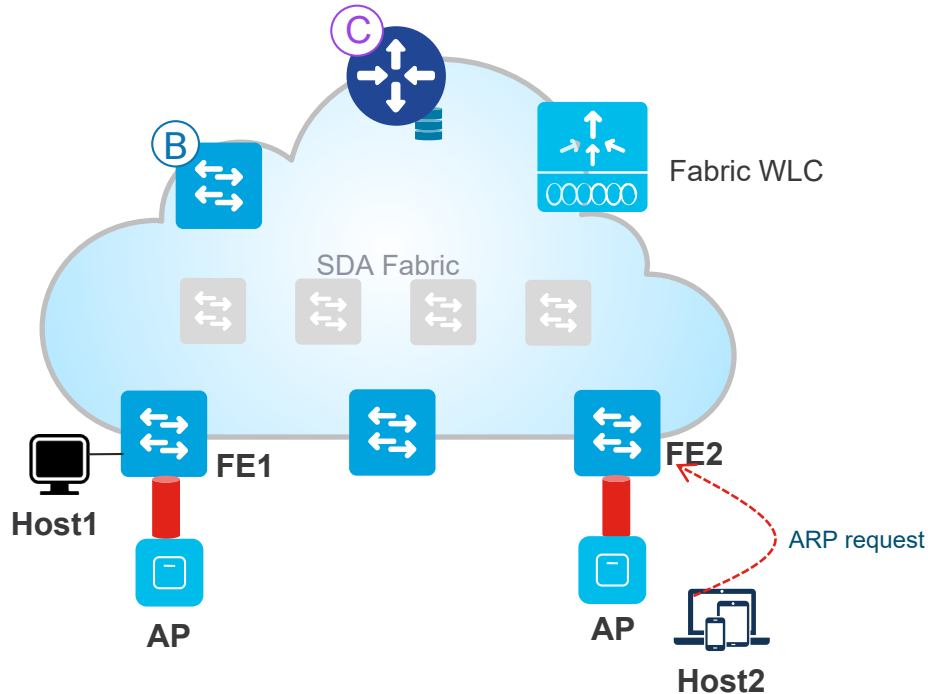
- Host H1 and Host H2 belong to the same subnet

### Control Plane state:

IP to RLOC Table	MAC to RLOC Table	Address Resolution Table
1.1.1.1 → FE1	aa.aa.aa.aa - → FE1	1.1.1.1 <-- →aa.aa.aa.aa
1.1.1.2 → FE2	bb.bb.bb.bb - → FE2	1.1.1.2 <-- →bb.bb.bb.bb

# Cisco SD-Access Fabric Architecture

## Unicast Forwarding in the same subnet



- 1 Host2 wants to communicate to Host1. Since they are in the same subnet Host2 sends an ARP request for the mac-address of Host1

The packet will contain below:

ARP REQUEST:

SRC IP: H2 IP

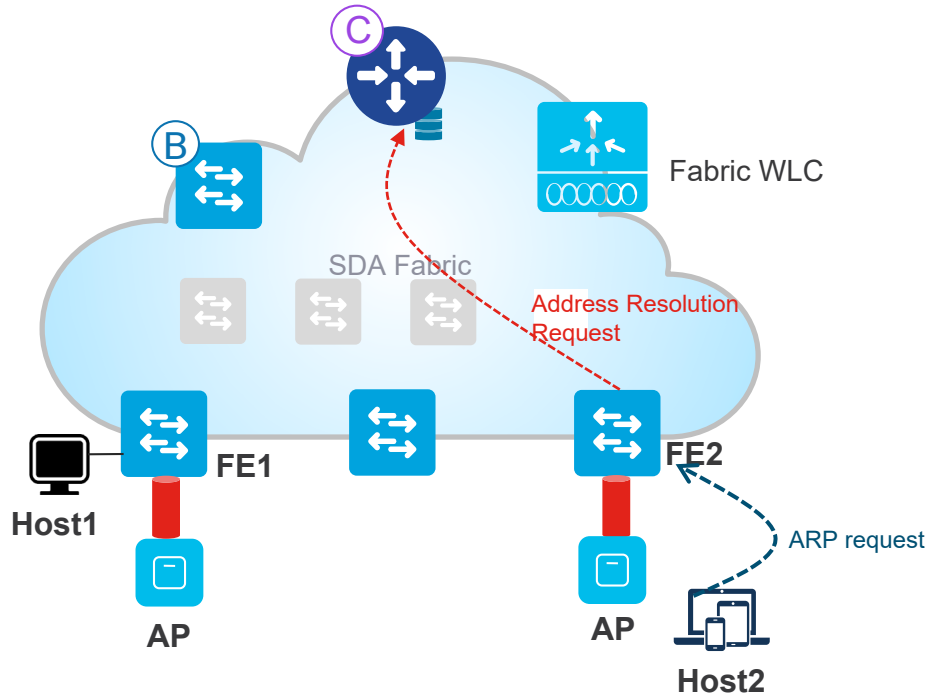
DST IP: H1 IP

SRC MAC: H2 MAC

DST MAC: FF:FF:FF:FF

# Cisco SD-Access Fabric Architecture

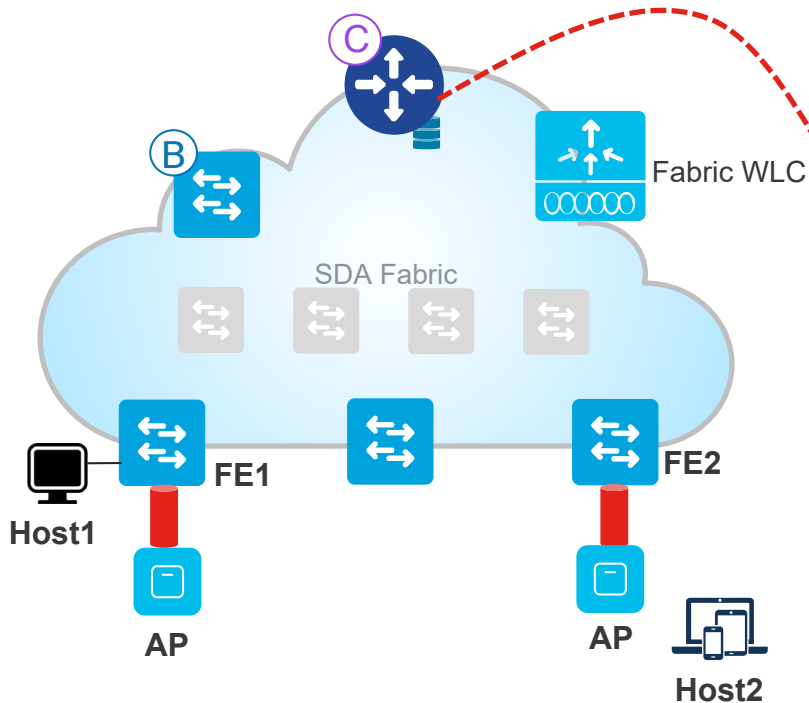
## Unicast Forwarding in the same subnet



- 2 Edge node FE2 will intercept the ARP request from Host2 and then will contact the control plane to ask for the MAC address of Host1.

# Cisco SD-Access Fabric Architecture

## Unicast Forwarding in the same subnet



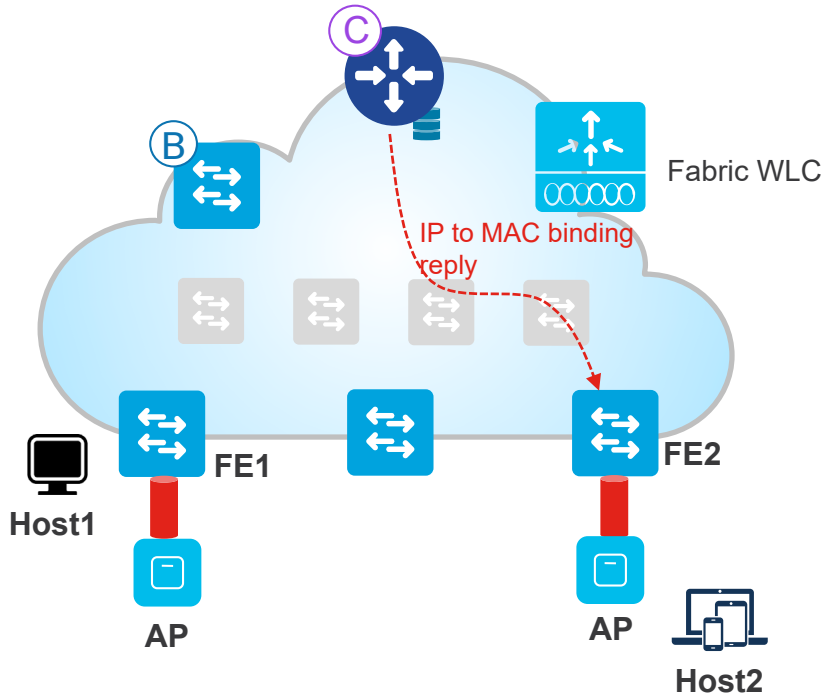
- 3 The control plane node will consult its Address resolution table to find the IP to MAC binding for Host1

Control Plane state:

IP to RLOC Table	MAC to RLOC Table	Address Resolution Table
1.1.1.1 → FE1	aa.aa.aa.aa - → FE1	1.1.1.1 <-- →aa.aa.aa.aa
1.1.1.2 → FE2	bb.bb.bb.bb - → FE2	1.1.1.2 <-- →bb.bb.bb.bb

# Cisco SD-Access Fabric Architecture

## Unicast Forwarding in the same subnet



- 4 If it finds the mapping it will reply back to the edge node with the MAC address of Host1.

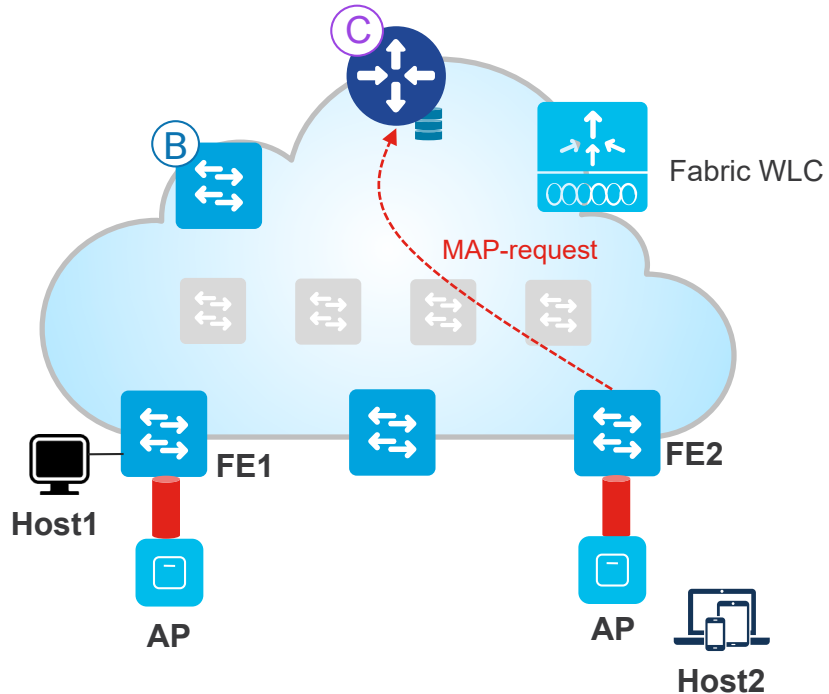
### Control Plane state:

IP to RLOC Table	MAC to RLOC Table	Address Resolution Table
1.1.1.1 → FE1	aa.aa.aa.aa - → FE1	1.1.1.1 <-- →aa.aa.aa.aa
1.1.1.2 → FE2	bb.bb.bb.bb - → FE2	1.1.1.2 <-- →bb.bb.bb.bb



# Cisco SD-Access Fabric Architecture

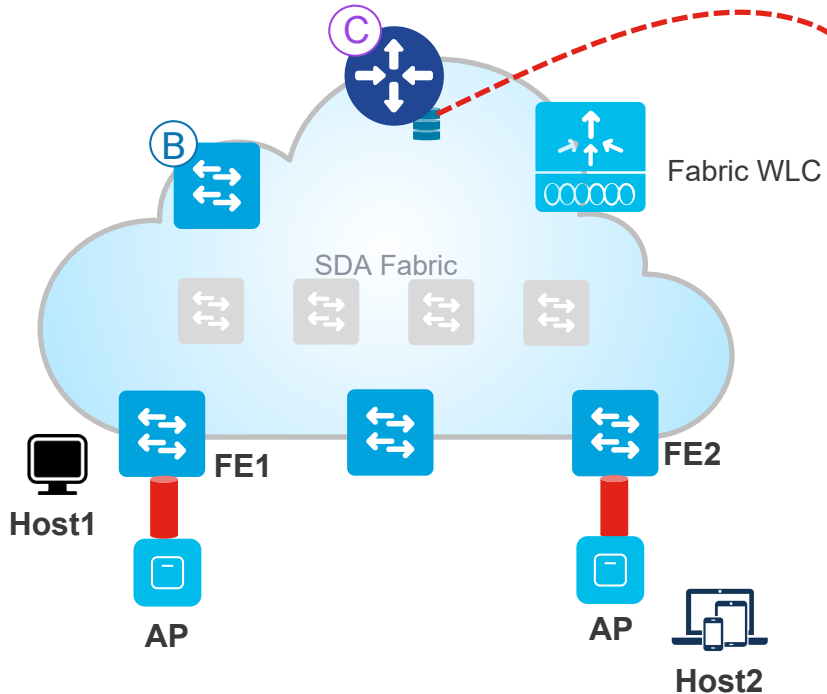
## Unicast Forwarding in the same subnet



- 5 Once the edge node FE2 gets the MAC address for Host1, it will again consult the control plane node to find out the location of Host1's MAC address.

# Cisco SD-Access Fabric Architecture

## Unicast Forwarding in the same subnet



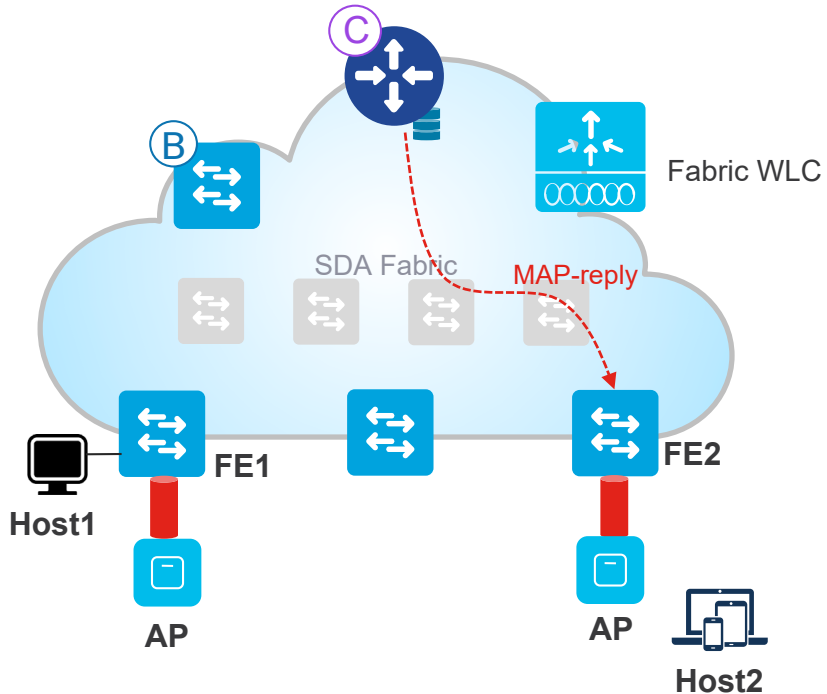
- 6 The control plane node after getting the map-request will consult its MAC to RLOC table for the location.

Control Plane state:

IP to RLOC Table	MAC to RLOC Table	Address Resolution Table
1.1.1.1 → FE1	aa.aa.aa.aa - → FE1	1.1.1.1 <-- → aa.aa.aa.aa
1.1.1.2 → FE2	bb.bb.bb.bb - → FE2	1.1.1.2 <-- → bb.bb.bb.bb

# Cisco SD-Access Fabric Architecture

## Unicast Forwarding in the same subnet



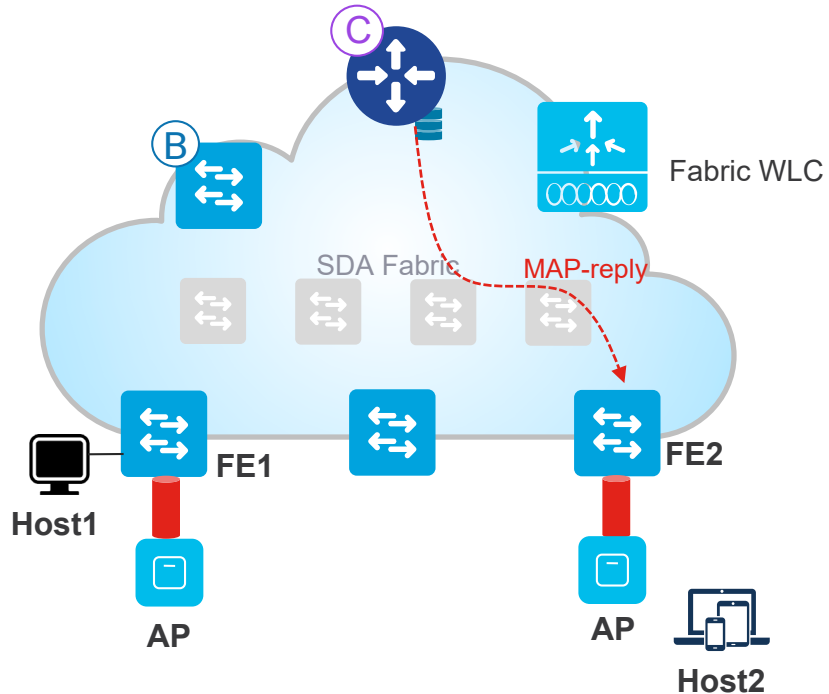
- 7 If the entry is found the information is sent to the edge node

### Control Plane state:

IP to RLOC Table	MAC to RLOC Table	Address Resolution Table
1.1.1.1 → FE1	aa.aa.aa.aa - → FE1	1.1.1.1 <-- → aa.aa.aa.aa
1.1.1.2 → FE2	bb.bb.bb.bb - → FE2	1.1.1.2 <-- → bb.bb.bb.bb

# Cisco SD-Access Fabric Architecture

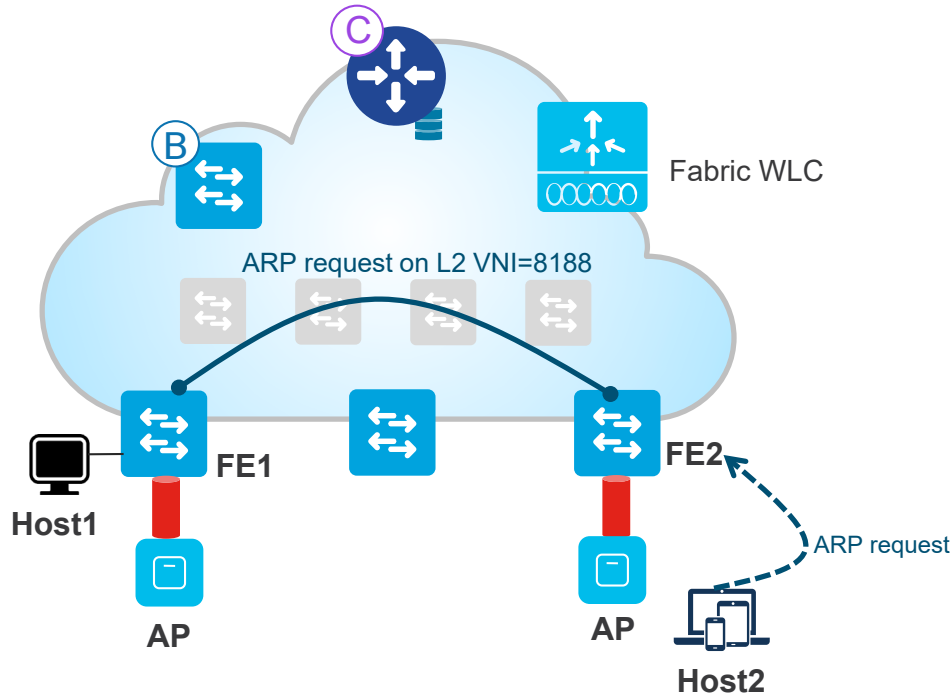
## Unicast Forwarding in the same subnet



- 8 The fabric edge node upon receiving the map-reply will install this entry into the Layer 2 forwarding table.

# Cisco SD-Access Fabric Architecture

## Unicast Forwarding in the same subnet



- 9 Edge node FE2 will now convert the ARP broadcast it received from Host2 to a directed unicast. It will send it to FE1 in the overlay encapsulating it with the L2VNI for that subnet/vlan.

The L2 VNI is derived from the L2 LISP configuration:  
instance-id 8188  
service ethernet  
eid-table vlan 1024

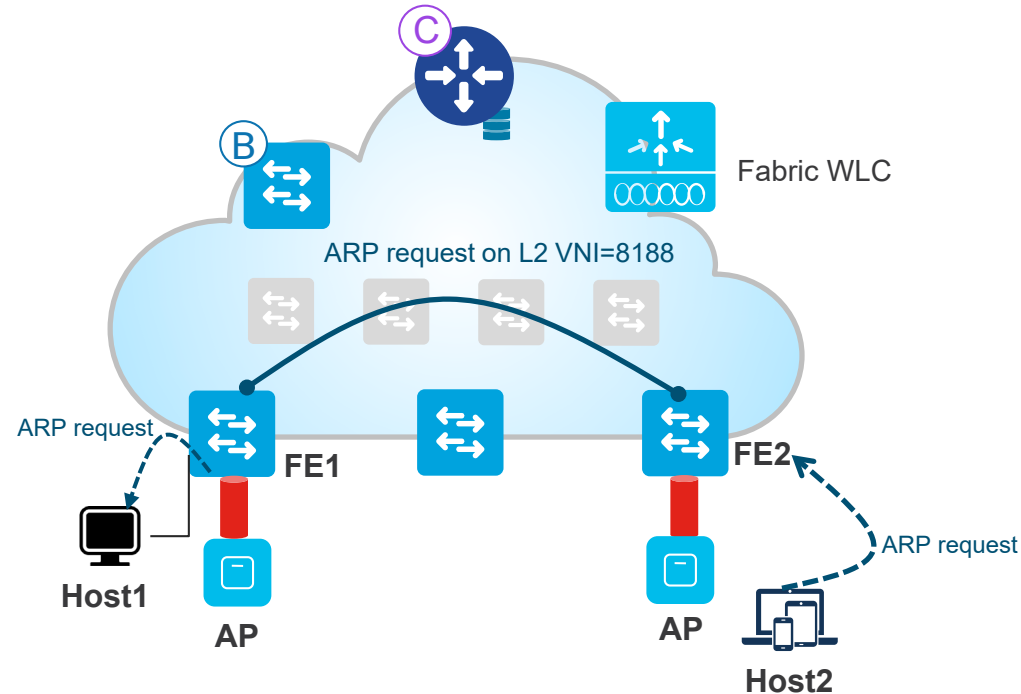
The packet will contain below:

ARP REQUEST:

SRC IP: H2 IP  
DST IP: H1 IP  
SRC MAC: H1 MAC  
DST MAC: H2 MAC << unicast ARP>>>

# Cisco SD-Access Fabric Architecture

## Unicast Forwarding in the same subnet

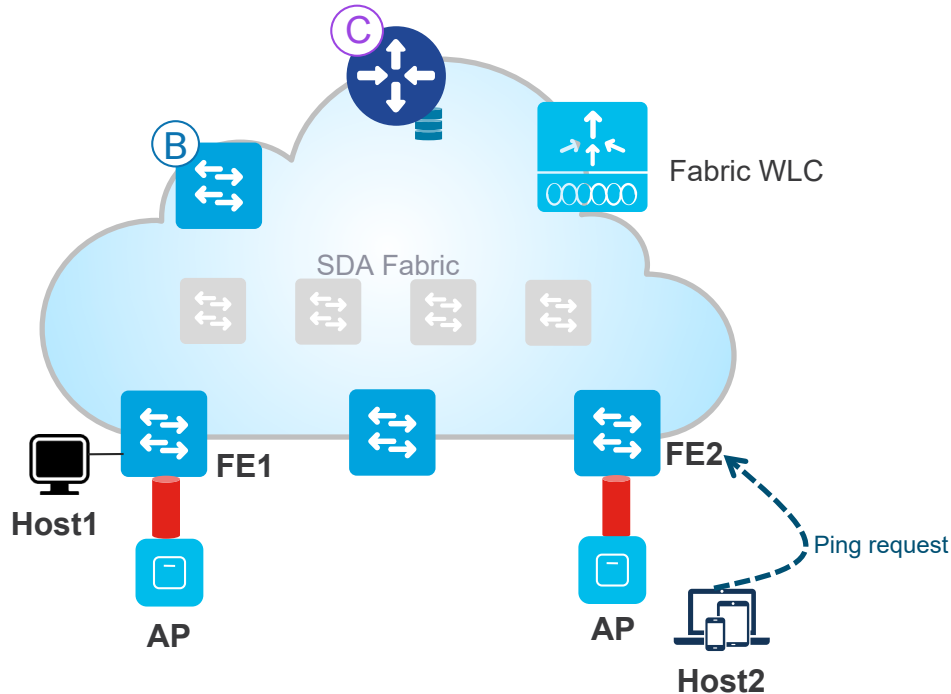


- 10 The edge node FE1 upon receiving the ARP request will decapsulate the packet. It will forward it to the destination host Host1.



# Cisco SD-Access Fabric Architecture

## Unicast Forwarding in the same subnet

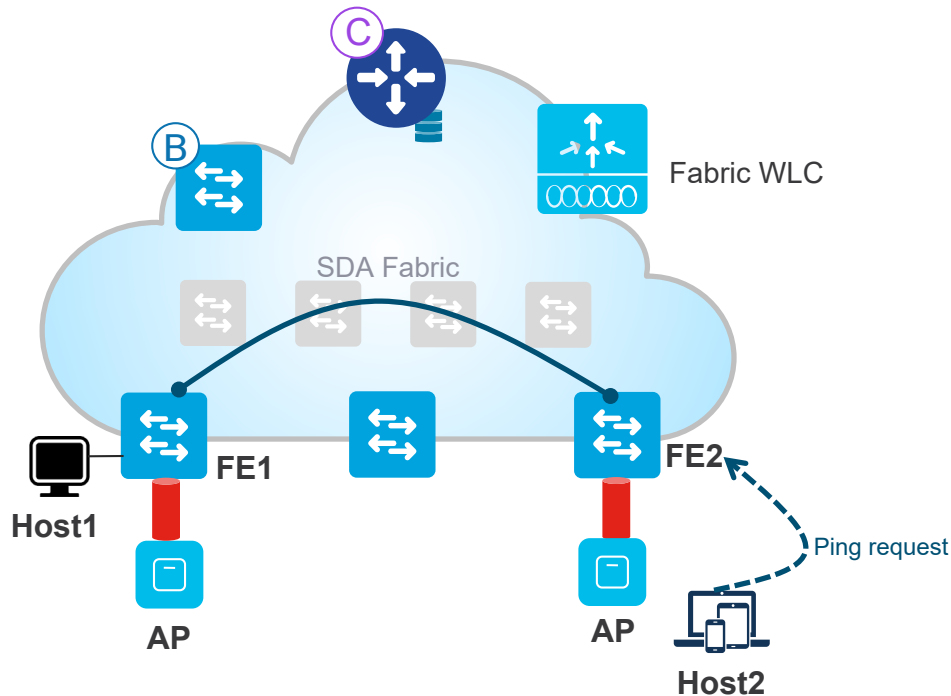


- 12 Now that ARP is resolved, we can test data connection (ping request) .



# Cisco SD-Access Fabric Architecture

## Unicast Forwarding in the same subnet



- 13 The fabric edge node will use the Layer 2 forwarding tables that have been already populated and will use the L2VNI to forward to the destination node.

INNER PACKET:

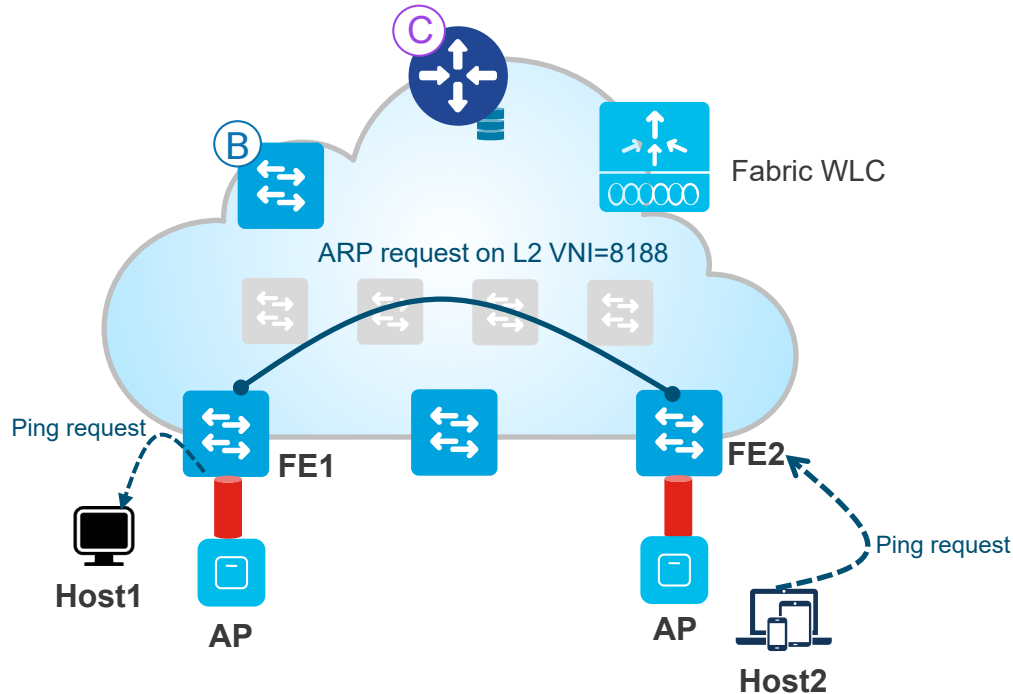
SRC IP: H2 IP  
DST IP: H1 IP  
SRC MAC: H2 MAC  
DST MAC: H1 MAC

OUTER PACKET:

SRC IP: Edge 2 IP  
SRC MAC: Edge 2 MAC  
DST IP: Edge 1 IP  
DST MAC: Next-Hop MAC (Intermediate node)

# Cisco SD-Access Fabric Architecture

## Unicast Forwarding in the same subnet

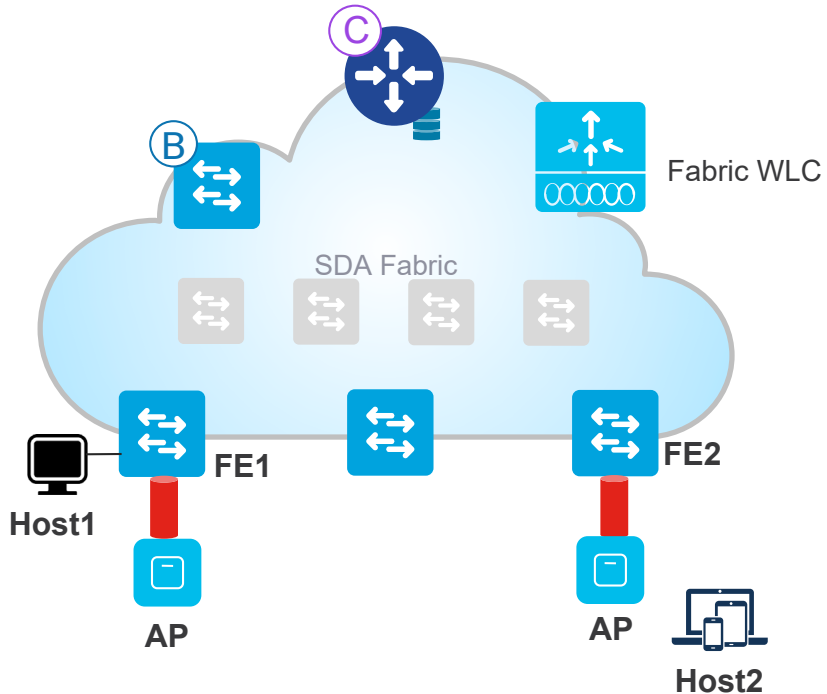


- 14 Once traffic reaches edge 1 it will decapsulate the VXLAN headers and based on the Inner packet details will do a switching lookup to send it to Host1.

# Forwarding across different subnets

# Cisco SD-Access Fabric Architecture

## Unicast Forwarding across different subnets



### Assumptions:

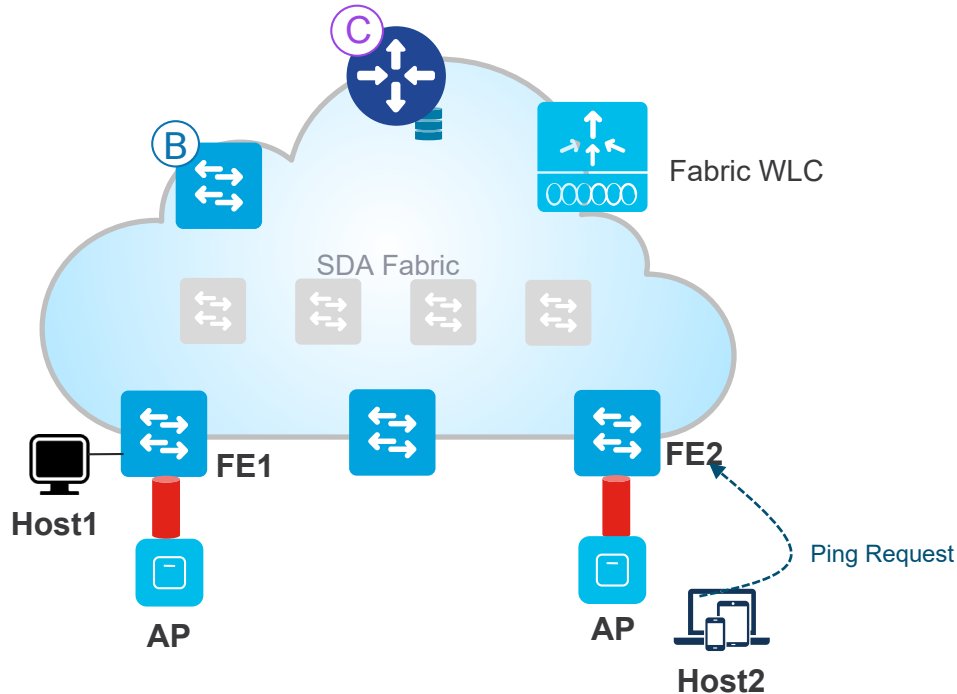
- Host H1 and Host H2 belong to different subnets

### Control Plane state:

IP to RLOC Table	MAC to RLOC Table	Address Resolution Table
1.1.1.1 → FE1	aa.aa.aa.aa - → FE1	1.1.1.1 <-- →aa.aa.aa.aa
2.2.2.2 → FE2	bb.bb.bb.bb - → FE2	2.2.2.2 <-- →bb.bb.bb.bb

# Cisco SD-Access Fabric Architecture

## Unicast Forwarding across different subnets



- 1 Host2 sends a ping packet to Host1.

The packet will contain below:

PING REQUEST:

SRC IP: H2 IP

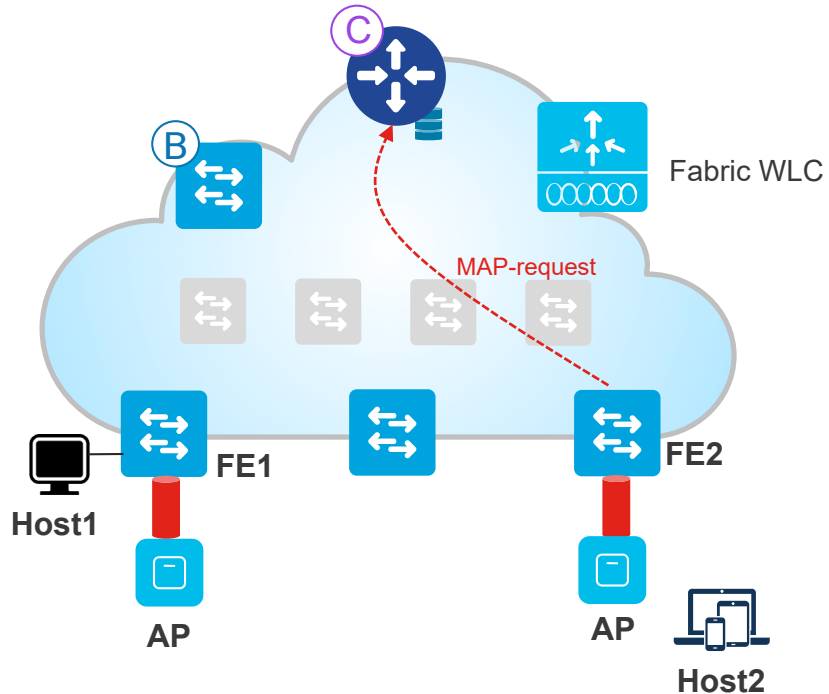
DST IP: H1 IP

SRC MAC: H2 MAC

DST MAC: DEFAULT GATEWAY MAC

# Cisco SD-Access Fabric Architecture

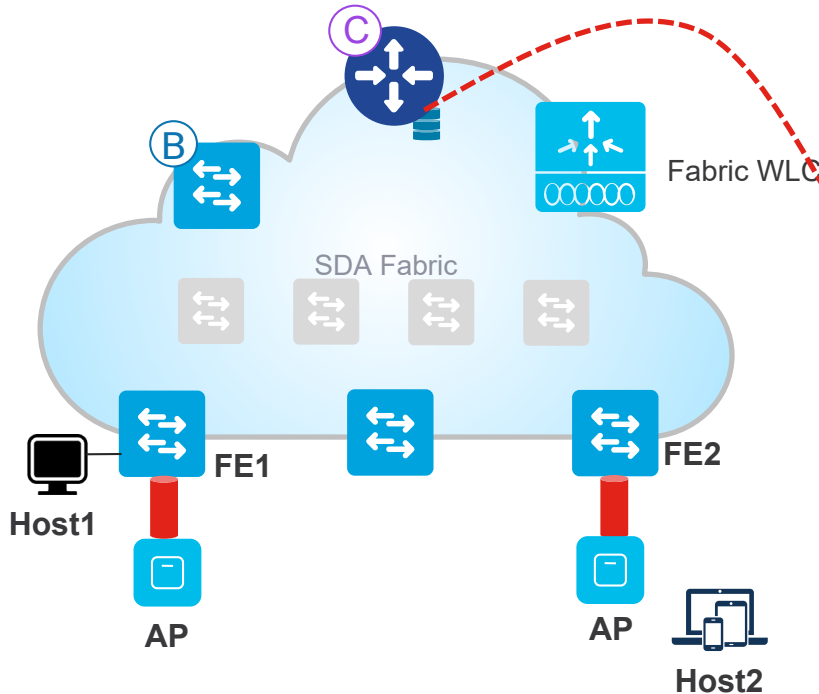
## Unicast Forwarding across different subnets



- 2 Edge node FE2 will intercept the ping request from Host2 destined to Host1 and then will contact the control plane to ask for the location of Host1.

# Cisco SD-Access Fabric Architecture

## Unicast Forwarding across different subnets



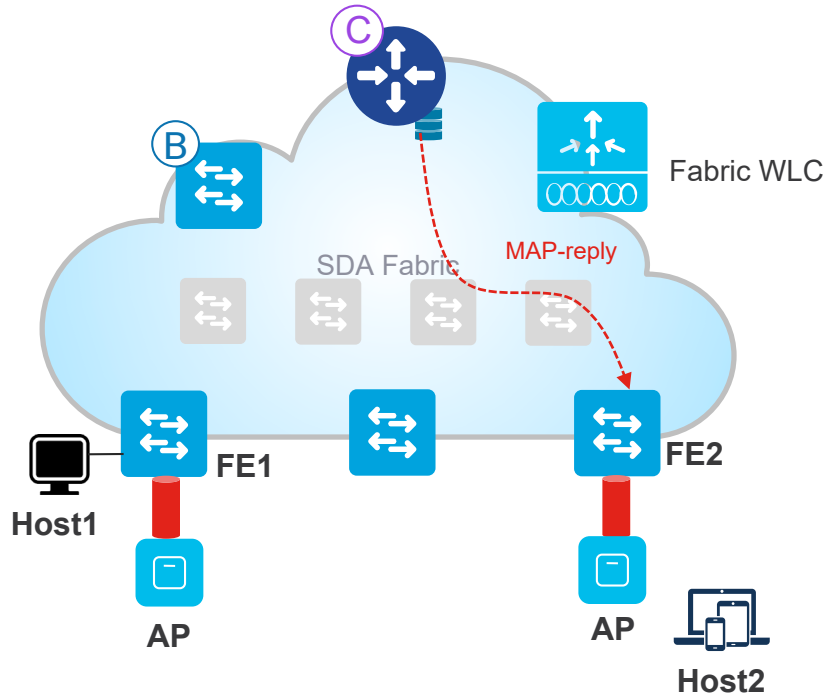
- 3 The control plane node will consult its IP to RLOC binding for location of Host1.

### Control Plane state:

IP to RLOC Table	MAC to RLOC Table	Address Resolution Table
1.1.1.1 → FE1	aa.aa.aa.aa - → FE1	1.1.1.1 <-- → aa.aa.aa.aa
2.2.2.2 → FE2	bb.bb.bb.bb - → FE2	2.2.2.2 <-- → bb.bb.bb.bb

# Cisco SD-Access Fabric Architecture

## Unicast Forwarding across different subnets



- 4 If it finds the mapping it will reply back to the edge node with the location of Host1

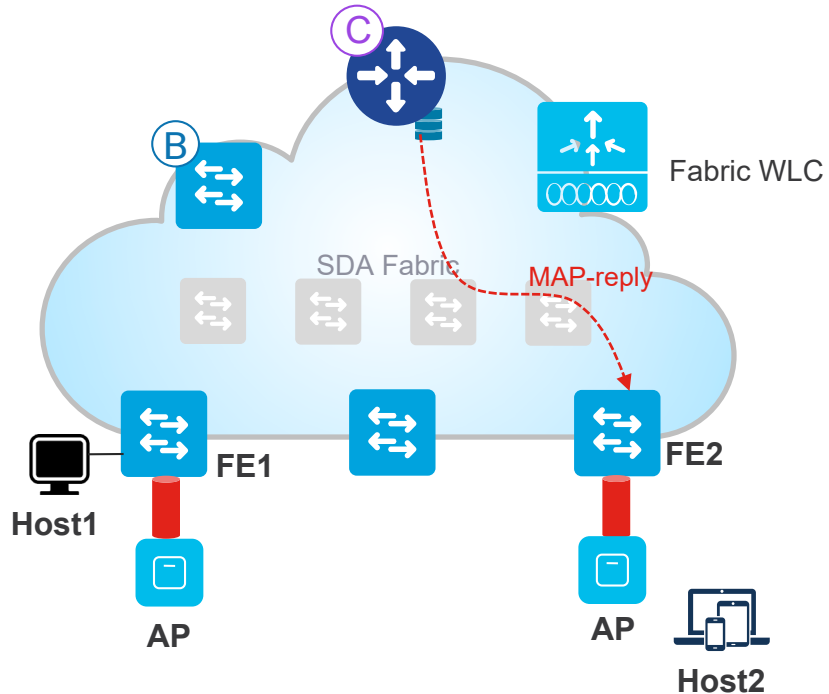
### Control Plane state:

IP to RLOC Table	MAC to RLOC Table	Address Resolution Table
1.1.1.1 → FE1	aa.aa.aa.aa - → FE1	1.1.1.1 <-- → aa.aa.aa.aa
2.2.2.2 → FE2	bb.bb.bb.bb - → FE2	2.2.2.2 <-- → bb.bb.bb.bb



# Cisco SD-Access Fabric Architecture

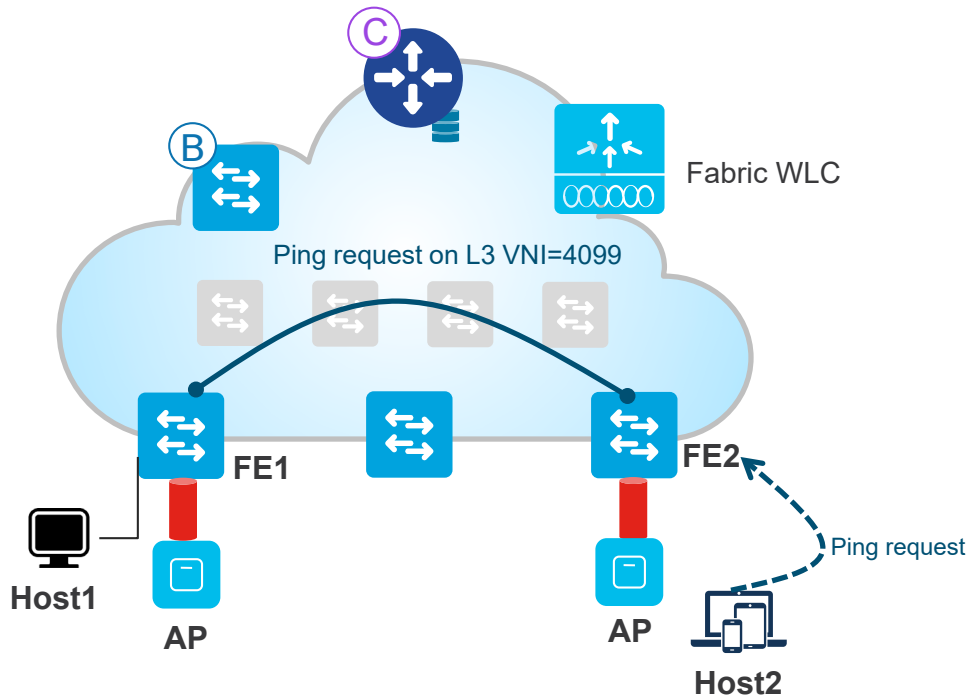
## Unicast Forwarding across different subnets



- 5 The fabric edge node upon receiving the map-reply will install this entry into the Layer 3 forwarding tables .

# Cisco SD-Access Fabric Architecture

## Unicast Forwarding across different subnets



- 6 Edge node FE2 will now send the ping request received from Host2 to edge node FE1 in the overlay encapsulating it with the L3VNI for that VRF.

The L3 VNI is derived from the LISP configuration:

```
instance-id 4099
```

```
remote-rlloc-probe on-route-change
```

```
service ipv4
```

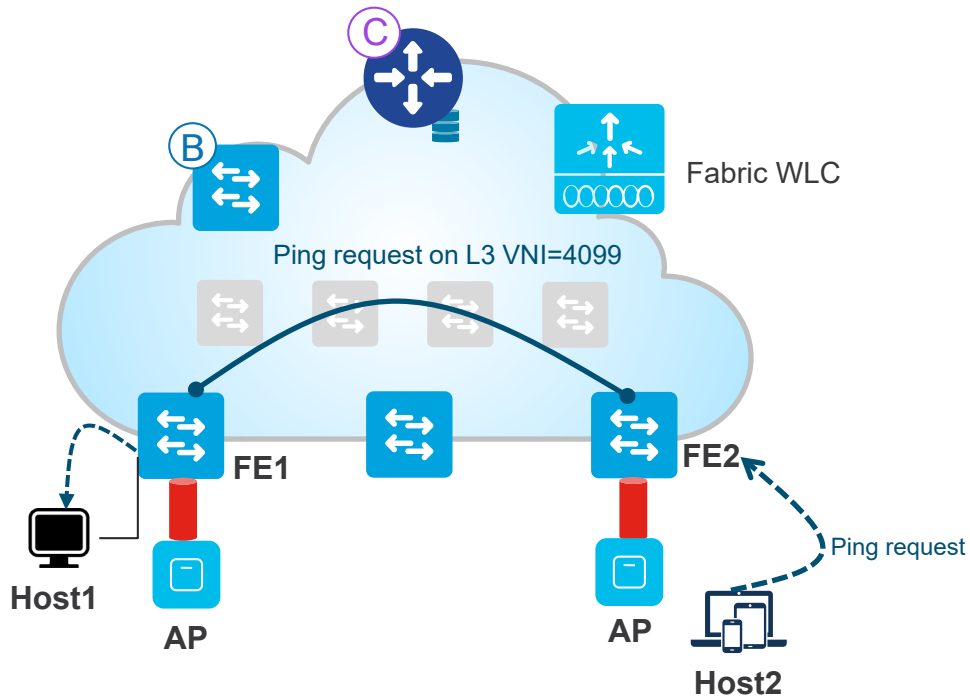
```
eid-table vrf Campus
```

```
map-cache 0.0.0.0/0 map-request
```

```
exit-service-ipv4
```

# Cisco SD-Access Fabric Architecture

## Unicast Forwarding across different subnets



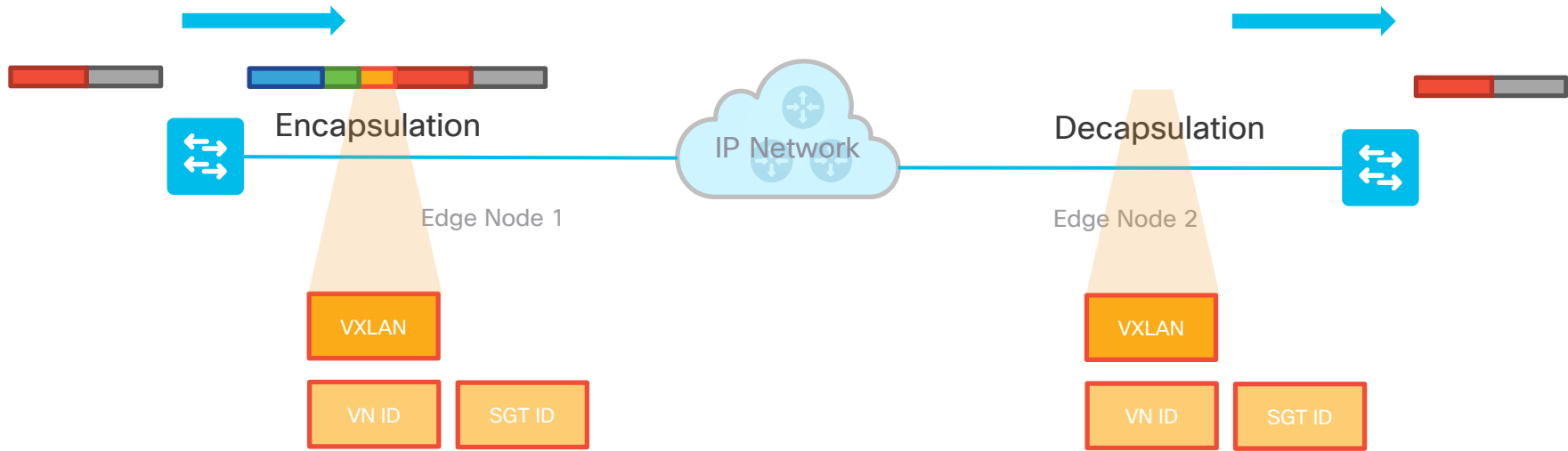
- 7 Once traffic reaches edge 1 it will decapsulate the VXLAN headers and based on the inner packet details will do a routing lookup to send it to Host1.

# Unicast Packet Forwarding & Access Control

1. Control Plane Lookup
2. Access Control Policy

# Group-Based Policy

## Ingress Classification & Egress Enforcement



### Classification

Static or Dynamic VN and SGT assignments



### Propagation

Carry VN and Group context across the network



### Enforcement

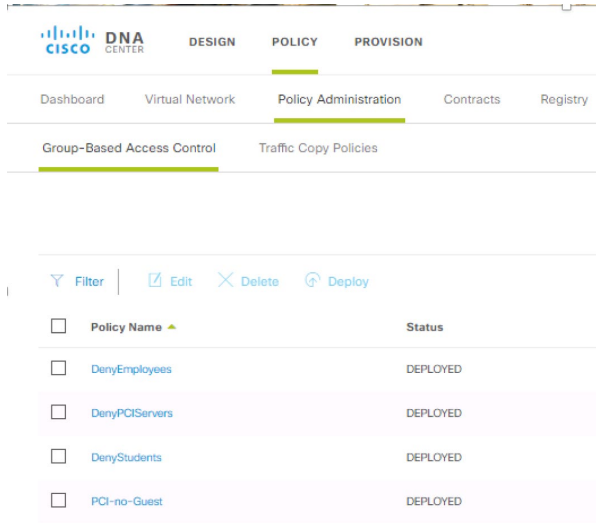
Group Based Policies  
ACLs, Firewall Rules

# Policy Definition in Cisco DNA Center

Policies in DNAC

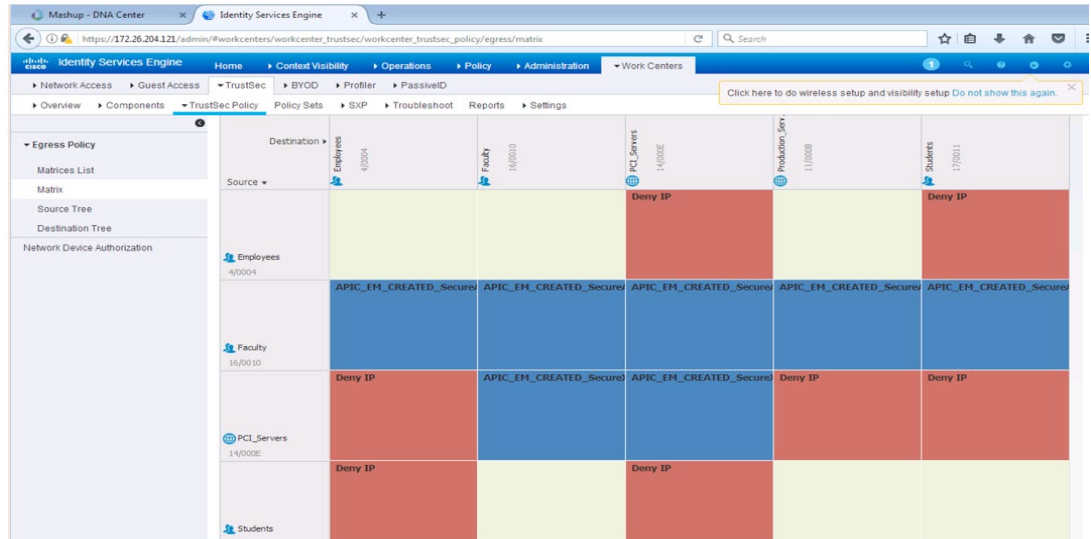
REST API

SGACL matrix in ISE for rendering (runtime download)



The screenshot shows the Cisco DNA Center interface. The top navigation bar includes 'DESIGN', 'POLICY', and 'PROVISION'. The 'POLICY' section is active, showing 'Policy Administration' and 'Traffic Copy Policies'. A table lists several policies with their status:

Policy Name	Status
DenyEmployees	DEPLOYED
DenyPCIServers	DEPLOYED
DenyStudents	DEPLOYED
PCI-no-Guest	DEPLOYED

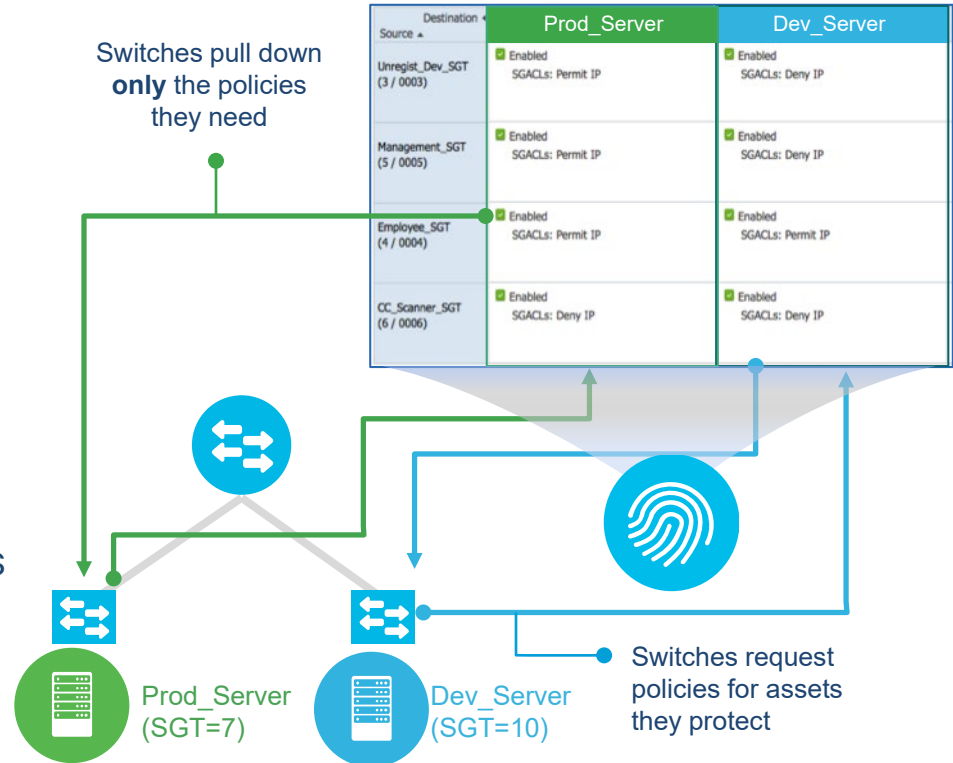


The screenshot shows the Identity Services Engine (ISE) interface displaying an Egress Policy matrix. The matrix is a grid where rows represent source groups and columns represent destination groups. The cells contain either 'Deny IP' (red) or 'APIC\_EH\_CREATED\_Secure' (blue).

Source	Employees (4/0004)	Faculty (16/0010)	PCI_Servers (14/000E)	Production_Serv (11/0008)	Students (17/0011)
Employees (4/0004)	Deny IP	Deny IP	Deny IP	Deny IP	Deny IP
Faculty (16/0010)	APIC_EH_CREATED_Secure	APIC_EH_CREATED_Secure	APIC_EH_CREATED_Secure	APIC_EH_CREATED_Secure	APIC_EH_CREATED_Secure
PCI_Servers (14/000E)	Deny IP	APIC_EH_CREATED_Secure	APIC_EH_CREATED_Secure	Deny IP	Deny IP
Students (17/0011)	Deny IP	Deny IP	Deny IP	Deny IP	Deny IP

# Dynamic Policy Download

- New User/Device/Server provisioned
- Switch requests policies for assets they protect
- Policies downloaded & applied dynamically
- Result: All controls centrally managed
  - Security policies de-coupled from network topology
  - No switch-specific security configs needed
  - One place to audit network-wide policies



# Downloaded Policy

```
Switch#show cts rbacl Permit_Email_Traffic
CTS RBACL Policy
=====
```

```
RBACL IP Version Supported: IPv4
name      = Permit_Email_Traffic-40
IP protocol version = IPV4
refcnt = 1
flag     = 0x40000000
stale    = FALSE
```

RBACL ACEs:

```
permit tcp dst eq 110
permit tcp dst eq 143
permit tcp dst eq 25
permit tcp dst eq 465
permit tcp dst eq 585
permit tcp dst eq 993
permit tcp dst eq 995
deny all log
```

ISE

Destination	Source	Policy
Contractors 30/001E	Contractors 30/001E	Cisco_Jabber_Access
Employee_FullAc... 10/000A	Employee_FullAc... 10/000A	Malware_Control_ACL
Mail_Servers 120/0078	Contractors 30/001E	Permit_Email_Traffic
Web_Servers 110/000E	Employee_FullAc... 10/000A	

Default  Enabled SGACLs : Permit IP Description : Default egress rule

Switch

```
Switch#show cts role-based permissions
```

```
IPv4 Role-based permissions default:
  Permit IP-00
...
IPv4 Role-based permissions from group 10:Employee_FullAccess to group
10:Employee_FullAccess:
  Malware_Control_ACL-10
IPv4 Role-based permissions from group 10:Employee_FullAccess to group 30:Contractors:
  Cisco_Jabber_Access-10
IPv4 Role-based permissions from group 30:Contractors to group 10:Employee_FullAccess:
  Cisco Jabber Access-10
IPv4 Role-based permissions from group 30:Contractors to group 120:Mail_Servers:
  Permit_Email_Traffic
...
```

IOS switch as enforcer



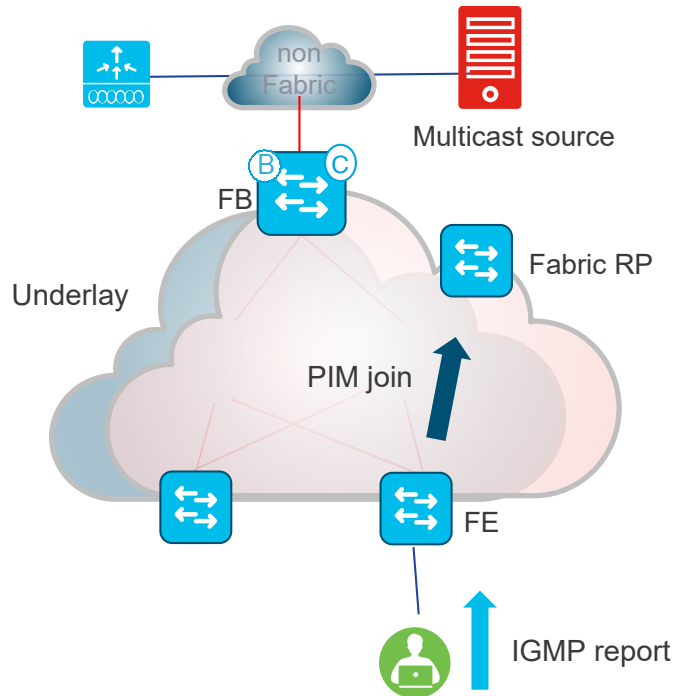
# Advanced Topics

1. Multicast modes
2. Broadcast forwarding

# PIM ASM/SSM Control plane Interaction with the Fabric

# Cisco SD-Access Fabric Architecture

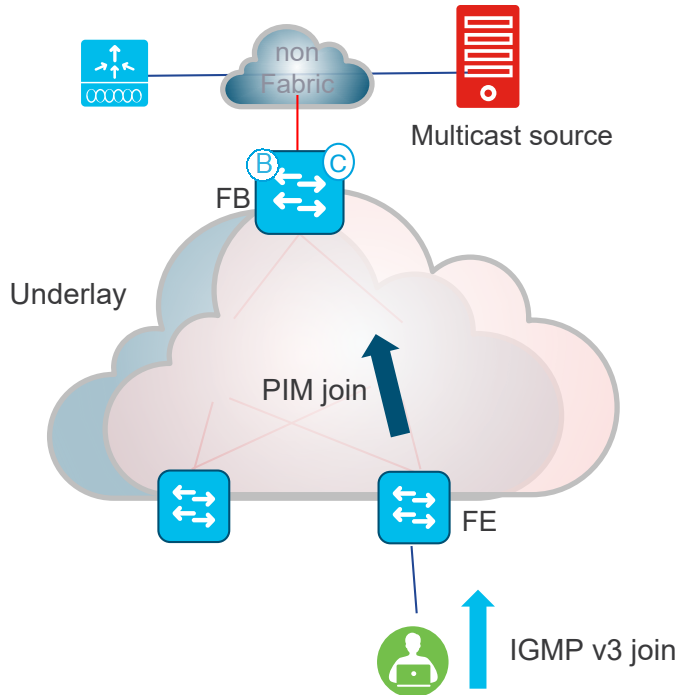
## Multicast with PIM-ASM – Multicast Receiver to RP



- Multicast client (receiver) is in the overlay, multicast source can be outside or inside the fabric
- PIM-ASM or SSM can be running in the overlay
- The client sends an IGMP report for a specific multicast group (G)
- The fabric Edge node (FE) sends a PIM join towards the Rendezvous Point RP
- The RP is registered as part of the end point IP (EID) space of the overlay.
- The edge node will ask the control plane for the location of the RP address (IP to RLOC Table) and based on the reply will send the PIM join in the overlay to the RP.

# Cisco SD-Access Fabric Architecture

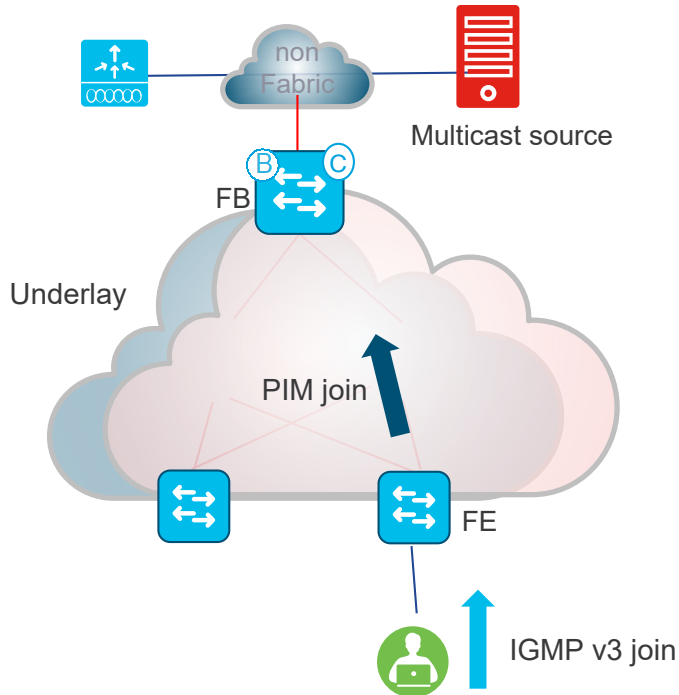
## How multicast works with PIM SSM



- Multicast client (receiver) is in the overlay, multicast source can be outside the fabric or in the overlay as well
- PIM-SSM needs to be running in the overlay
- An RP is not used in a PIM SSM deployment
- The client sends an IGMP v3 report for a specific multicast group (G)
- The fabric edge node (FE) receives it and since the IGMP v3 report has the source address information for that multicast group it sends a PIM join towards the source directly. In our case since the source is reachable through the border it sends the PIM join to the border.

# Cisco SD-Access Fabric Architecture

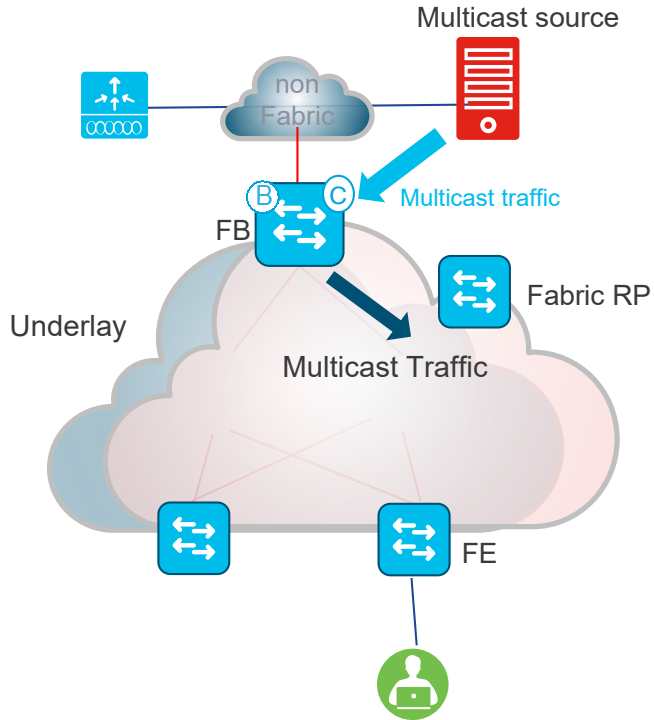
## How multicast works with PIM SSM – Multicast Receiver to RP



- Since in a SSM deployment the source address is part of IGMP v3 join the edge node will ask the control plane for the location of the source address (IP to RLOC Table) and based on the reply will send the PIM join in the overlay to the destination node.
- If Border registered that source then the PIM join is directly sent to the Border (the Border may forward the PIM join upstream towards the source, if the source is not directly connected).
- If the source is not known in the fabric, the PIM join is also sent to the border as it is the default exit point of the fabric.

# Cisco SD-Access Fabric Architecture

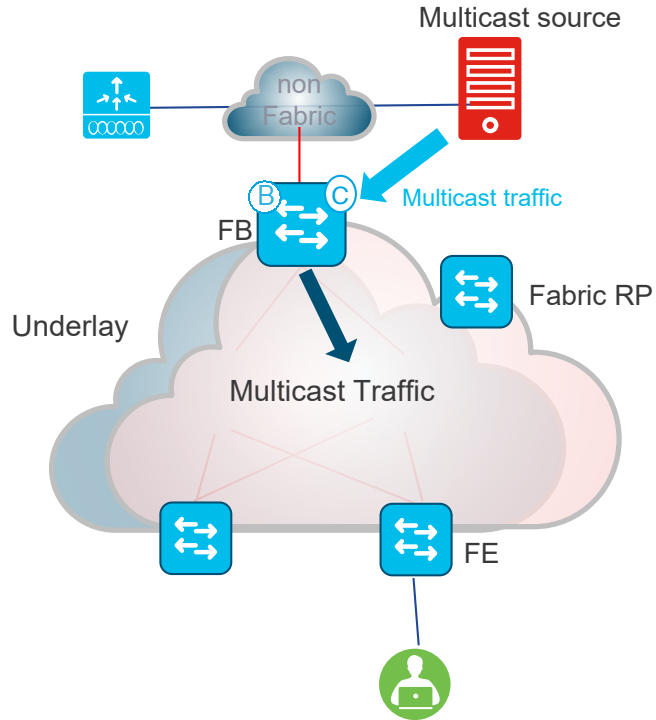
## How multicast works with PIM ASM – Multicast Source to RP



- The multicast source will send the multicast traffic on the interfaces towards the RP via the fabric Border(FB)
- If the Source is directly connected to the border, the border is the DR for that segment
- If the Source is not directly connected, the border is simply along the routed path towards the RP
- The FB receives it and sends the traffic towards the RP
- The Border node will ask the control plane for the location of the RP address (IP to RLOC Table) and based on the reply will send the traffic in the overlay to the RP
- The RP now has the source and receiver information for that multicast group

# Cisco SD-Access Fabric Architecture

## How multicast works with PIM SSM – Multicast Source to RP



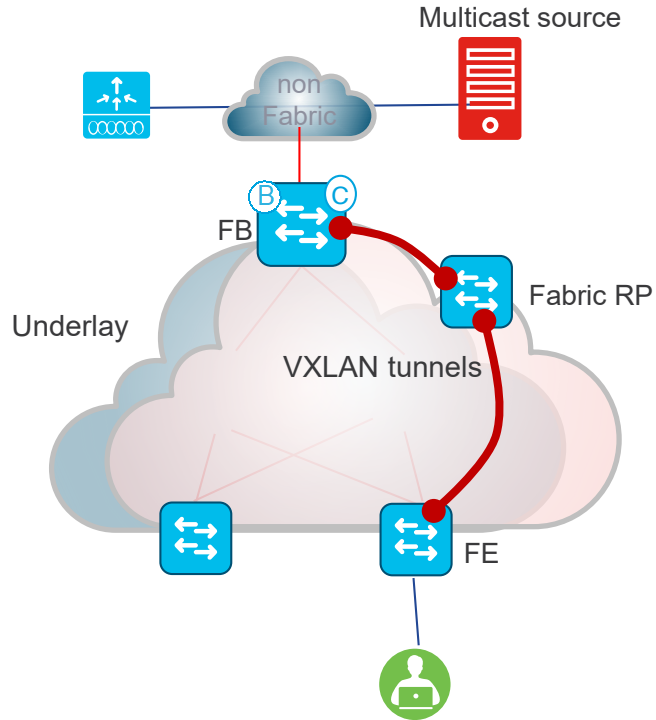
- The multicast source will send the multicast traffic on the interfaces towards the listeners via the fabric Border(FB)
- If the Source is directly connected to the border, the border is the DR for the segment
- If the Source is not directly connected, an SSM tree was formed when the PIM joins where forwarded previously and the Fabric Border is on the path of that tree.
- The FB receives it and sends the traffic towards the fabric Edge(s) as the PIM join is directly coming from the fabric Edge(s) to the Border in a SSM deployment.
- In a PIM SSM deployment there isn't an RP anchored shared tree .

# Head End Replication Multicast Data Plane in Fabric



# Cisco SD-Access Fabric Architecture

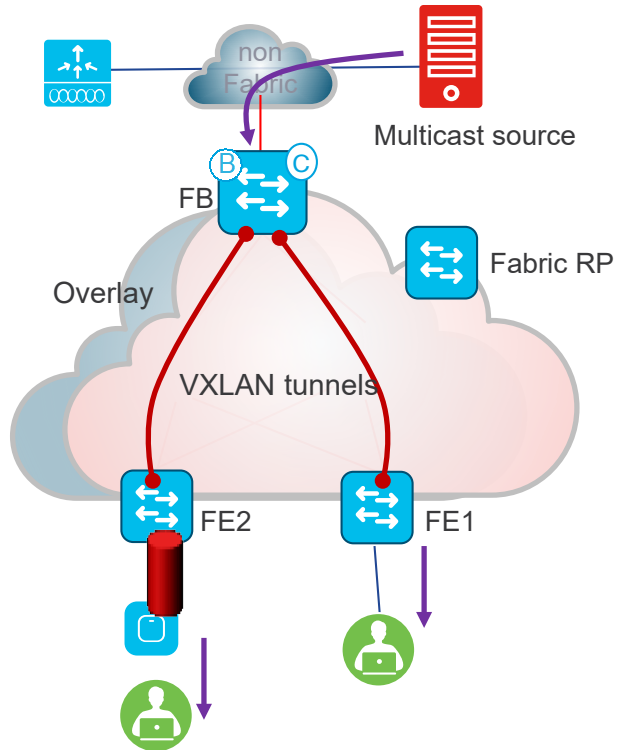
## How multicast works with PIM ASM – Data Plane



- The RP has the source and receiver information for a particular multicast group
- The fabric Border will send the multicast source traffic over a VXLAN tunnel to the RP. The RP will forward that traffic to FE over another VXLAN tunnel
- FE receives the VXLAN packets, decapsulates, applies policy and sends original IP multicast packet to the port on which the receiver is connected

# Cisco SD-Access Fabric Architecture

## Multicast using PIM-ASM – Data Plane

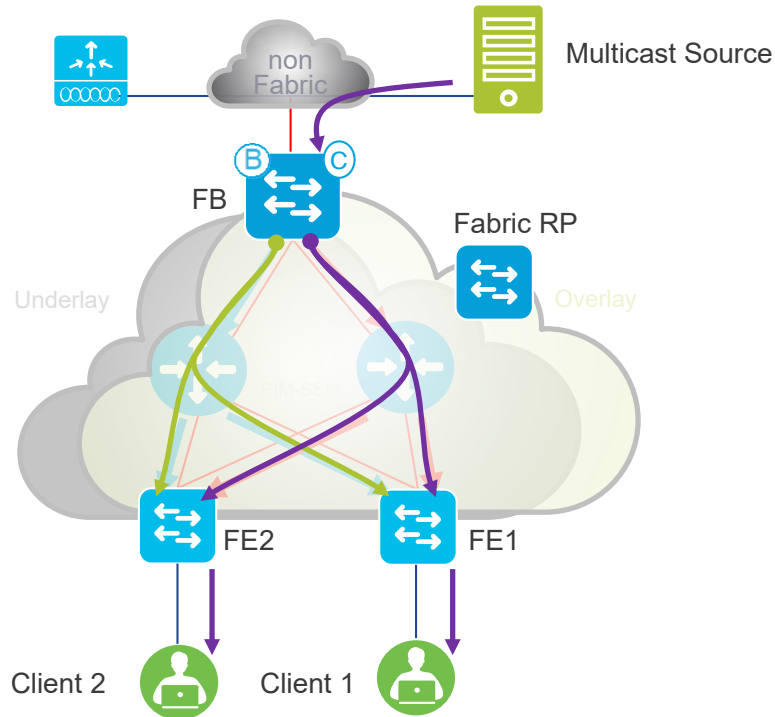


- Once the first multicast packet is delivered to the FE, the shortest path failover (SPT) takes place. Traffic is forwarded directly between the Border and the Edge.
- The FE knows that the Border owns the multicast source based on the first multicast packet received and send a PIM join directly to the Border for that multicast group.
- FB now knows which FEs have clients that requested the specific multicast group.
- It performs headend replication and VXLAN encapsulates the multicast traffic and unicasts it to the interested FEs
- The multicast traffic is sent in the overlay
- FE receives the VXLAN packets, decapsulates, policy and then sends original IP multicast packet to the port on which the receiver is connected.

# Native Multicast

# Cisco SD-Access Fabric Architecture

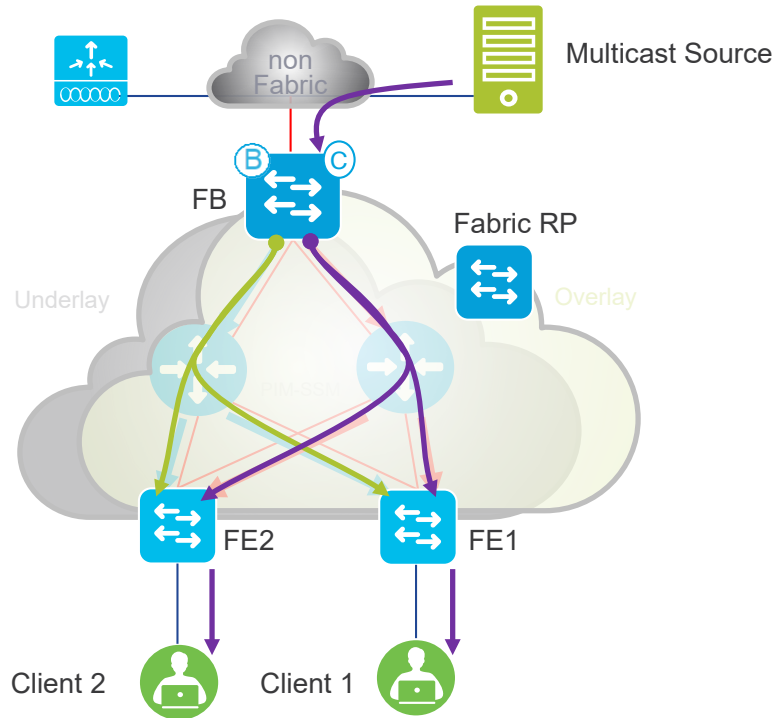
## Native Multicast



- ✓ Significantly reduces replication load at the Head-End
- ✓ Significantly improves overall scale and reduces latency

# Cisco SD-Access Fabric Architecture

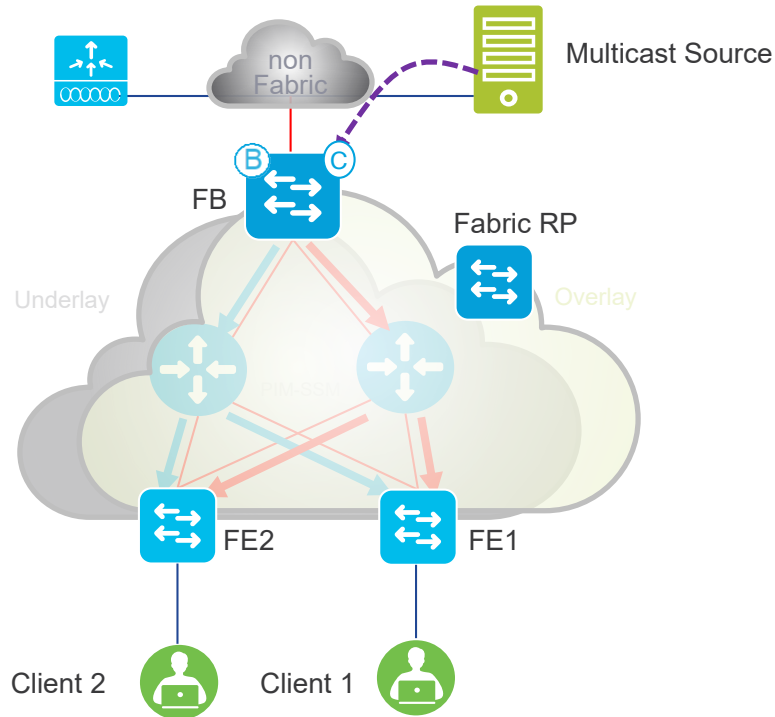
## Native Multicast



- All existing multicast control plane overlay behavior is the same
- PIM ASM and SSM can be used in the overlay as before.
- Each multicast group in the overlay is mapped to a corresponding (PIM SSM) multicast underlay group
- Multicast distribution (replication) occurs natively within the underlay network (e.g. intermediate nodes)
- Incoming multicast traffic for a given VN is encapsulated in VXLAN, and then sent with {Source IP = FE node RLOC, Destination IP = Underlay Multicast Group} as the outer IP addresses.
- PIM SSM is used in the underlay for multicast transport

# Cisco SD-Access Fabric Architecture

## Native Multicast



When the native multicast knob is turned on for a given fabric site the VNs where multicast is turned on will be instructed to move over to native multicast for the data path.

The configuration is pushed under the LISP interface for the respective VNs and the multicast groups in that VN will be mapped to underlay SSM groups for data transport.

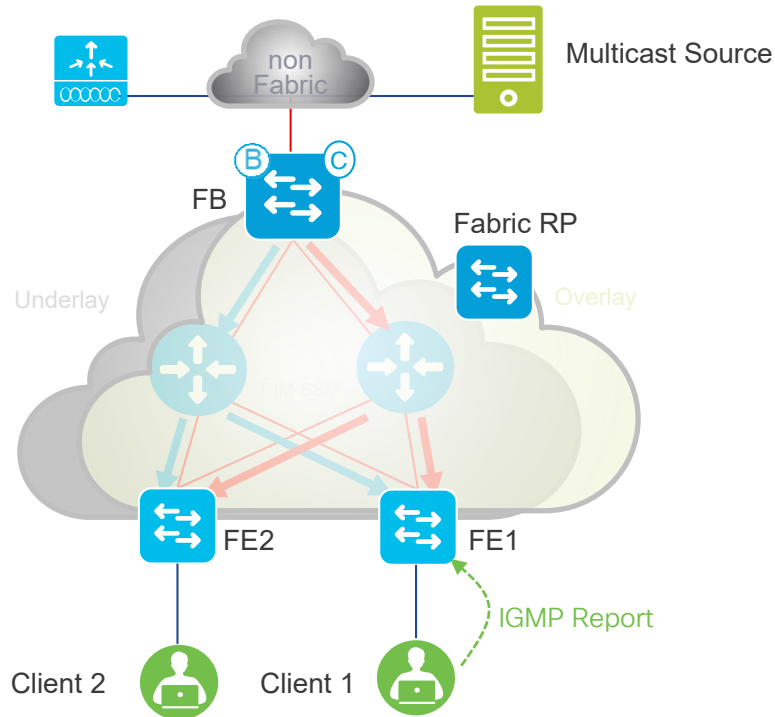
Interface LISP0.4096

```
ip pim lisp transport multicast
```

```
ip pim lisp core-group-range 232.0.0.1 1000
```

# Cisco SD-Access Fabric Architecture

## Native Multicast



1

Client 1, a multicast receiver sends an IGMP report for group 238.0.0.1 to the fabric Edge

In this example we assume that ASM is used in the overlay and the group address is 238.0.0.1.

The overlay ASM group is mapped to SSM group (RP-RLOC, 232.0.0.9) in the underlay.

This is derived based on the configuration

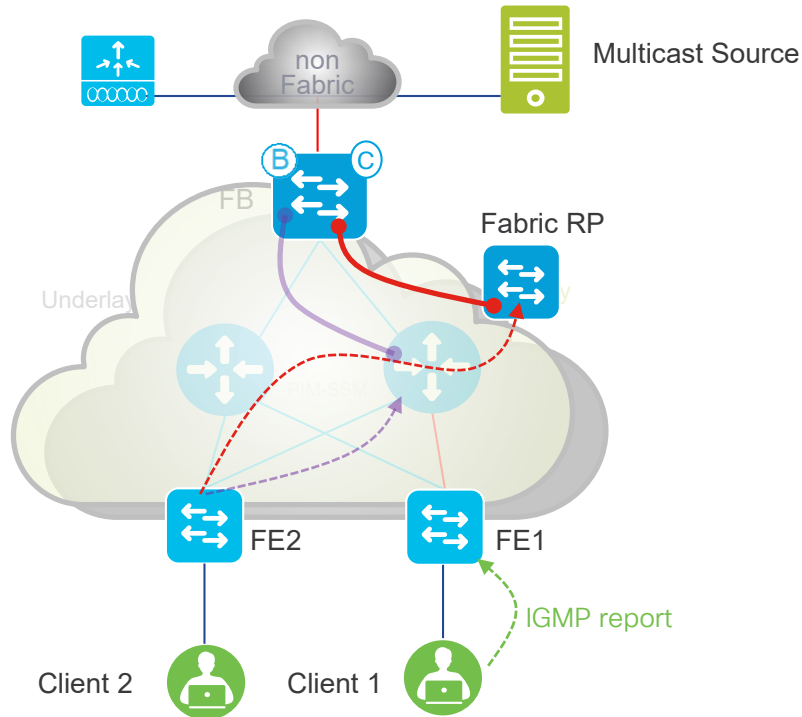
Since we are using SSM in the underlay for native multicast there is no pre built multicast tree for any given group in overlay.





# Cisco SD-Access Fabric Architecture

## Native Multicast



3

The multicast source starts sending traffic.

The overlay multicast traffic is unicast encapsulated to the RP-RLOC, it is then multicast encapsulated in the (S,G) tree from the RP-RLOC onwards.

The iTR (in our case the fabric Border) will send a source registration message in the overlay on the group address 238.0.0.1 to the RP and also sends the 238.0.0.1 multicast traffic in the overlay to the RP-RLOC (unicast encapsulated).

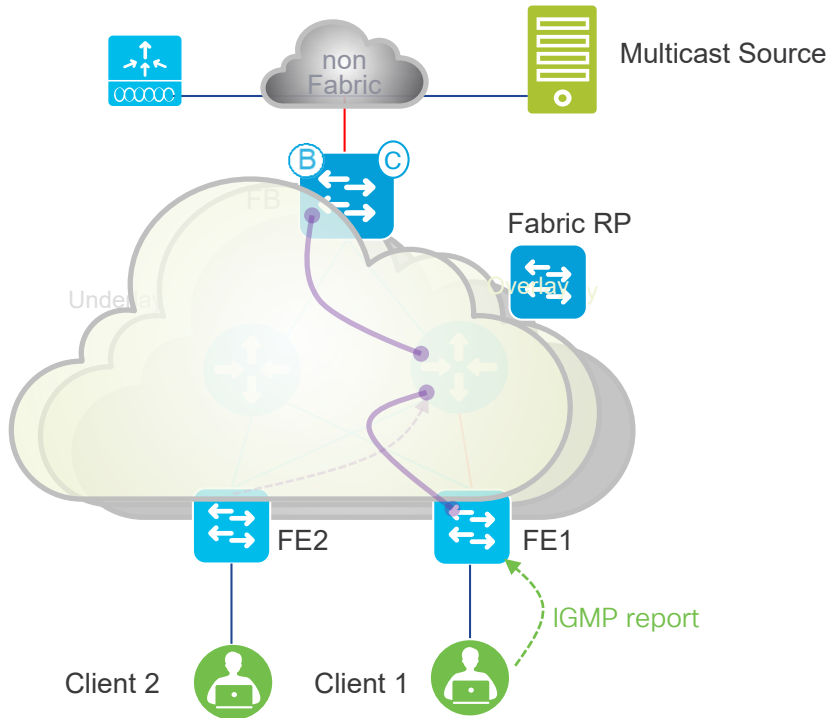
The RP-RLOC xTR sends the traffic in the underlay on the mapped group 232.0.0.9.

The information of the RLOC address is provided by control plane node. The control plane gives the IP to RLOC information for the RP address.

Up to this point an RP anchored ASM tree has been created

# Cisco SD-Access Fabric Architecture

## Native Multicast



4

The underlay now has enough information to replicate the traffic to the needed devices.

When SPT failover takes place, the source specific PIM joins in the overlay are mapped to source specific PIM joins in the underlay and an SSM SPT to the actual source is formed in the underlay.

Pre-SPT: Underlay (RP-RLOC, 232.0.0.9)

Post SPT: Underlay (Source-RLOC, 232.0.0.9)

# Summary

# SD-Access Support

## Digital Platforms for your Cisco Digital Network Architecture



For more details: [cs.co/sda-compatibility-matrix](https://cs.co/sda-compatibility-matrix)

### Switching

Catalyst 9600



Catalyst 9400



Catalyst 9500



Catalyst 9300



Catalyst 9200



Catalyst 4500E



Catalyst 6800



Nexus 7700



Catalyst 3850 & 3650

### Routing

ASR-1000-HX



ASR-1000-X



ISR 4451



ISR 4430



ISR 4330



ENCS 5400



### Wireless

Catalyst 9800



Catalyst 9100 APs



AIR-CT8540



AIR-CT3504



AIR-CT5520



Aironet Wave 1 APs\*



Aironet Wave 2 APs

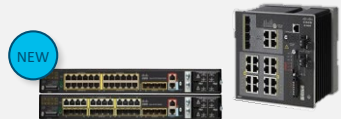
### Extended <sup>BETA</sup>



Cisco Digital Building



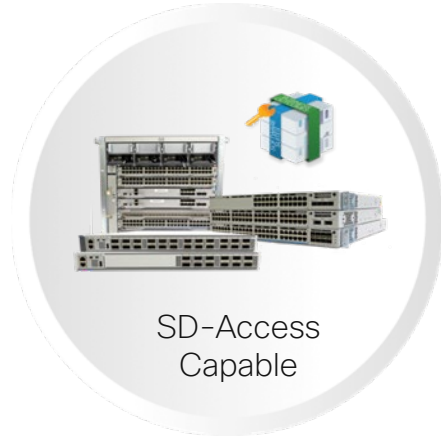
Catalyst 3560-CX



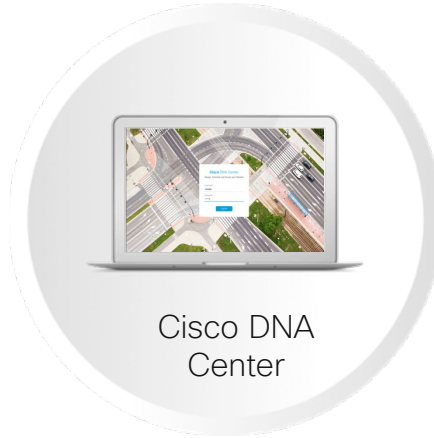
Cisco IE 4K/5K

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SD-Access  
Capable



Cisco DNA  
Center



Cisco  
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with **Cisco DNA Advantage Licenses**

Deploy the  
Cisco DNA Center

Get **Cisco DNA Center Appliances**  
with **Cisco DNA Center Software**

Engage with  
Cisco Services

**Cisco Services** can help you  
**Test - Migrate - Deploy - Manage**

# SD-Access Resources

Would you like to know more?



[cisco.com/go/dna](https://cisco.com/go/dna)

[cisco.com/go/sdaccess](https://cisco.com/go/sdaccess)

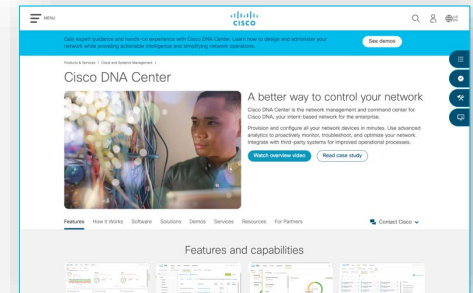
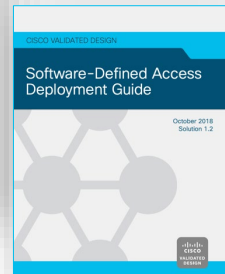
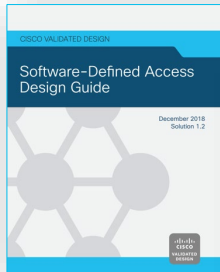
- [SD-Access At-A-Glance](#)
- [SD-Access Ordering Guide](#)
- [SD-Access Solution Data Sheet](#)
- [SD-Access Solution White Paper](#)

[cisco.com/go/cvd](https://cisco.com/go/cvd)

- [SD-Access Design Guide](#)
- [SD-Access Deployment Guide](#)
- [SD-Access Segmentation Guide](#)

[cisco.com/go/dnacenter](https://cisco.com/go/dnacenter)

- [Cisco DNA Center At-A-Glance](#)
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- [Cisco DNA Center Data Sheet](#)
- [Cisco DNA Center 'How To' Video Resources](#)





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# More Resources

Understand the underpinnings of your SDA fabric

Understand the potential of your LISP enabled network



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**Victor Moreno, CCIE No. 6908**

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